

CprE 281: Digital Logic

Instructor: Alexander Stoytchev

<http://www.ece.iastate.edu/~alexs/classes/>

Signed Numbers

CprE 281: Digital Logic
Iowa State University, Ames, IA
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Administrative Stuff

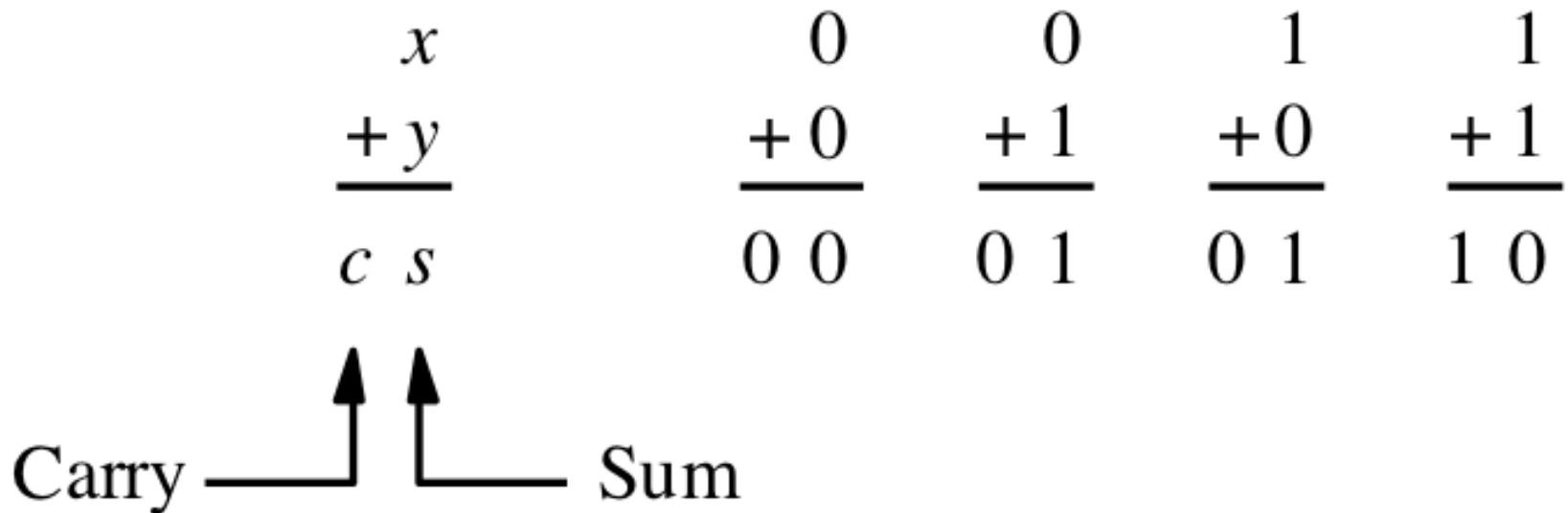
- **HW5 is out**
- **It is due on Monday Oct 5 @ 4pm.**
- **Please write clearly on the first page (in block capital letters) the following three things:**
 - **Your First and Last Name**
 - **Your Student ID Number**
 - **Your Lab Section Letter**

Administrative Stuff

- **Labs Next Week**
- **Mini-Project**
- **This one is worth 3% of your grade.**
- **Make sure to get all the points.**
- **http://www.ece.iastate.edu/~alexs/classes/2015_Fall_281/labs/Project-Mini/**

Quick Review

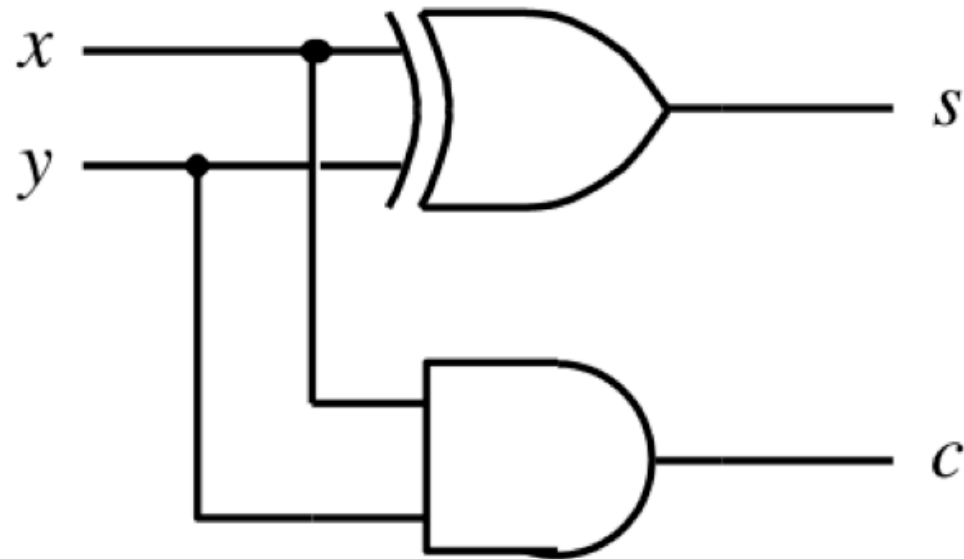
Adding two bits (there are four possible cases)



Adding two bits (the truth table)

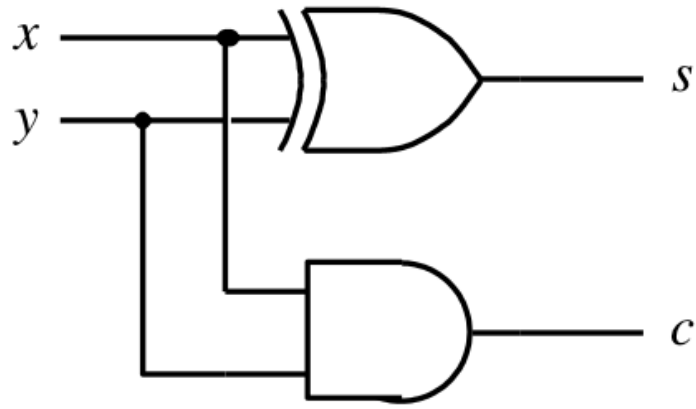
x	y	Carry c	Sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

Adding two bits (the logic circuit)

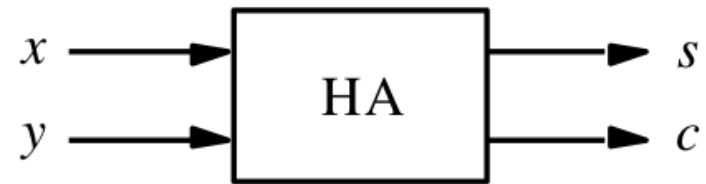


[Figure 3.1c from the textbook]

The Half-Adder



(c) Circuit

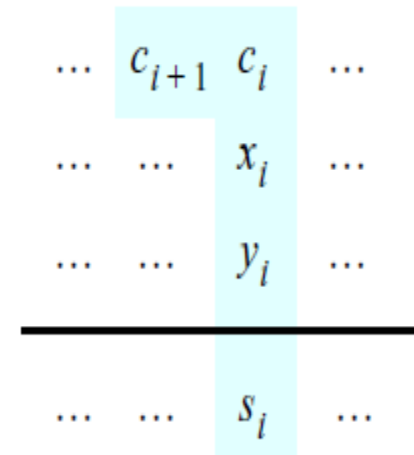


(d) Graphical symbol

Addition of multibit numbers

Generated carries \longrightarrow 1110

$$\begin{array}{r}
 X = x_4x_3x_2x_1x_0 \quad 01111 \quad (15)_{10} \\
 + Y = y_4y_3y_2y_1y_0 \quad + 01010 \quad + (10)_{10} \\
 \hline
 S = s_4s_3s_2s_1s_0 \quad 11001 \quad (25)_{10}
 \end{array}$$



Bit position i

Analogy with addition in base 10

$$\begin{array}{r} + \quad \quad \quad x_2 \quad x_1 \quad x_0 \\ \quad \quad \quad y_2 \quad y_1 \quad y_0 \\ \hline \quad \quad \quad s_2 \quad s_1 \quad s_0 \end{array}$$

Analogy with addition in base 10

$$\begin{array}{r} + \quad 3 \quad 8 \quad 9 \\ \quad 1 \quad 5 \quad 7 \\ \hline \quad 5 \quad 4 \quad 6 \end{array}$$

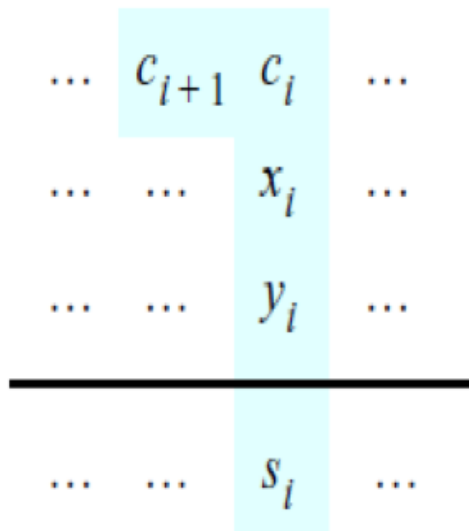
Analogy with addition in base 10

carry	0	1	1	0
		3	8	9
+		1	5	7
		<hr/>		
		5	4	6

Analogy with addition in base 10

$$\begin{array}{rcccc} & C_3 & C_2 & C_1 & C_0 \\ + & & X_2 & X_1 & X_0 \\ & & Y_2 & Y_1 & Y_0 \\ \hline & & S_2 & S_1 & S_0 \end{array}$$

Problem Statement and Truth Table



c_i	x_i	y_i	c_{i+1}	s_i
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

[Figure 3.2b from the textbook]

[Figure 3.3a from the textbook]

Let's fill-in the two K-maps

c_i	x_i	y_i	c_{i+1}	s_i
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

		$x_i y_i$			
		00	01	11	10
c_i	0				
	1				

$s_i =$

		$x_i y_i$			
		00	01	11	10
c_i	0				
	1				

$c_{i+1} =$

[Figure 3.3a-b from the textbook]

Let's fill-in the two K-maps

c_i	x_i	y_i	c_{i+1}	s_i
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$c_i \backslash x_i y_i$	00	01	11	10
0		1		1
1	1		1	

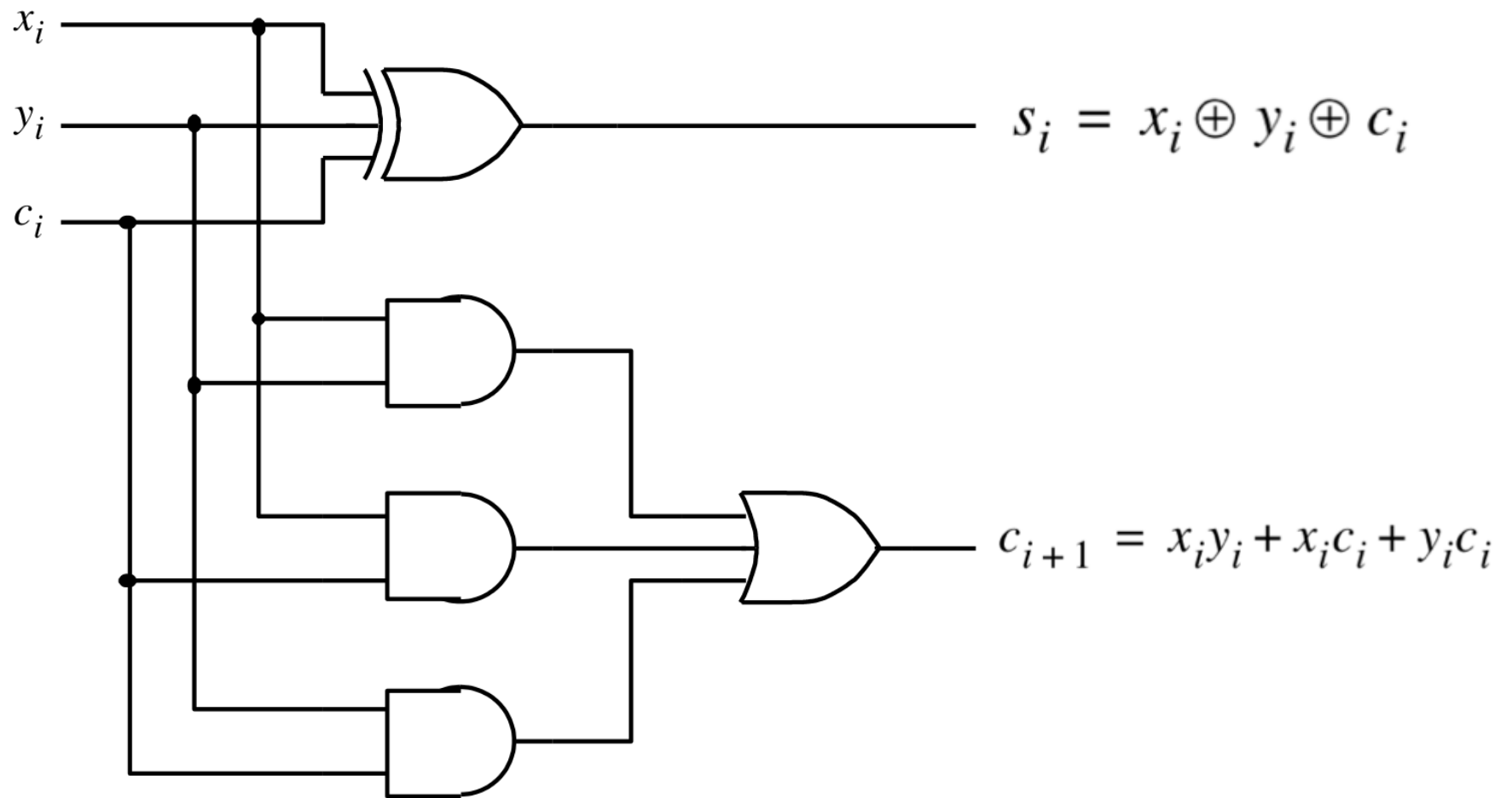
$$s_i = x_i \oplus y_i \oplus c_i$$

$c_i \backslash x_i y_i$	00	01	11	10
0			1	
1		1	1	1

$$c_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

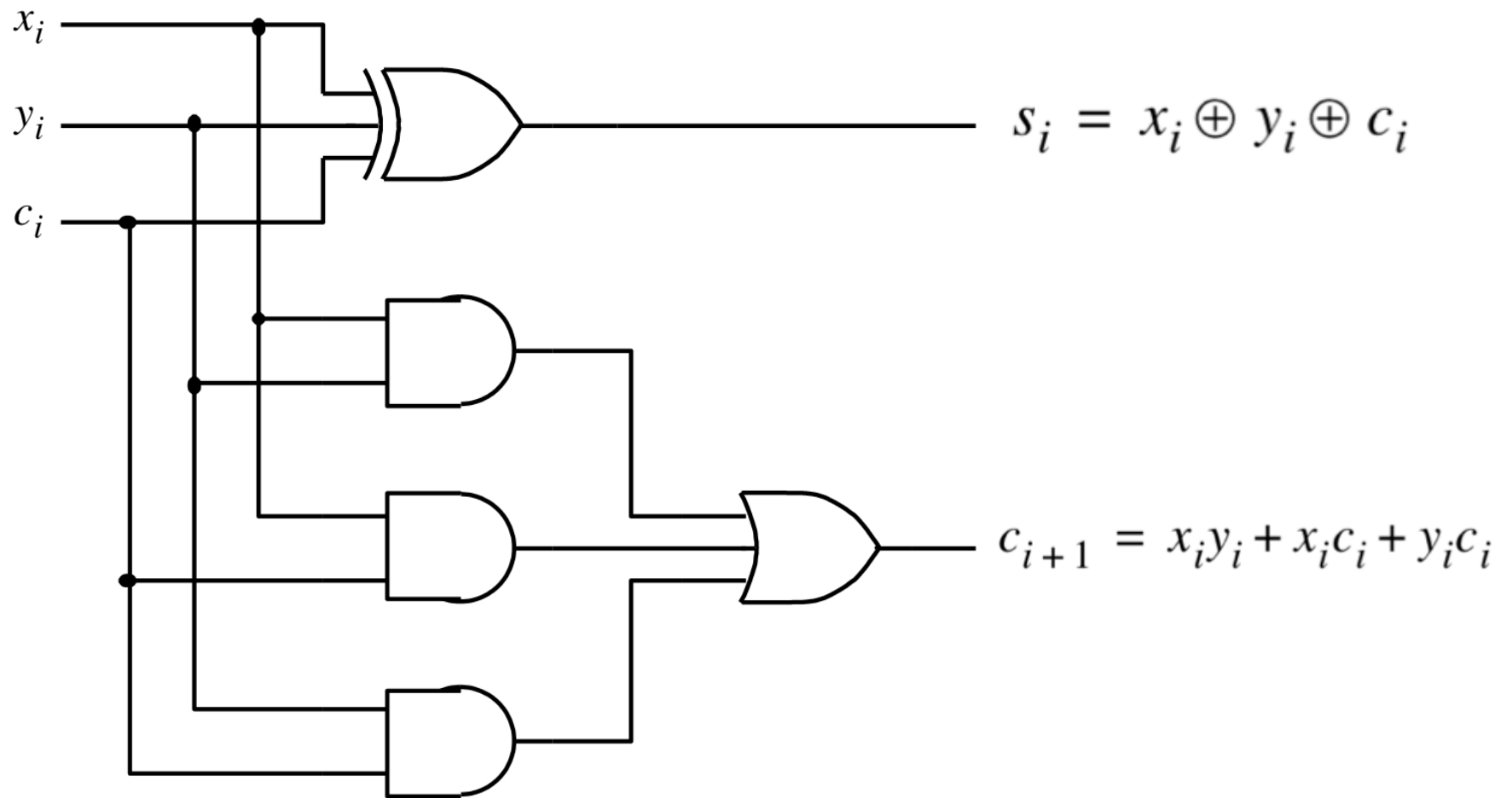
[Figure 3.3a-b from the textbook]

The circuit for the two expressions



[Figure 3.3c from the textbook]

This is called the Full-Adder

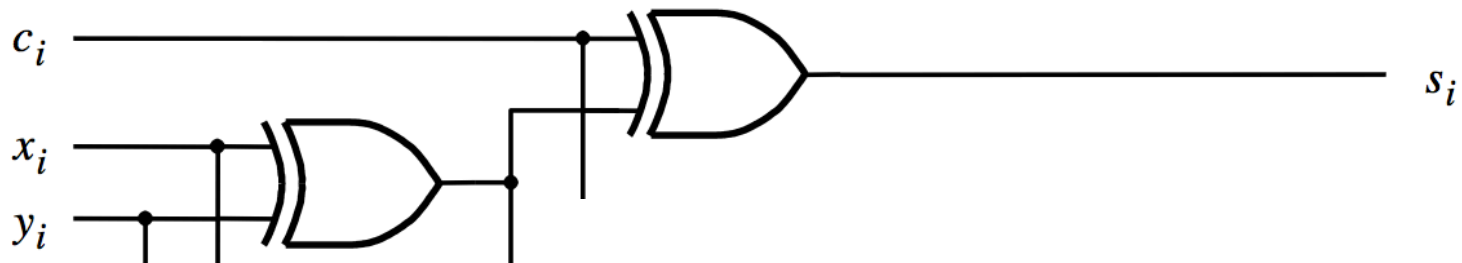
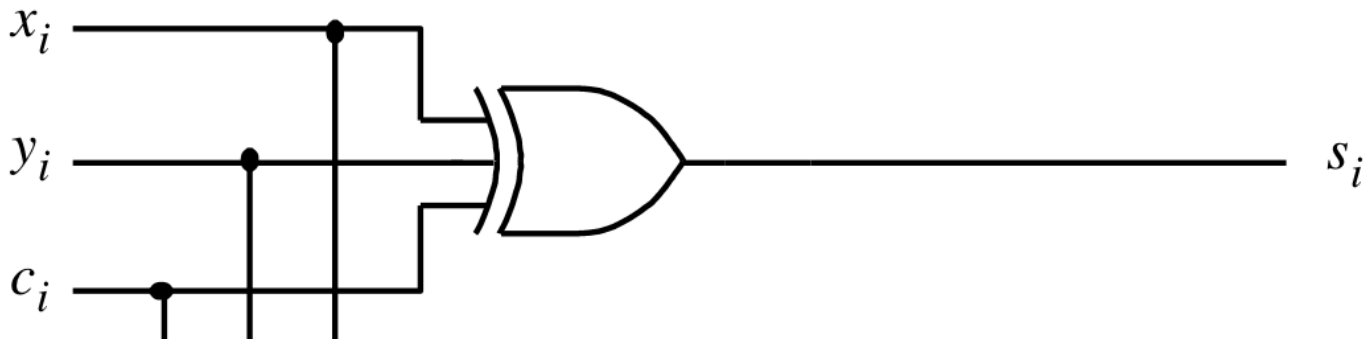


[Figure 3.3c from the textbook]

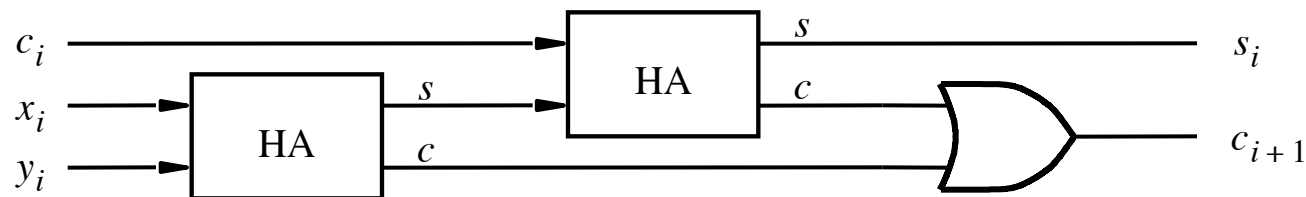
XOR Magic

(s_i can be implemented in two different ways)

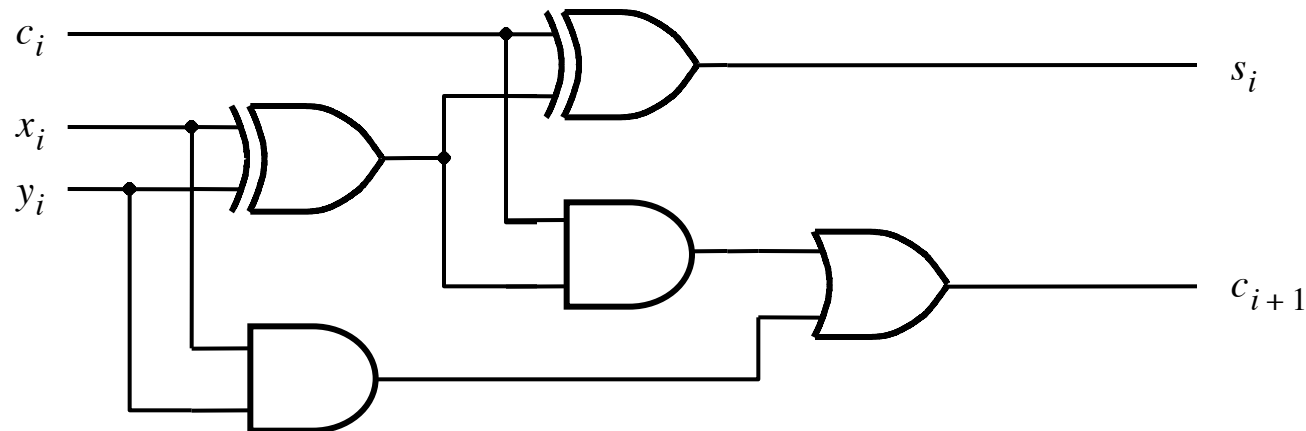
$$s_i = x_i \oplus y_i \oplus c_i$$



A decomposed implementation of the full-adder circuit

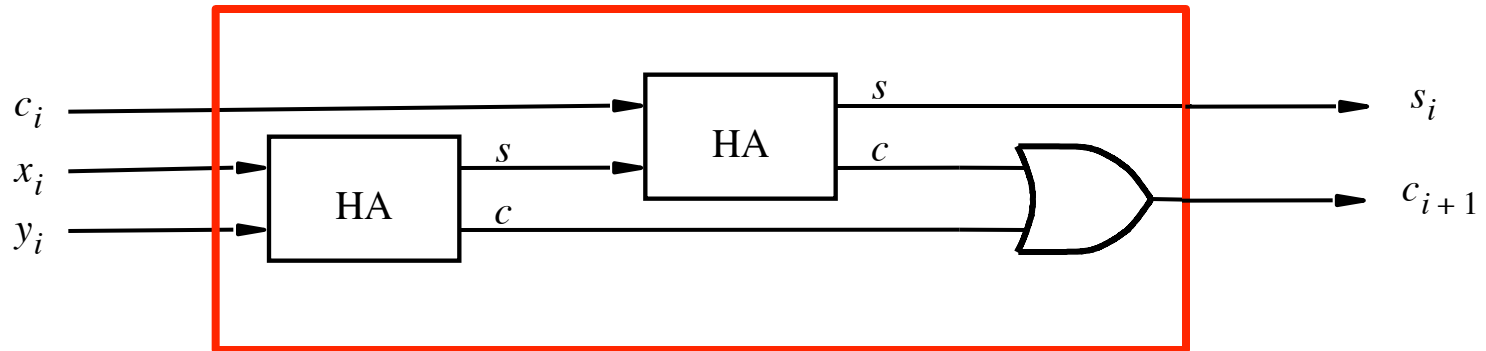


(a) Block diagram

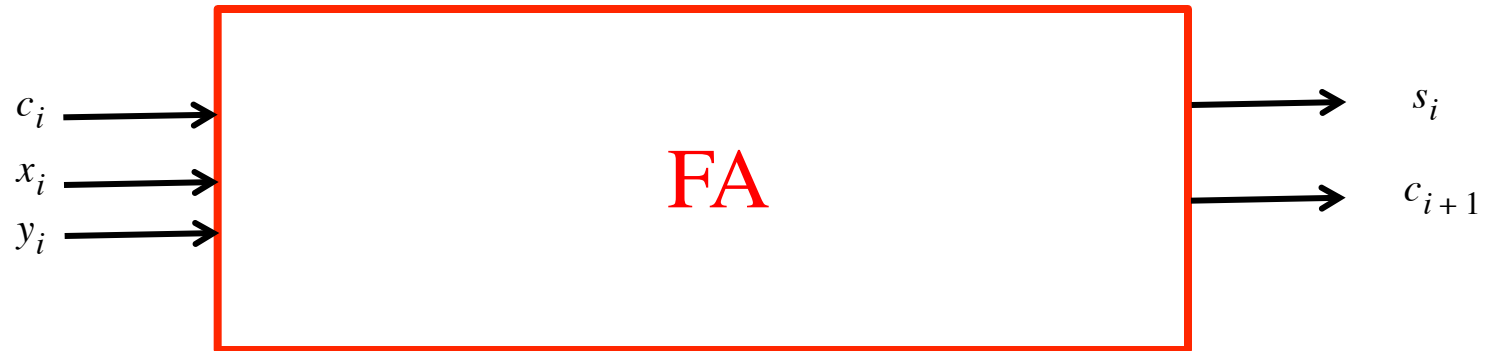


(b) Detailed diagram

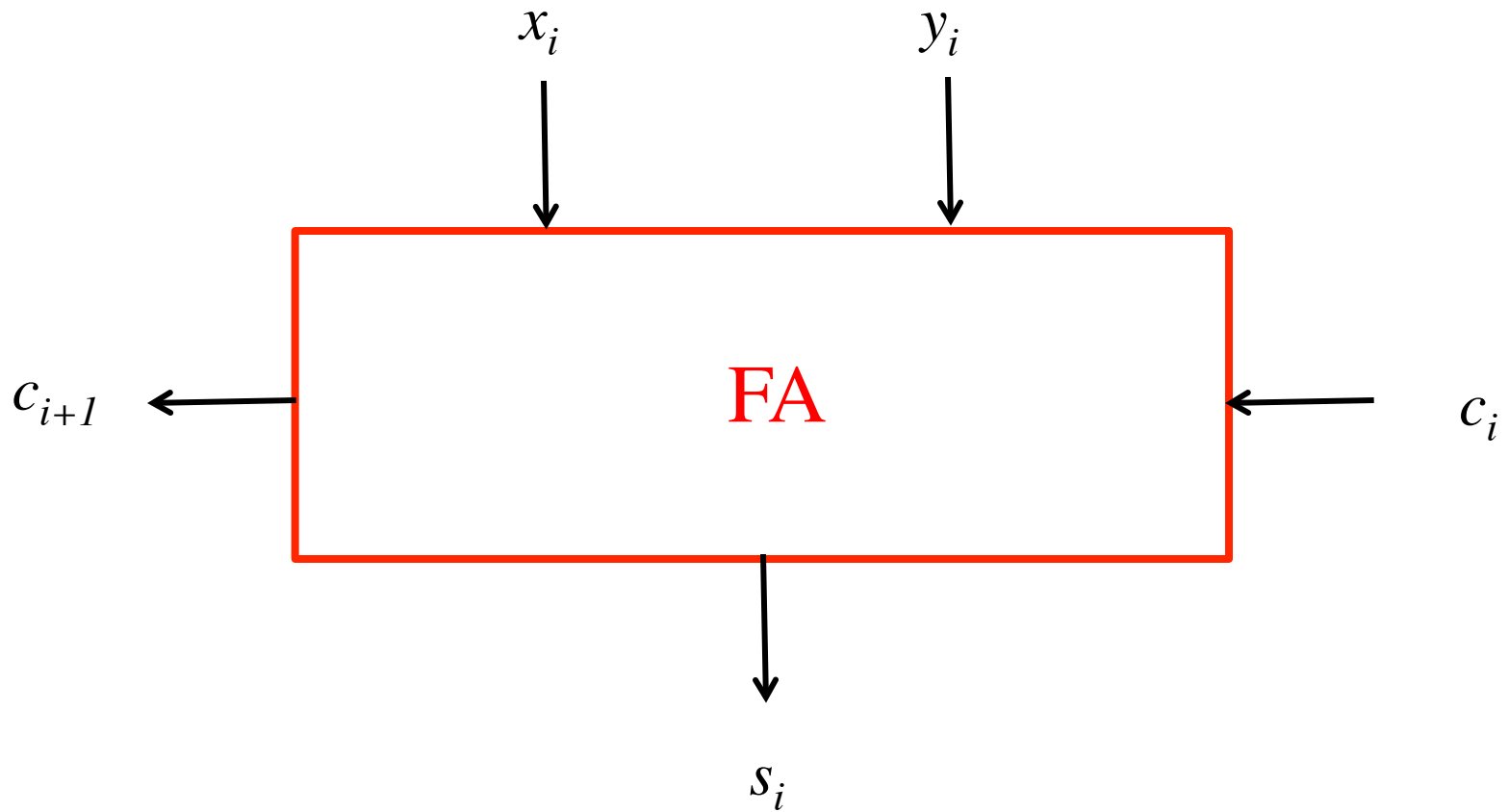
The Full-Adder Abstraction



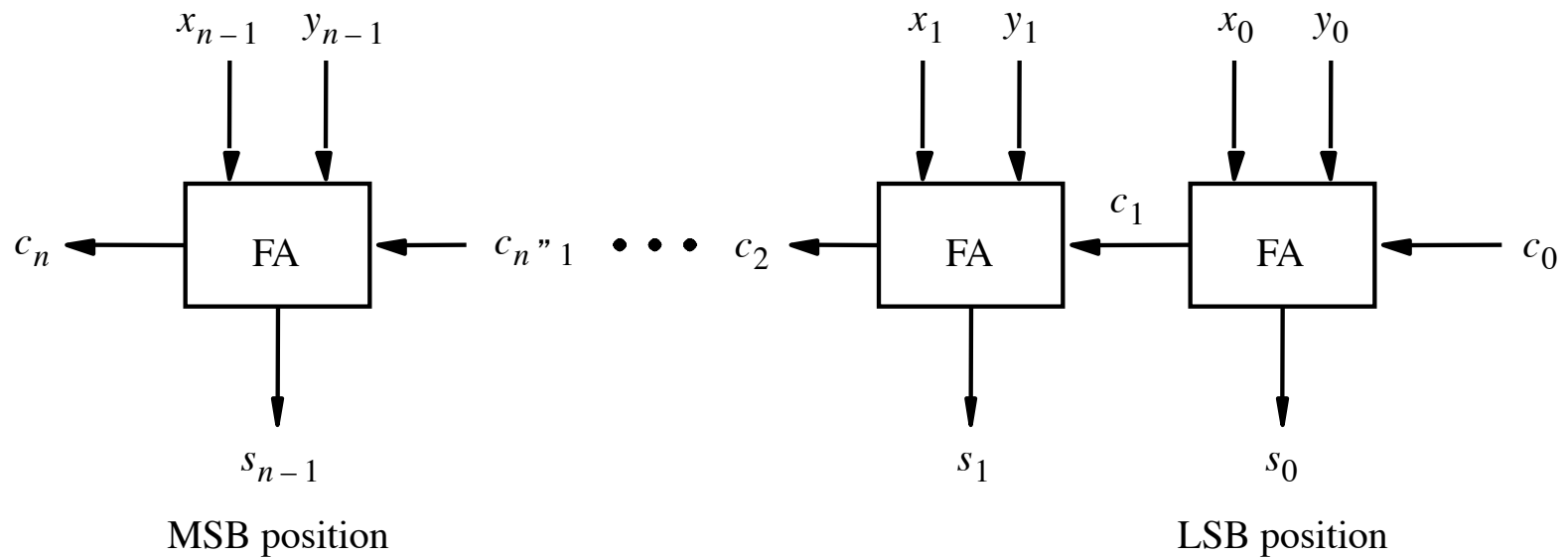
The Full-Adder Abstraction



We can place the arrows anywhere

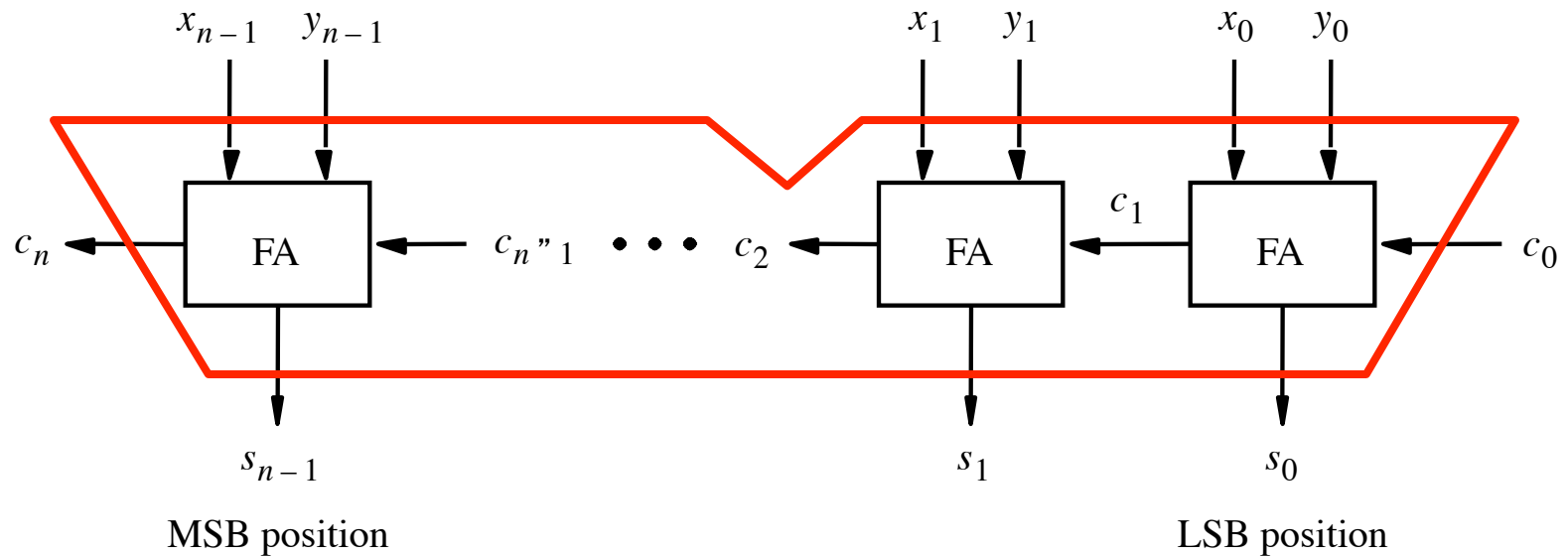


n-bit ripple-carry adder

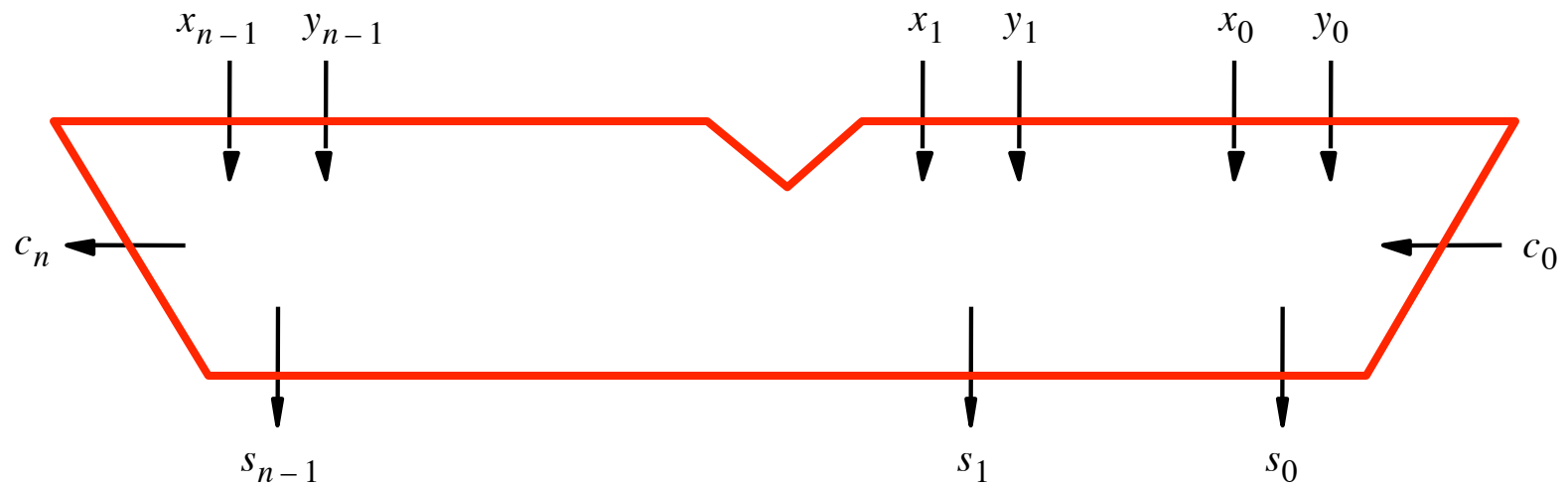


[Figure 3.5 from the textbook]

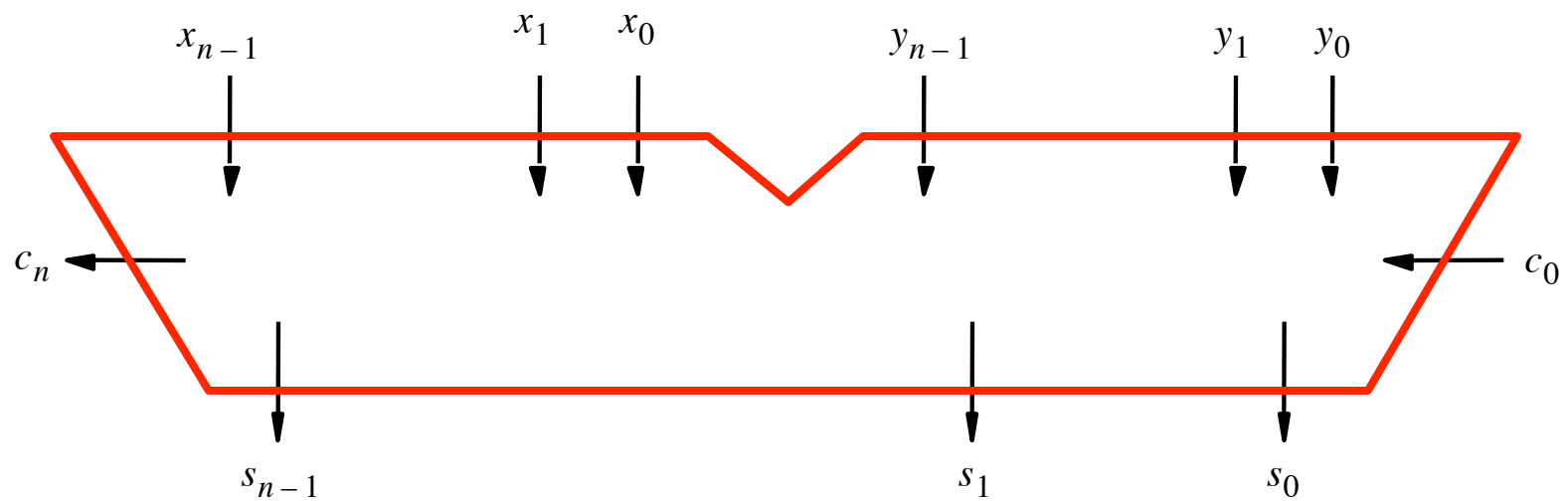
n-bit ripple-carry adder abstraction



n-bit ripple-carry adder abstraction



The x and y lines are typically grouped together for better visualization, but the underlying logic remains the same



Math Review: Subtraction

$$\begin{array}{r} 39 \\ - 15 \\ \hline ?? \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 39 \\ - 15 \\ \hline 24 \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ - 61 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 48 \\ - 26 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 32 \\ - 11 \\ \hline ?? \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ - 61 \\ \hline 21 \end{array}$$

$$\begin{array}{r} 48 \\ - 26 \\ \hline 22 \end{array}$$

$$\begin{array}{r} 32 \\ - 11 \\ \hline 21 \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ - 64 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 48 \\ - 29 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 32 \\ - 13 \\ \hline ?? \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ - 64 \\ \hline 18 \end{array}$$

$$\begin{array}{r} 48 \\ - 29 \\ \hline 19 \end{array}$$

$$\begin{array}{r} 32 \\ - 13 \\ \hline 19 \end{array}$$

The problems in which row are easier to calculate?

$$\begin{array}{r} 82 \\ - 61 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 48 \\ - 26 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 32 \\ - 11 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 82 \\ - 64 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 48 \\ - 29 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 32 \\ - 13 \\ \hline ?? \end{array}$$

The problems in which row are easier to calculate?

$$\begin{array}{r} 82 \\ - 61 \\ \hline 21 \end{array}$$

$$\begin{array}{r} 48 \\ - 26 \\ \hline 22 \end{array}$$

$$\begin{array}{r} 32 \\ - 11 \\ \hline 21 \end{array}$$

Why?

$$\begin{array}{r} 82 \\ - 64 \\ \hline 18 \end{array}$$

$$\begin{array}{r} 48 \\ - 29 \\ \hline 19 \end{array}$$

$$\begin{array}{r} 32 \\ - 13 \\ \hline 19 \end{array}$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

Another Way to Do Subtraction

$$\begin{aligned} 82 - 64 &= 82 + 100 - 100 - 64 \\ &= 82 + (100 - 64) - 100 \end{aligned}$$

Another Way to Do Subtraction

$$\begin{aligned}82 - 64 &= 82 + 100 - 100 - 64 \\ &= 82 + (100 - 64) - 100 \\ &= 82 + (99 + 1 - 64) - 100\end{aligned}$$

Another Way to Do Subtraction

$$\begin{aligned}82 - 64 &= 82 + 100 - 100 - 64 \\ &= 82 + (100 - 64) - 100 \\ &= 82 + (99 + 1 - 64) - 100 \\ &= 82 + (99 - 64) + 1 - 100\end{aligned}$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

$$= 82 + (100 - 64) - 100$$

$$= 82 + (99 + 1 - 64) - 100$$

Does not require borrows

$$= 82 + (99 - 64) + 1 - 100$$

9's Complement (subtract each digit from 9)

$$\begin{array}{r} 99 \\ - 64 \\ \hline 35 \end{array}$$

10's Complement

(subtract each digit from 9 and add 1 to the result)

$$\begin{array}{r} 99 \\ - 64 \\ \hline 35 + 1 = 36 \end{array}$$

Another Way to Do Subtraction

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$\begin{aligned}82 - 64 &= 82 + (99 - 64) + 1 - 100 \\ &= 82 + 35 + 1 - 100\end{aligned}$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

$$= 82 + (35 + 1) - 100$$

10's complement

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

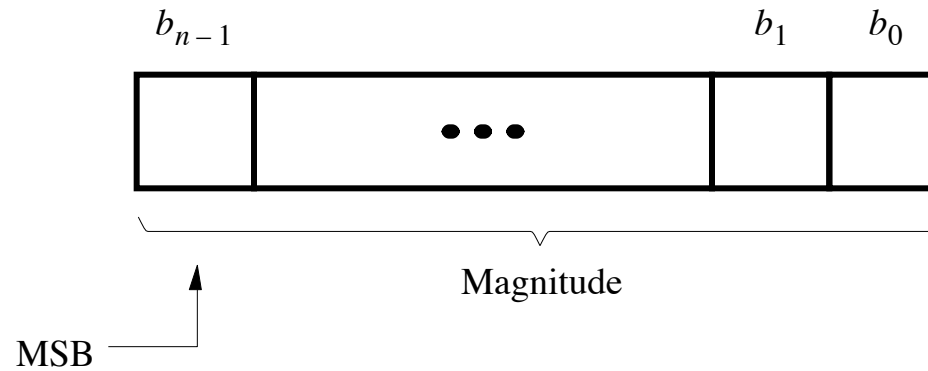
$$= 82 + (35 + 1) - 100$$

$$= 82 + 36 - 100 \quad // \text{ add the first two}$$

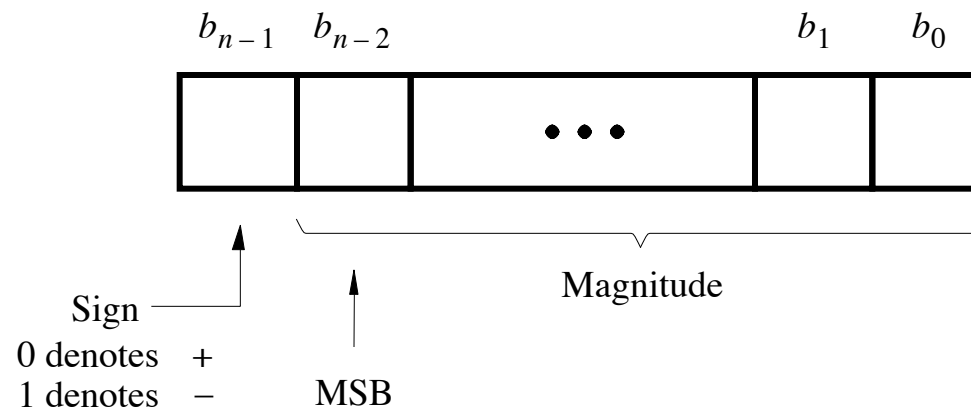
Another Way to Do Subtraction

$$\begin{aligned} 82 - 64 &= 82 + \overset{\text{9's complement}}{(99 - 64)} + 1 - 100 \\ &= 82 + \overset{\text{10's complement}}{(35 + 1)} - 100 \\ &= 82 + 36 - 100 \quad // \text{ add the first two} \\ &= 118 - 100 \quad // \text{ delete the leading 1} \\ &= 18 \end{aligned}$$

Formats for representation of integers



(a) Unsigned number



(b) Signed number

Negative numbers can be represented in following ways

- Sign and magnitude
- 1' s complement
- 2' s complement

1' s complement

Let K be the negative equivalent of an n -bit positive number P .

Then, in 1' s complement representation K is obtained by subtracting P from $2^n - 1$, namely

$$K = (2^n - 1) - P$$

This means that K can be obtained by inverting all bits of P .

Find the 1' s complement of ...

0 1 0 1

0 0 1 0

0 0 1 1

0 1 1 1

Find the 1's complement of ...

0 1 0 1

1 0 1 0

0 0 1 0

1 1 0 1

0 0 1 1

1 1 0 0

0 1 1 1

1 0 0 0

Just flip 1's to 0's and vice versa.

A) Example of 1's complement addition

$$\begin{array}{r}
 (+5) \\
 +(+2) \\
 \hline
 (+7)
 \end{array}
 \qquad
 \begin{array}{r}
 0101 \\
 +0010 \\
 \hline
 0111
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

[Figure 3.8 from the textbook]

A) Example of 1's complement addition

$$\begin{array}{r}
 (+5) \\
 +(+2) \\
 \hline
 (+7)
 \end{array}
 \quad
 \begin{array}{r}
 \color{red}{0101} \\
 + \color{green}{0010} \\
 \hline
 \color{blue}{0111}
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

B) Example of 1's complement addition

$$\begin{array}{r}
 (-5) \\
 +(+2) \\
 \hline
 (-3)
 \end{array}
 \qquad
 \begin{array}{r}
 1010 \\
 +0010 \\
 \hline
 1100
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

B) Example of 1's complement addition

$$\begin{array}{r}
 (-5) \\
 +(+2) \\
 \hline
 (-3)
 \end{array}
 \quad
 \begin{array}{r}
 \color{red}{1010} \\
 + \color{green}{0010} \\
 \hline
 \color{blue}{1100}
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

C) Example of 1's complement addition

$$\begin{array}{r}
 (+5) \quad 0101 \\
 +(-2) \quad +1101 \\
 \hline
 (+3) \quad 10010
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

C) Example of 1's complement addition

$$\begin{array}{r}
 (+5) \\
 +(-2) \\
 \hline
 (+3)
 \end{array}
 \quad
 \begin{array}{r}
 \underline{0101} \\
 + \underline{1101} \\
 \hline
 1 \underline{0010}
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

C) Example of 1's complement addition

$$\begin{array}{r}
 (+5) \\
 +(-2) \\
 \hline
 (+3)
 \end{array}
 \quad
 \begin{array}{r}
 \quad 0101 \\
 + \quad 1101 \\
 \hline
 1\quad 0010
 \end{array}$$

But this is 2!

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

C) Example of 1's complement addition

$$\begin{array}{r}
 (+5) \quad 0101 \\
 +(-2) \quad +1101 \\
 \hline
 (+3) \quad 10010 \\
 \quad \underline{1} \\
 \quad 0011
 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

C) Example of 1's complement addition

$$\begin{array}{r}
 (+5) \quad 0101 \\
 +(-2) \quad +1101 \\
 \hline
 (+3) \quad 10010 \\
 \text{Carry } \leftarrow 1 \\
 \hline
 0011
 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

D) Example of 1's complement addition

$$\begin{array}{r}
 (-5) \\
 + (-2) \\
 \hline
 (-7)
 \end{array}
 \qquad
 \begin{array}{r}
 1010 \\
 + 1101 \\
 \hline
 10111
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

[Figure 3.8 from the textbook]

D) Example of 1's complement addition

$$\begin{array}{r}
 (-5) \\
 + (-2) \\
 \hline
 (-7)
 \end{array}
 \quad
 \begin{array}{r}
 + \begin{array}{c} \color{red}{1010} \\ \color{green}{1101} \\ \hline \color{blue}{10111} \end{array}
 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

D) Example of 1's complement addition

$$\begin{array}{r}
 (-5) \\
 + (-2) \\
 \hline
 (-7)
 \end{array}
 \quad
 \begin{array}{r}
 1010 \\
 + 1101 \\
 \hline
 10111
 \end{array}$$

But this is +7!

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

D) Example of 1's complement addition

$$\begin{array}{r}
 (-5) \\
 + (-2) \\
 \hline
 (-7)
 \end{array}
 \qquad
 \begin{array}{r}
 1010 \\
 + 1101 \\
 \hline
 10111 \\
 \text{⌞} \rightarrow 1 \\
 \hline
 1000
 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

D) Example of 1's complement addition

$$\begin{array}{r}
 (-5) \\
 + (-2) \\
 \hline
 (-7)
 \end{array}
 \qquad
 \begin{array}{r}
 1010 \\
 + 1101 \\
 \hline
 10111 \\
 \text{⌞} \quad \text{⌞} \quad \text{⌞} \quad \text{⌞} \quad \text{⌞} \\
 \text{⌞} \quad \text{⌞} \quad \text{⌞} \quad \text{⌞} \quad \text{⌞} \\
 \hline
 1000
 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-7
1001	-6
1010	-5
1011	-4
1100	-3
1101	-2
1110	-1
1111	-0

2' s complement

Let K be the negative equivalent of an n -bit positive number P .

Then, in 2' s complement representation K is obtained by subtracting P from 2^n , namely

$$K = 2^n - P$$

Deriving 2' s complement

For a positive n-bit number P, let K_1 and K_2 denote its 1' s and 2' s complements, respectively.

$$K_1 = (2^n - 1) - P$$
$$K_2 = 2^n - P$$

Since $K_2 = K_1 + 1$, it is evident that in a logic circuit the 2' s complement can be computed by inverting all bits of P and then adding 1 to the resulting 1' s-complement number.

Find the 2' s complement of ...

0 1 0 1

0 0 1 0

0 0 1 1

0 1 1 1

Find the 2' s complement of ...

0 1 0 1

1 0 1 1

0 0 1 0

1 1 1 0

0 0 1 1

1 1 0 1

0 1 1 1

1 0 0 1

Quick Way to find 2's complement

- **Scan the binary number from right to left**
- **Copy all bits that are 0 from right to left**
- **Stop at the first 1**
- **Copy that 1 as well**
- **Invert all remaining bits**

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	-0	-7	-8
1001	-1	-6	-7
1010	-2	-5	-6
1011	-3	-4	-5
1100	-4	-3	-4
1101	-5	-2	-3
1110	-6	-1	-2
1111	-7	-0	-1

[Table 3.2 from the textbook]

A) Example of 2's complement addition

$$\begin{array}{r}
 (+5) \quad 0101 \\
 + (+2) \quad 0010 \\
 \hline
 (+7) \quad 0111
 \end{array}$$

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-8
1001	-7
1010	-6
1011	-5
1100	-4
1101	-3
1110	-2
1111	-1

[Figure 3.9 from the textbook]

B) Example of 2's complement addition


$$\begin{array}{r}
 (-5) \quad 1011 \\
 + (+2) \quad + 0010 \\
 \hline
 (-3) \quad 1101
 \end{array}$$

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-8
1001	-7
1010	-6
1011	-5
1100	-4
1101	-3
1110	-2
1111	-1

[Figure 3.9 from the textbook]

C) Example of 2's complement addition

$$\begin{array}{r}
 (+5) \quad 0101 \\
 + (-2) \quad +1110 \\
 \hline
 (+3) \quad 10011
 \end{array}$$




 ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-8
1001	-7
1010	-6
1011	-5
1100	-4
1101	-3
1110	-2
1111	-1

[Figure 3.9 from the textbook]

D) Example of 2's complement addition

$$\begin{array}{r}
 (-5) \quad \quad 1\ 0\ 1\ 1 \\
 +\ (-2) \quad +\ 1\ 1\ 1\ 0 \\
 \hline
 (-7) \quad \quad 1\ 1\ 0\ 0\ 1
 \end{array}$$


 ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	-8
1001	-7
1010	-6
1011	-5
1100	-4
1101	-3
1110	-2
1111	-1

Example of 2' s complement subtraction

$$\begin{array}{r} (+5) \quad 0101 \\ - (+2) \quad \underline{0010} \\ \hline (+3) \end{array} \quad \Rightarrow \quad \begin{array}{r} 0101 \\ + 1110 \\ \hline 10011 \\ \uparrow \\ \text{ignore} \end{array}$$

Example of 2's complement subtraction

$$\begin{array}{r} (-5) \\ - (+2) \\ \hline (-7) \end{array} \quad \begin{array}{r} 1011 \\ - 0010 \\ \hline \end{array} \quad \Rightarrow \quad \begin{array}{r} 1011 \\ + 1110 \\ \hline 11001 \\ \uparrow \\ \text{ignore} \end{array}$$

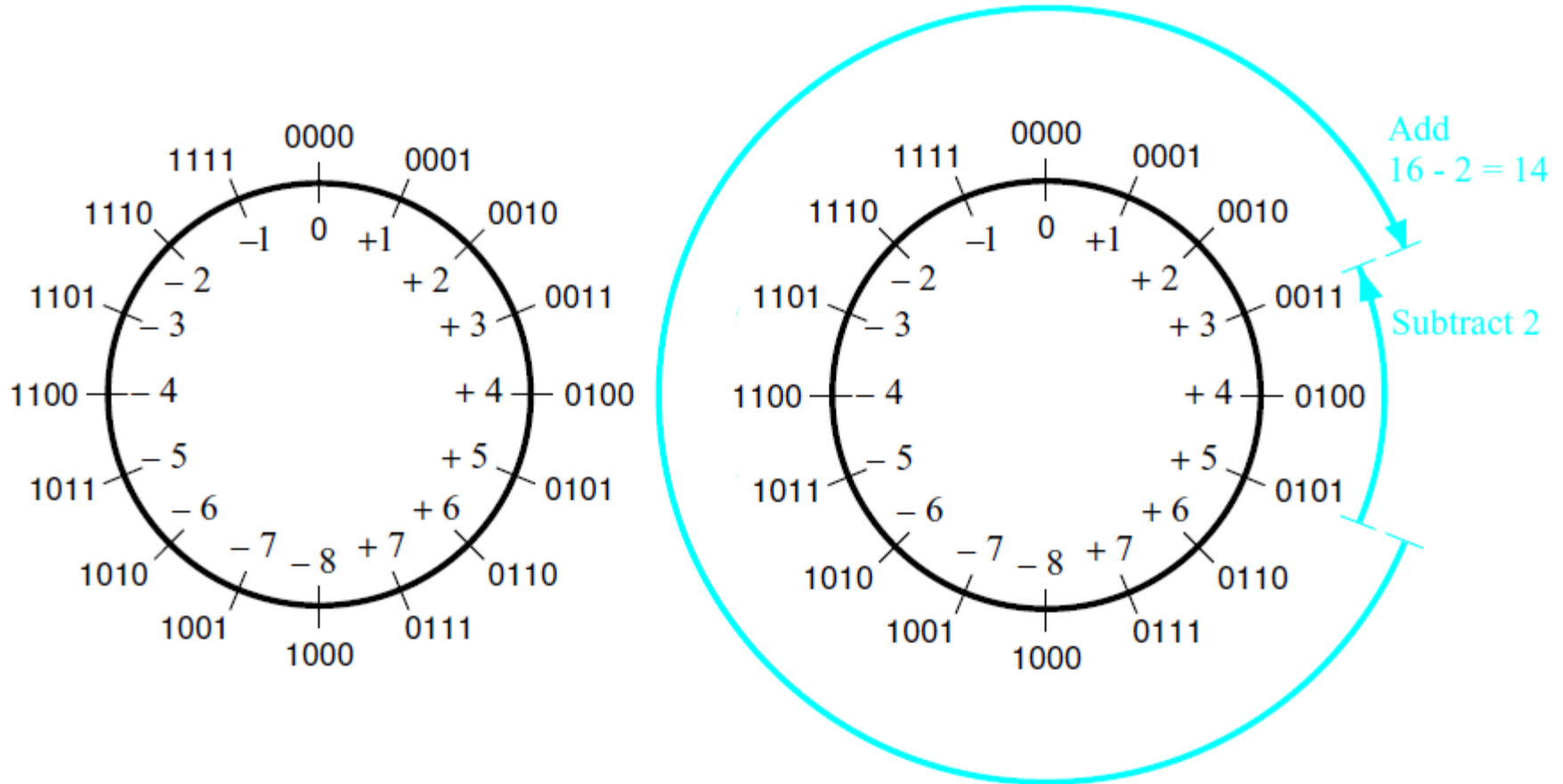
Example of 2' s complement subtraction

$$\begin{array}{r} (+5) \quad 0101 \\ - (-2) \quad \underline{1110} \\ \hline (+7) \end{array} \quad \Rightarrow \quad \begin{array}{r} 0101 \\ + \underline{0010} \\ \hline 0111 \end{array}$$

Example of 2's complement subtraction

$$\begin{array}{r} (-5) \\ - (-2) \\ \hline (-3) \end{array} \quad \begin{array}{r} 1011 \\ - 1110 \\ \hline \end{array} \quad \Rightarrow \quad \begin{array}{r} 1011 \\ + 0010 \\ \hline 1101 \end{array}$$

Graphical interpretation of four-bit 2's complement numbers



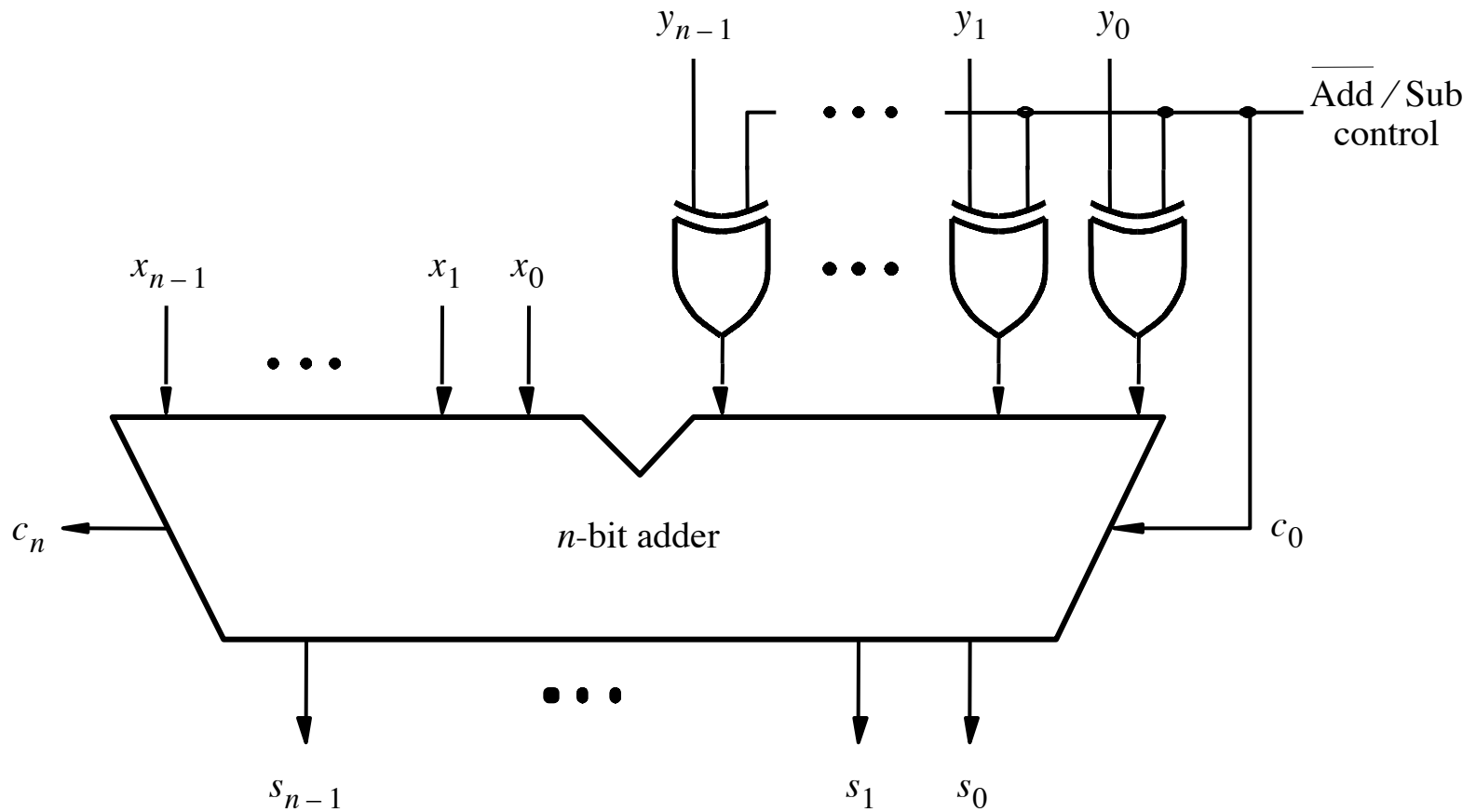
(a) The number circle

(b) Subtracting 2 by adding its 2's complement

Take-Home Message

- **Subtraction can be performed by simply adding the 2's complement of the second number, regardless of the signs of the two numbers.**
- **Thus, the same adder circuit can be used to perform both addition and subtraction !!!**

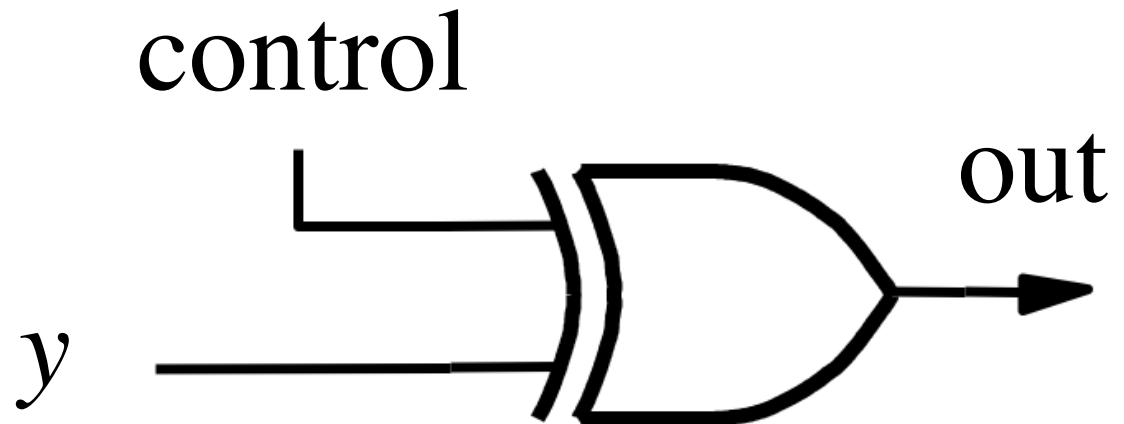
Adder/subtractor unit



[Figure 3.12 from the textbook]

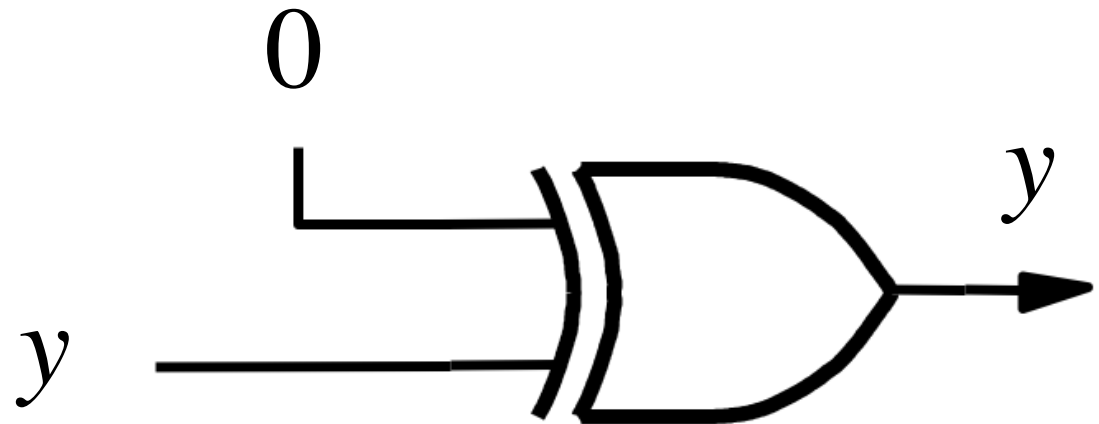
XOR Tricks

control	y	out
0	0	0
0	1	1
1	0	1
1	1	0



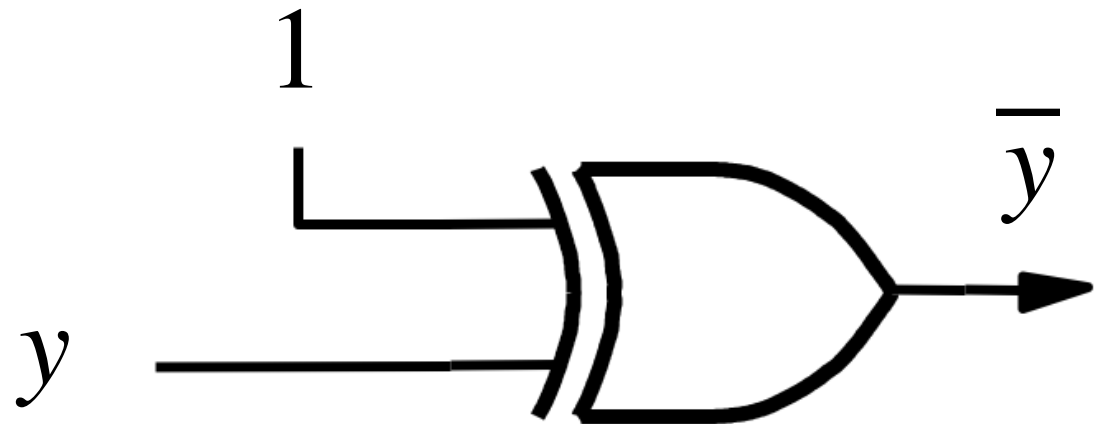
XOR as a repeater

control	y	out
0	0	0
0	1	1
1	0	1
1	1	0

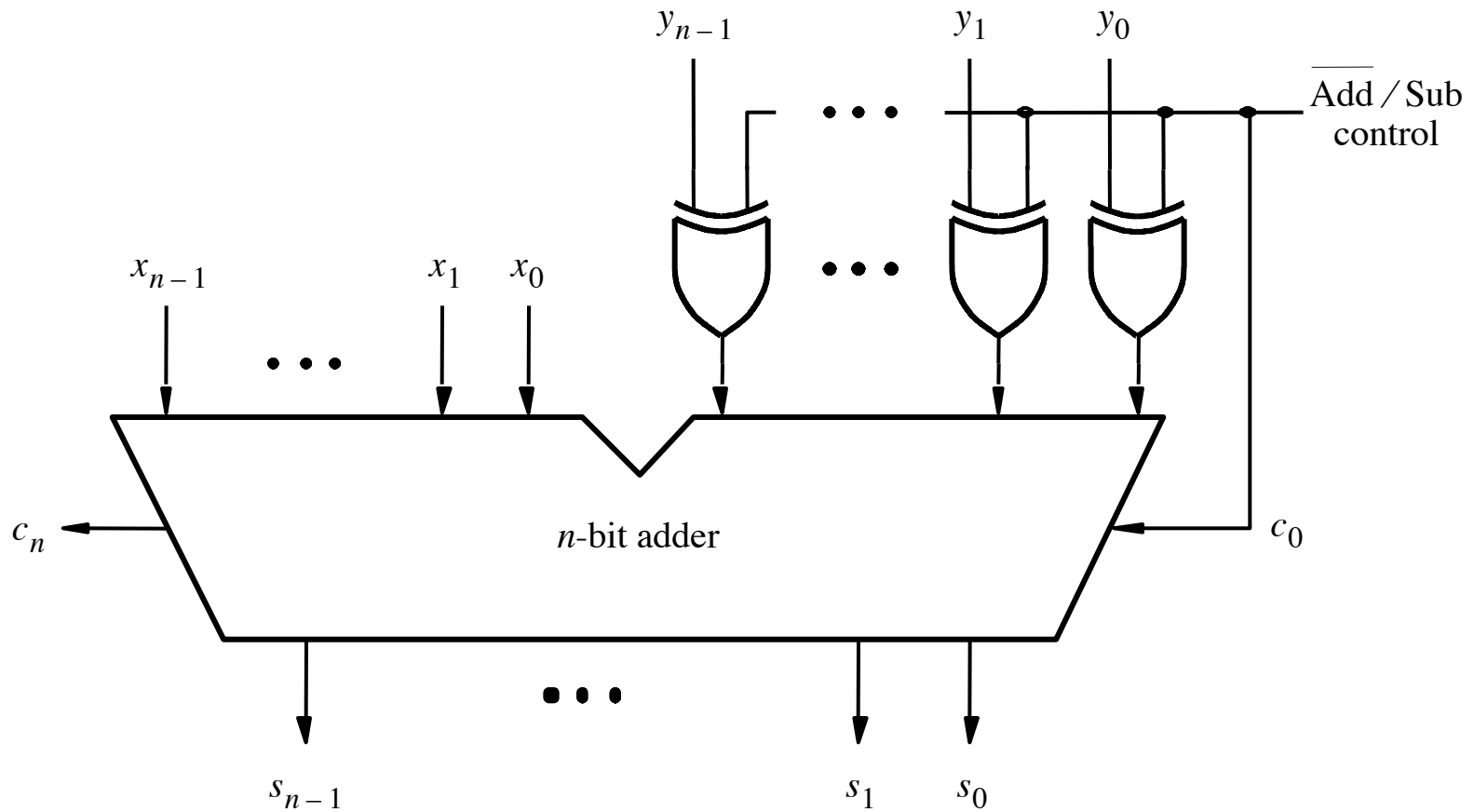


XOR as an inverter

control	y	out
0	0	0
0	1	1
1	0	1
1	1	0

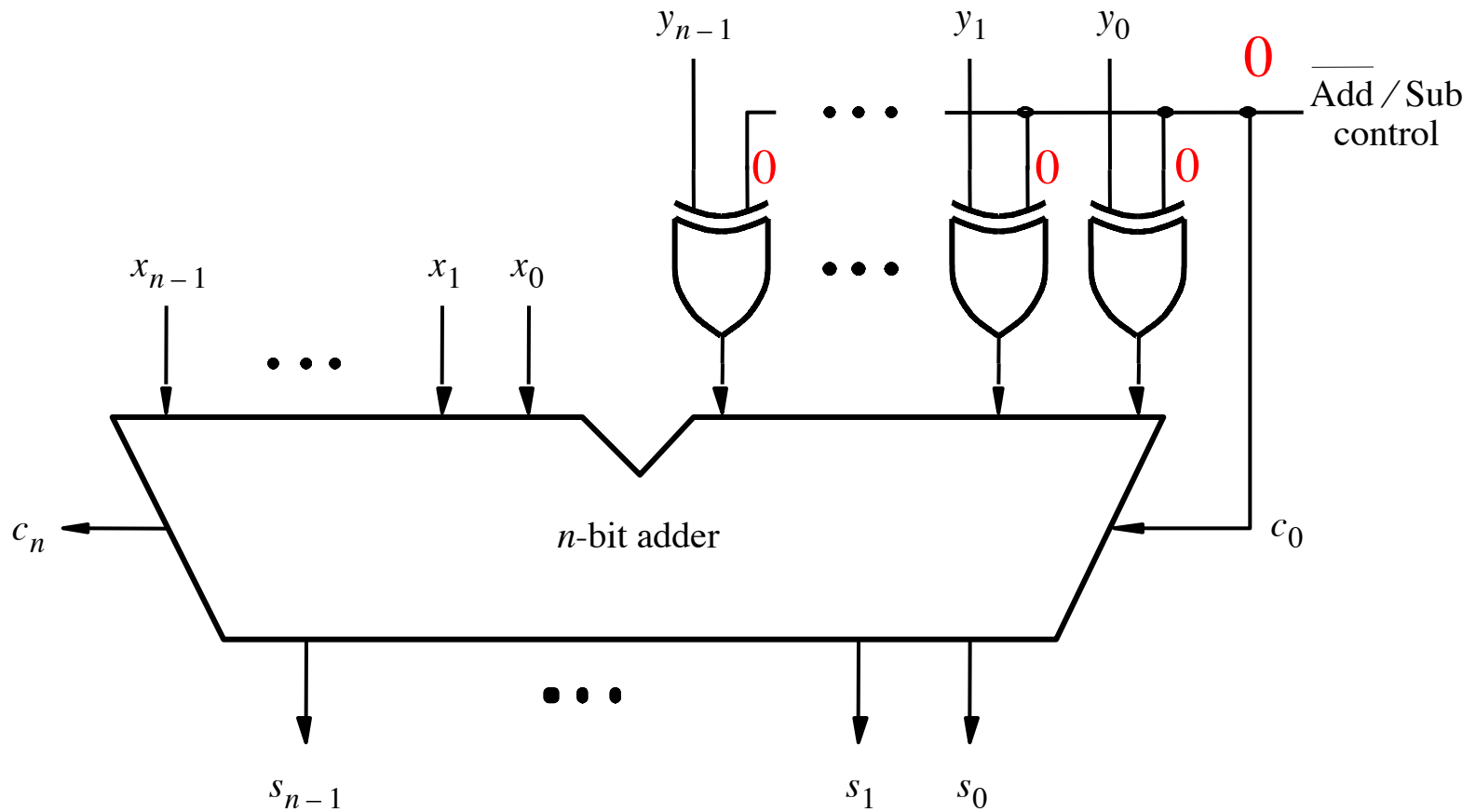


Addition: when control = 0



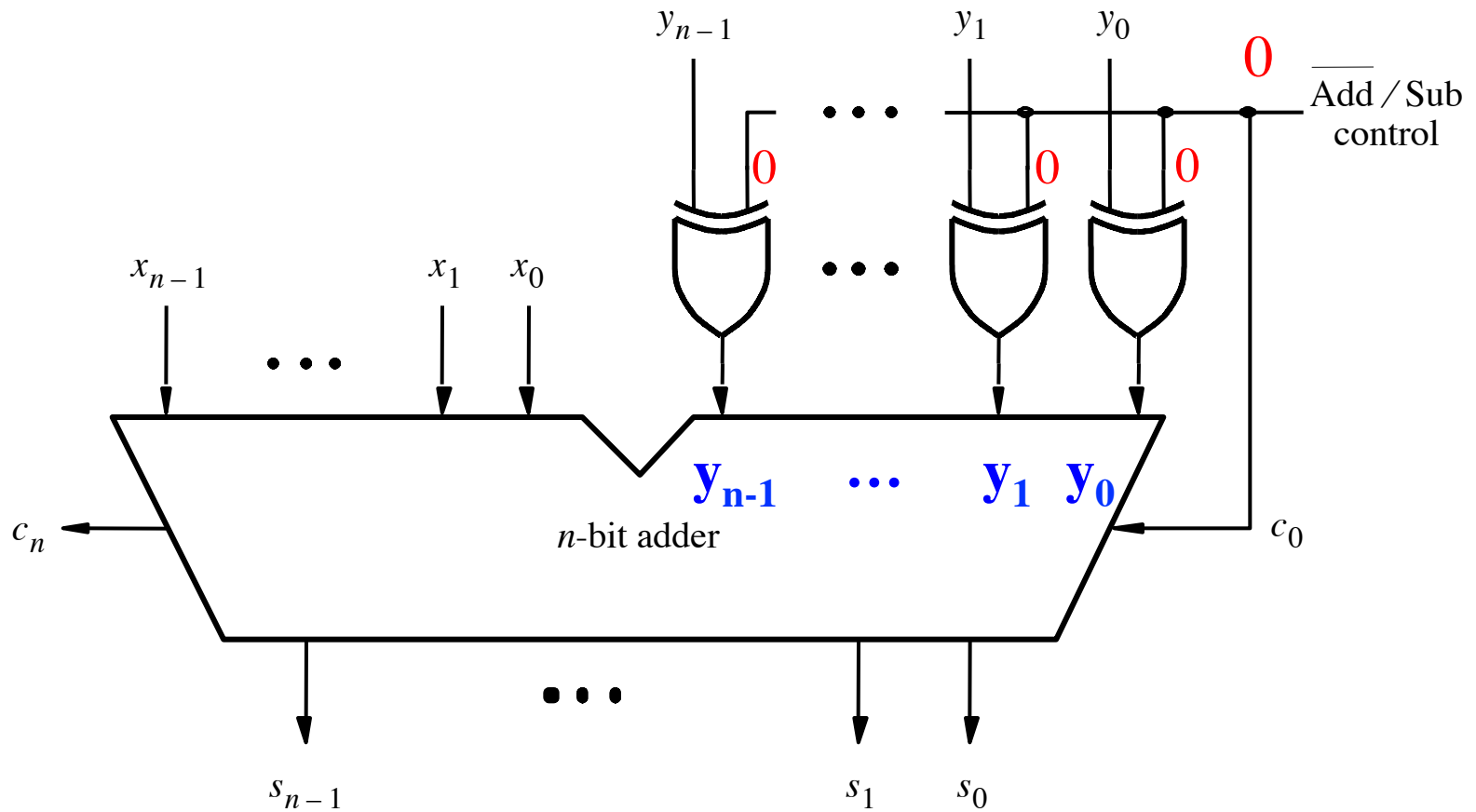
[Figure 3.12 from the textbook]

Addition: when control = 0



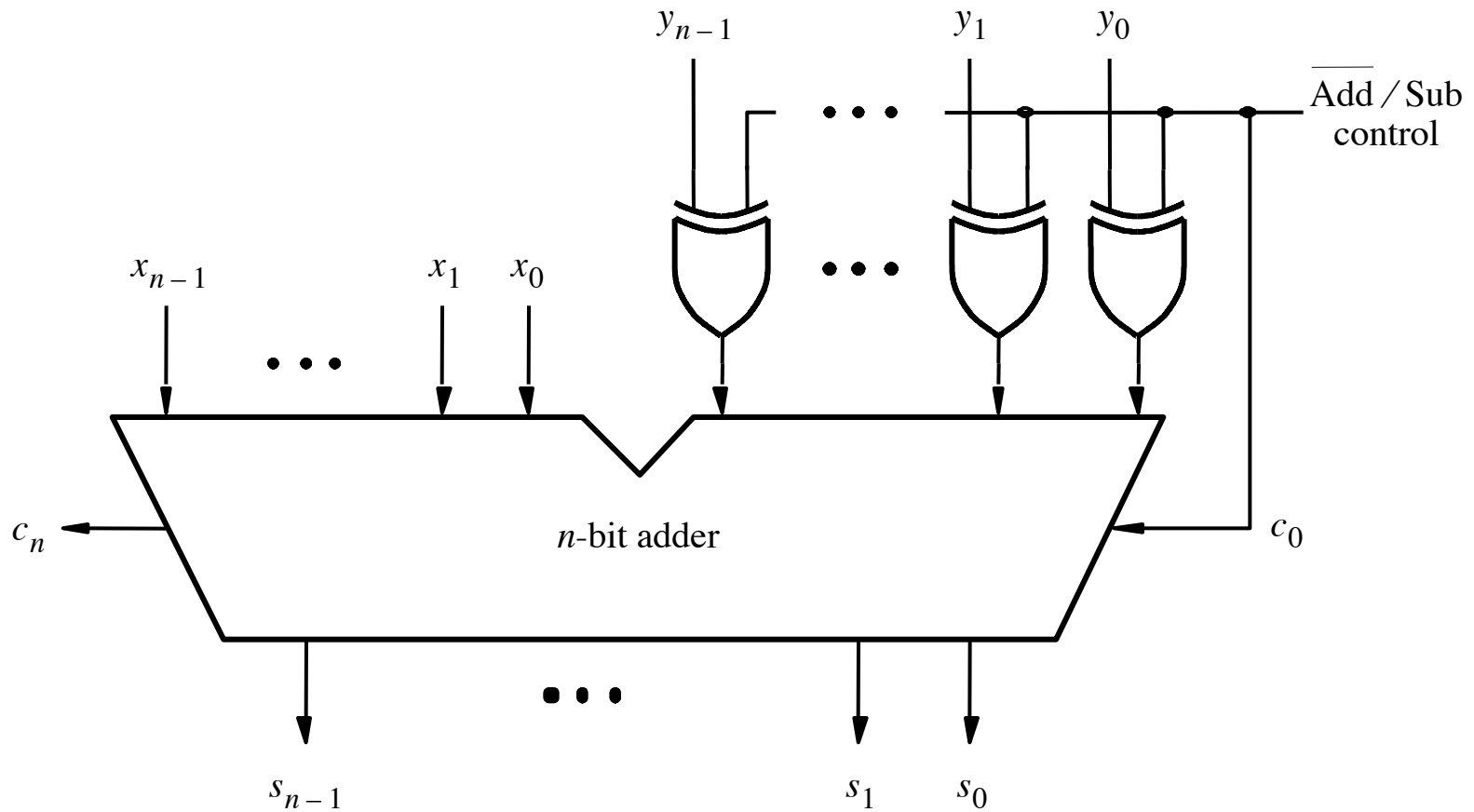
[Figure 3.12 from the textbook]

Addition: when control = 0



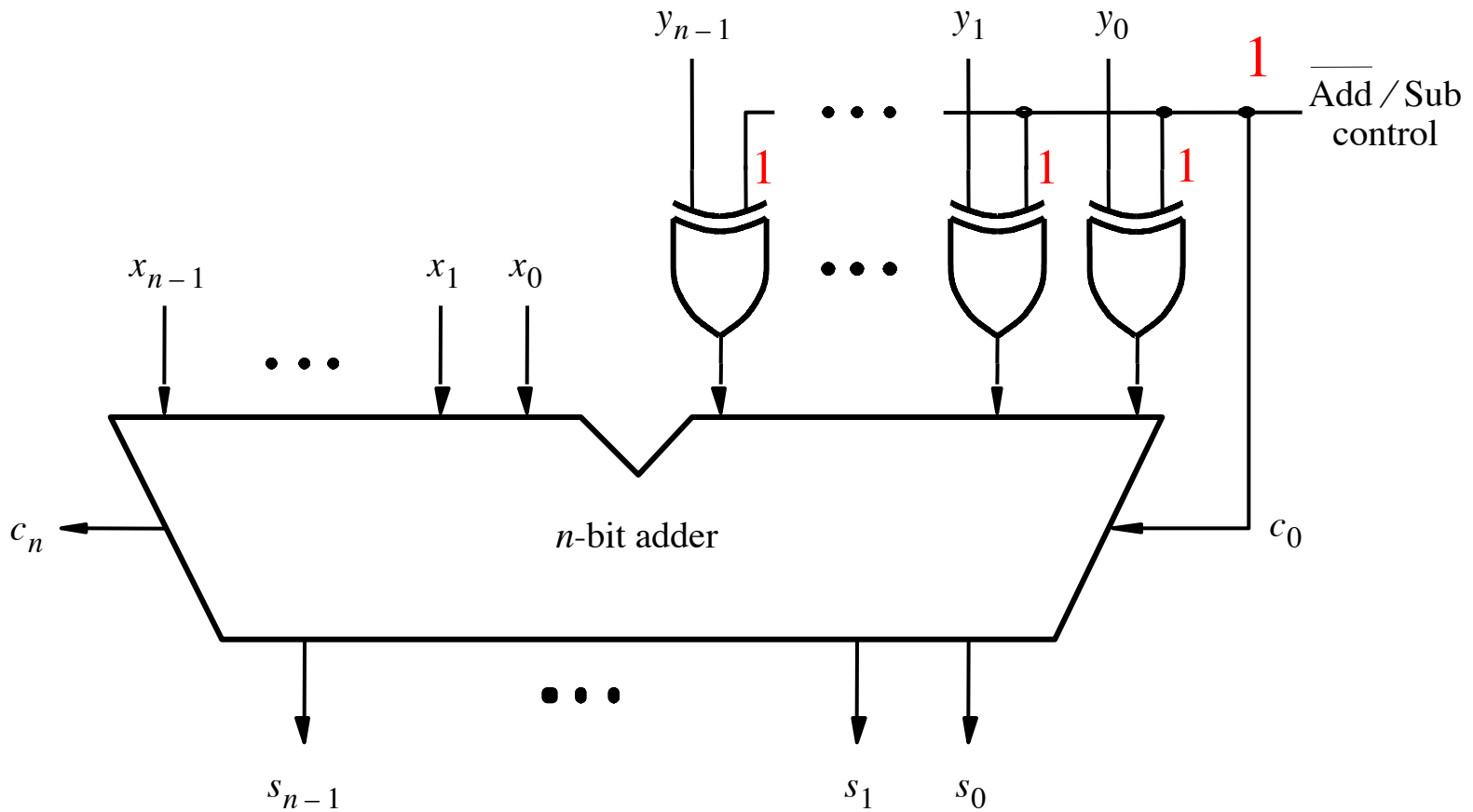
[Figure 3.12 from the textbook]

Subtraction: when control = 1



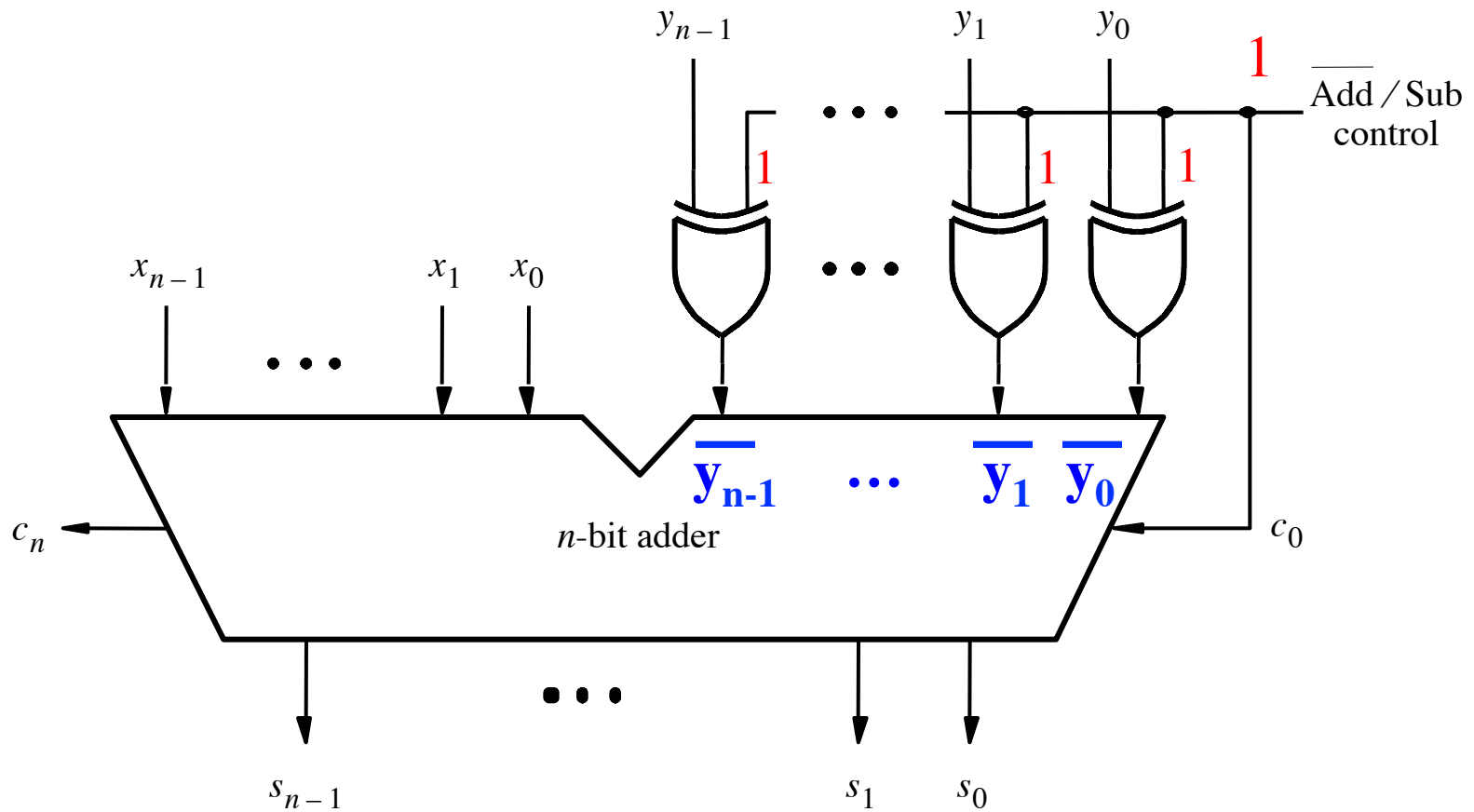
[Figure 3.12 from the textbook]

Subtraction: when control = 1



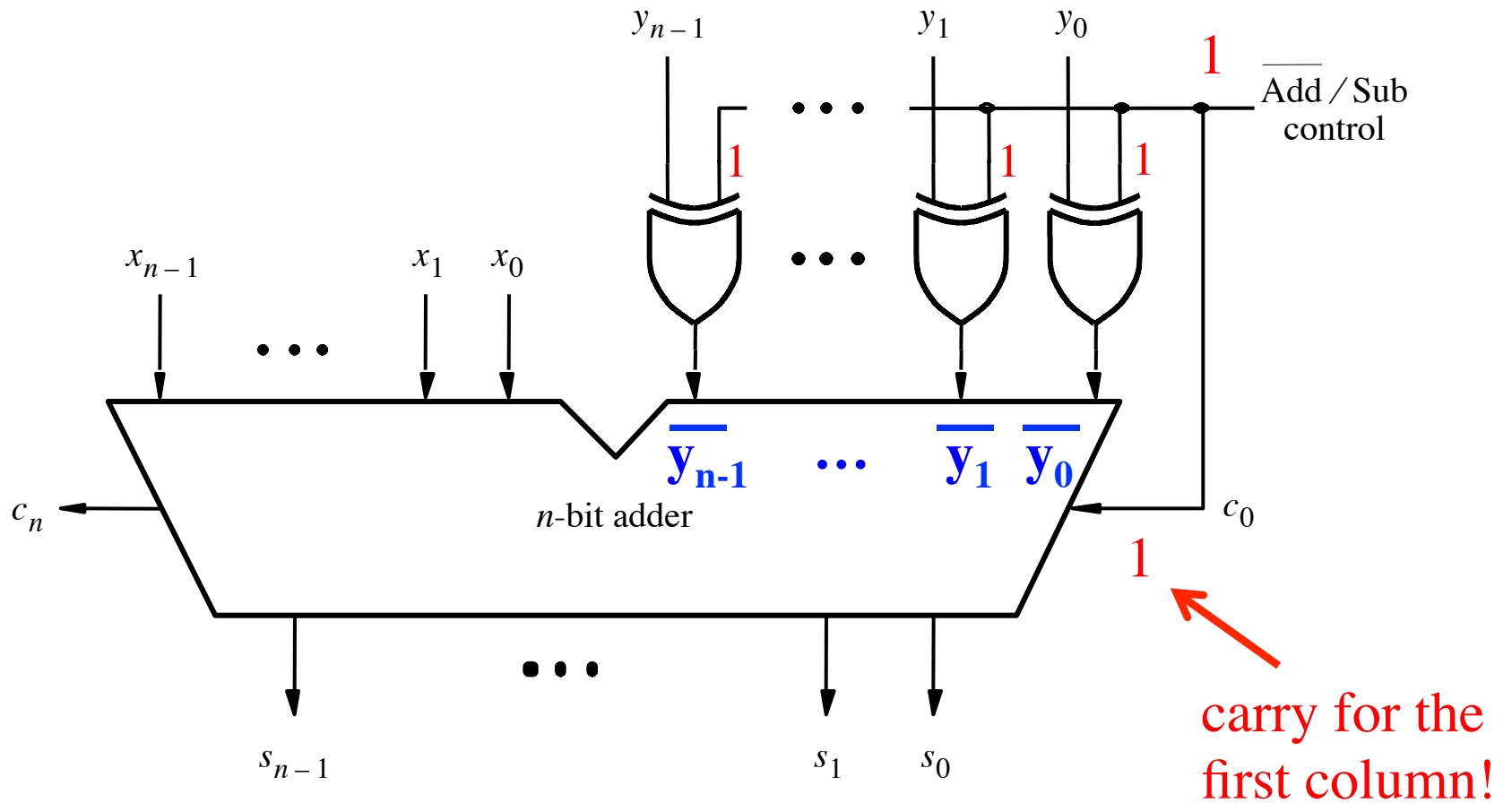
[Figure 3.12 from the textbook]

Subtraction: when control = 1



[Figure 3.12 from the textbook]

Subtraction: when control = 1



[Figure 3.12 from the textbook]

Examples of determination of overflow

$$\begin{array}{r}
 (+7) \quad 0111 \\
 + (+2) \quad +0010 \\
 \hline
 (+9) \quad 1001 \\
 c_4 = 0 \\
 c_3 = 1
 \end{array}$$

$$\begin{array}{r}
 (-7) \quad 1001 \\
 + (+2) \quad +0010 \\
 \hline
 (-5) \quad 1011 \\
 c_4 = 0 \\
 c_3 = 0
 \end{array}$$

$$\begin{array}{r}
 (+7) \quad 0111 \\
 + (-2) \quad +1110 \\
 \hline
 (+5) \quad 10101 \\
 c_4 = 1 \\
 c_3 = 1
 \end{array}$$

$$\begin{array}{r}
 (-7) \quad 1001 \\
 + (-2) \quad +1110 \\
 \hline
 (-9) \quad 10111 \\
 c_4 = 1 \\
 c_3 = 0
 \end{array}$$

Examples of determination of overflow

$$\begin{array}{r}
 (+7) \quad 0111 \\
 + (+2) \quad +0010 \\
 \hline
 (+9) \quad 1001 \\
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 (-7) \quad 1001 \\
 + (+2) \quad +0010 \\
 \hline
 (-5) \quad 1011 \\
 c_4 = 0 \\
 c_3 = 0
 \end{array}$$

Overflow occurs only in these two cases.

$$\begin{array}{r}
 (+7) \quad 0111 \\
 + (-2) \quad +1110 \\
 \hline
 (+5) \quad 10101 \\
 c_4 = 1 \\
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 \hline
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 (-9) \quad 10111 \\
 \\
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 c_3 = 0
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$$\text{Overflow} = c_3 \bar{c}_4 + \bar{c}_3 c_4$$

Examples of determination of overflow

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 (-7) \quad 1001 \\
 + (-2) \quad +1110 \\
 \hline
 (-9) \quad 10111 \\
 c_4 = 1 \\
 c_3 = 0
 \end{array}$$

$$\text{Overflow} = \underbrace{c_3 \bar{c}_4 + \bar{c}_3 c_4}_{\text{XOR}}$$

Calculating overflow for 4-bit numbers with only three significant bits

$$\begin{aligned}\text{Overflow} &= c_3 \bar{c}_4 + \bar{c}_3 c_4 \\ &= c_3 \oplus c_4\end{aligned}$$

Calculating overflow for n-bit numbers with only n-1 significant bits

$$\text{Overflow} = c_{n-1} \oplus c_n$$

Another way to look at the overflow issue

$$X = x_3x_2x_1x_0$$

$$Y = y_3y_2y_1y_0$$

$$S = s_3s_2s_1s_0$$

Another way to look at the overflow issue

$$X = x_3x_2x_1x_0$$

$$Y = y_3y_2y_1y_0$$

$$S = s_3s_2s_1s_0$$

If both numbers that we are adding have the same sign but the sum does not, then we have an overflow.

$$\text{Overflow} = x_3y_3\bar{s}_3 + \bar{x}_3\bar{y}_3s_3$$

Questions?

THE END