

CprE 281: Digital Logic

Instructor: Alexander Stoytchev

http://www.ece.iastate.edu/~alexs/classes/

Registers

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Administrative Stuff

- The second midterm is next Friday.
- Homework 8 is due next Monday.

Administrative Stuff

- Midterm Exam #2
- When: Friday October 30 @ 4pm.
- Where: This classroom
- What: Chapters 1, 2, 3, 4 and 5.1-5.7
- The exam will be open book and open notes (you can bring up to 3 pages of handwritten notes).

Midterm 2: Format

- The exam will be out of 130 points
- You need 95 points to get an A
- It will be great if you can score more than 100 points.
 - but you can't roll over your extra points ⊗

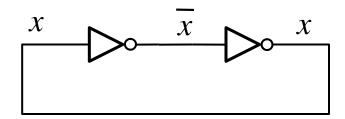
Midterm 2: Topics

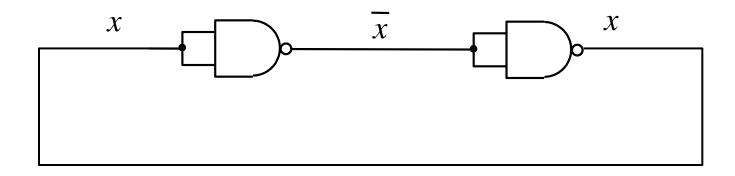
- Binary Numbers and Hexadecimal Numbers
- 1's complement and 2's complement representation
- Addition and subtraction of binary numbers
- Circuits for adders and fast adders
- Single and Double precision IEEE floating point formats
- Converting a real number to the IEEE format
- Converting a floating point number to base 10
- Multiplexers (circuits and function)
- Synthesis of logic functions using multiplexers
- Shannon's Expansion Theorem

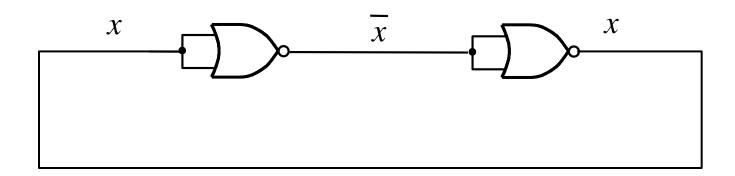
Midterm 2: Topics

- Decoders (circuits and function)
- Demultiplexers
- Encoders (binary and priority)
- Code Converters
- K-maps for 2, 3, and 4 variables
- Synthesis of logic circuits using adders, multiplexers, encoders, decoders, and basic logic gates
- Synthesis of logic circuits given constraints on the available building blocks that you can use
- Latches (circuits, behavior, timing diagrams)
- Flip-Flops (circuits, behavior, timing diagrams)

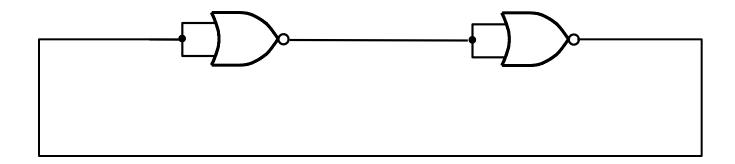
Review of Flip-Flops

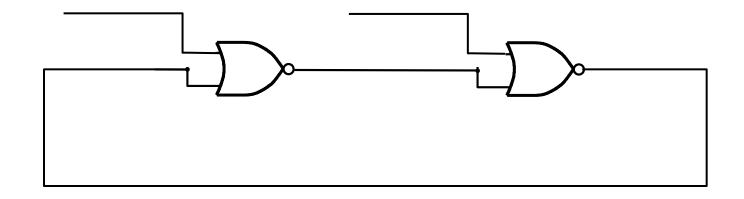


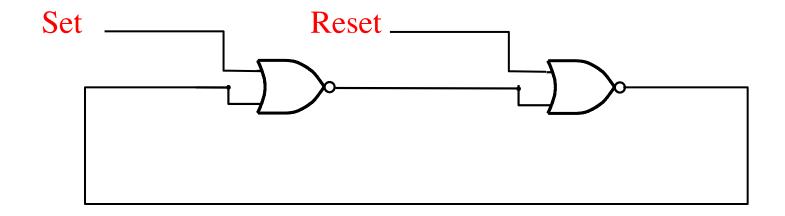




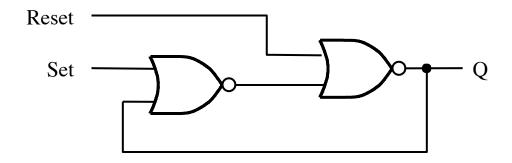
Basic Latch



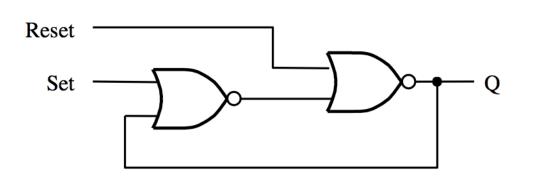


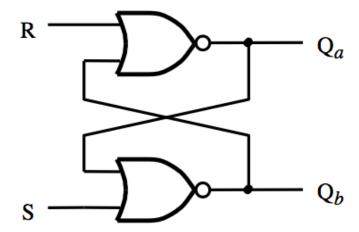


A memory element with NOR gates

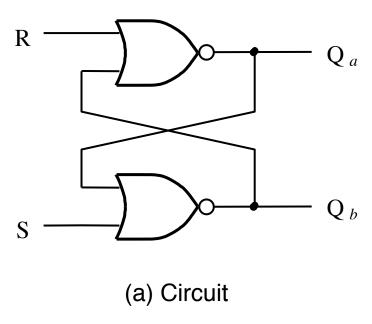


Two Different Ways to Draw the Same Circuit





SR Latch: Circuit and Truth Table

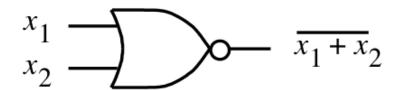


S	R	Q_a	Q_b	_
0	0	0/1	1/0	(no change)
0	1	0	1	
1	0	1	0	
1	1	0	0	(Undesirable)

(b) Truth table

[Figure 5.4a,b from the textbook]

NOR Gate

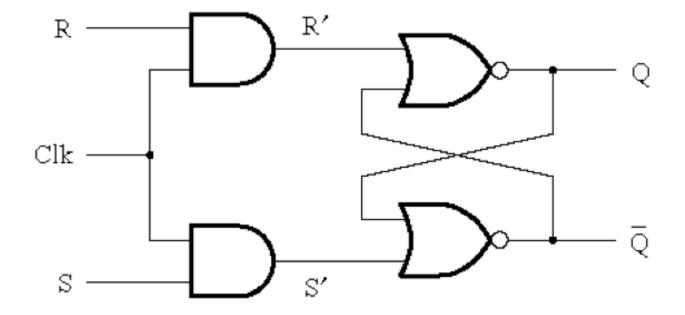


NOR Gate Truth table

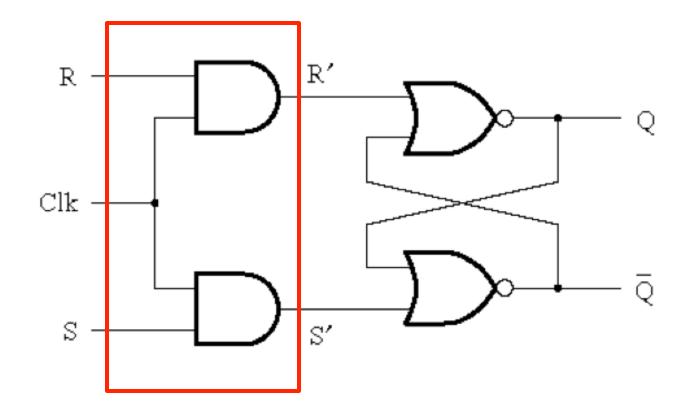
x_1	x_2	f
0	0	1
0	1	0
1	0	0
1	1	0

Gated SR Latch

Circuit Diagram for the Gated SR Latch

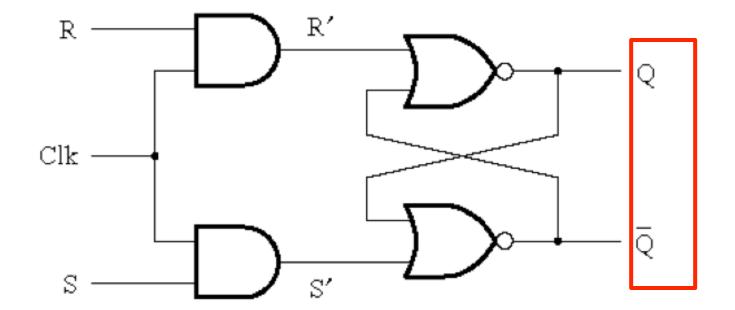


Circuit Diagram for the Gated SR Latch



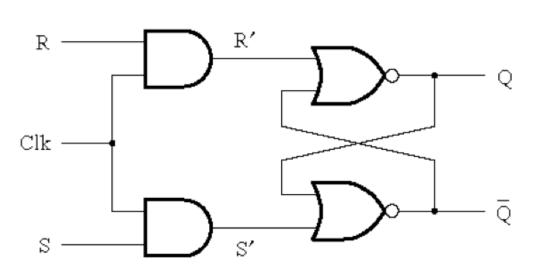
This is the "gate" of the gated latch

Circuit Diagram for the Gated SR Latch

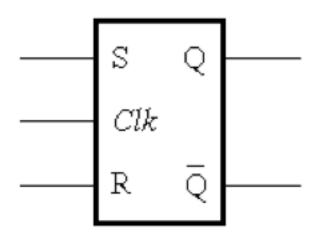


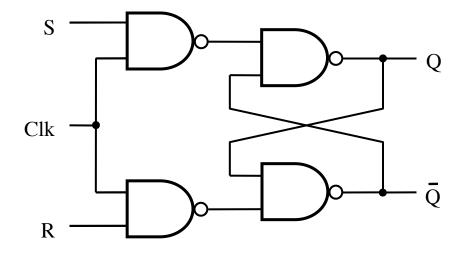
Notice that these are complements of each other

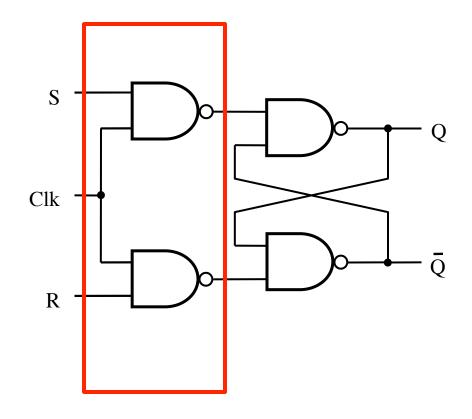
Gated SR Latch: Circuit Diagram, Characteristic Table, and Graphical Symbol



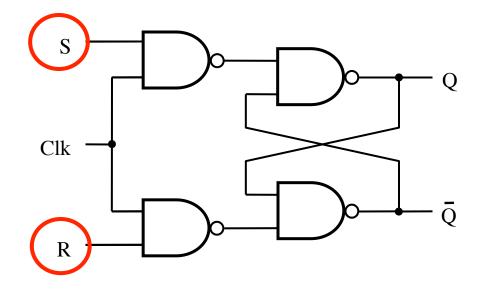
Clk	S	R	Q(t+1)
0	x	x	Q(t) (no change)
1	0	0	Q(t) (no change)
1	0	1	0
1	1	0	1
1	1	1	x (Undesirable)



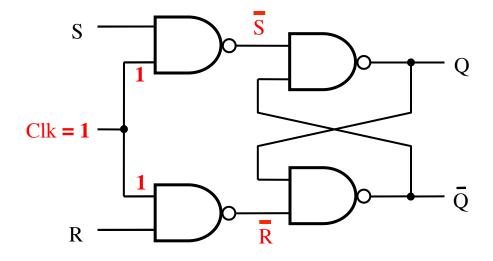




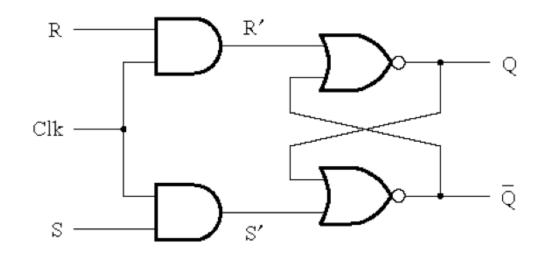
In this case the "gate" is constructed using NAND gates! Not AND gates.



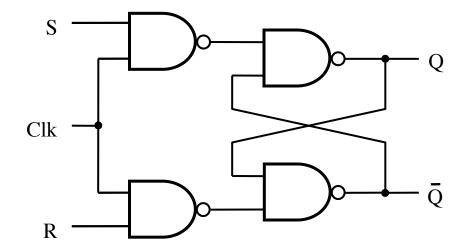
Also, notice that the positions of S and R are now swapped.

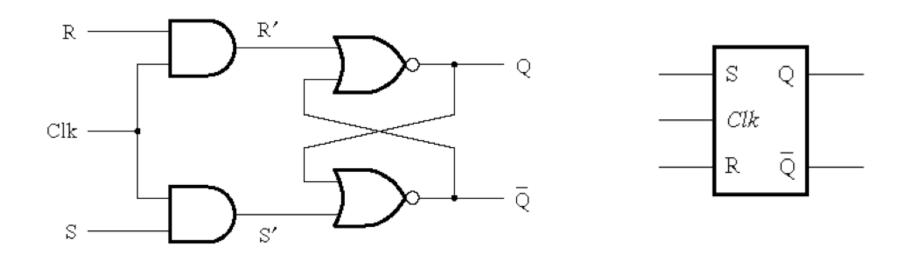


Finally, notice that when Clk=1 this turns into the basic latch with NAND gates, i.e., the \overline{SR} Latch.

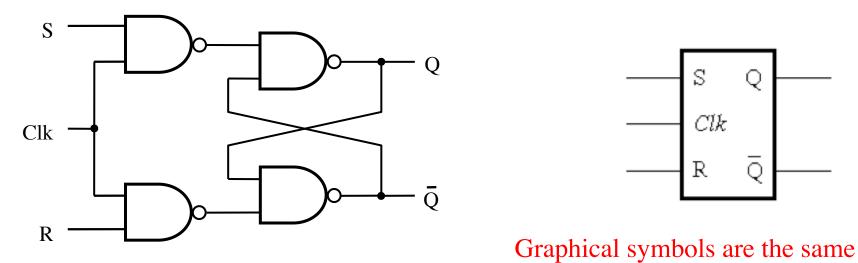


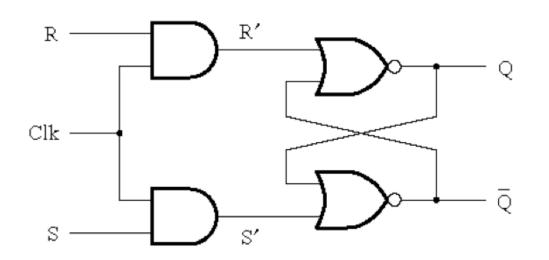
Gated SR latch with NAND gates





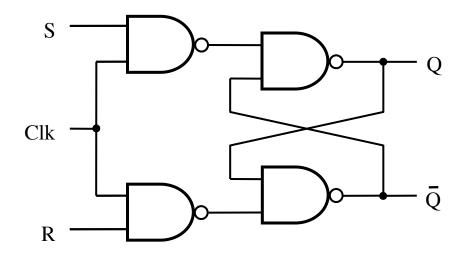
Gated SR latch with NAND gates





Clk	S	R	Q(t+1)
0	x	x	Q(t) (no change)
1	0	0	Q(t) (no change)
1	0	1	0
1	1	0	1
1	1	1	x (undesirable)

Gated SR latch with NAND gates

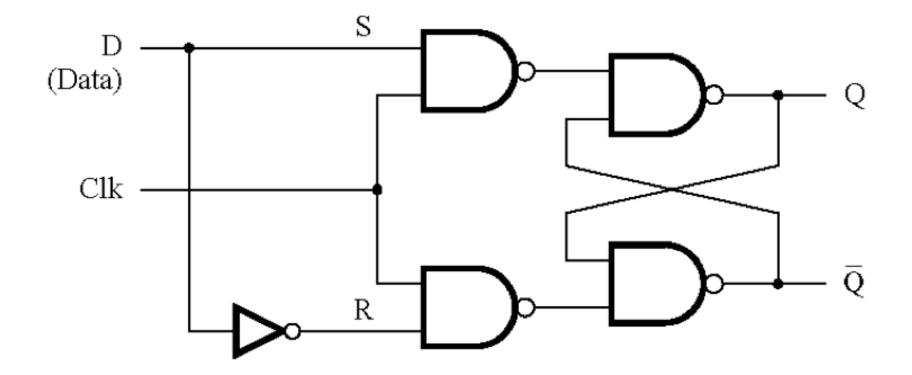


Clk	S	R	Q(t+1)
0	х	x	Q(t) (no change)
1	0	0	Q(t) (no change)
1	0	1	0
1	1	0	1
1	1	1	x (undesirable)

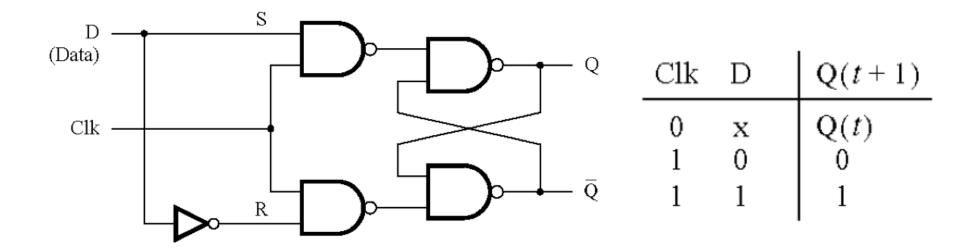
Characteristic tables are the same

Gated D Latch

Circuit Diagram for the Gated D Latch

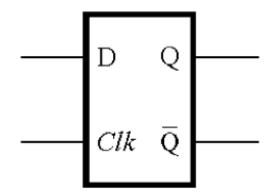


Gated D Latch: Circuit Diagram, Characteristic Table, and Graphical Symbol

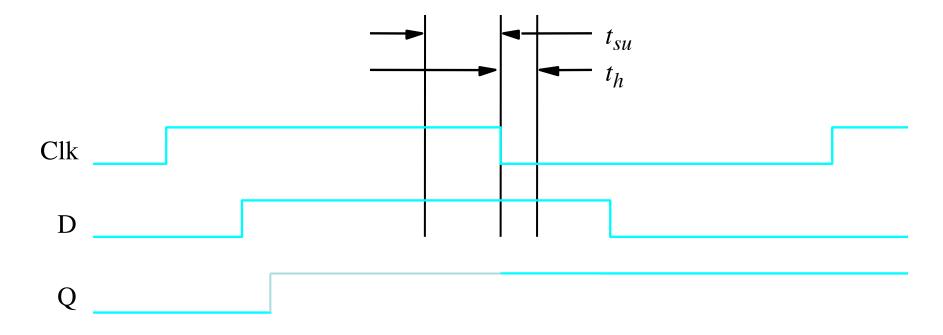


Note that it is now impossible to have S=R=1.

When Clk=1 the output follows the D input. When Clk=0 the output cannot be changed.



Setup and hold times for Gated D latch



Setup time (t_{su}) – the minimum time that the D signal must be stable prior to the negative edge of the Clock signal

Hold time (t_h) – the minimum time that the D signal must remain stable after the negative edge of the Clock signal

Master-Slave D Flip-Flop

Constructing a Master-Slave D Flip-Flop From Two D Latches

Master Slave

(Data)

Clk

R

Q

Clk

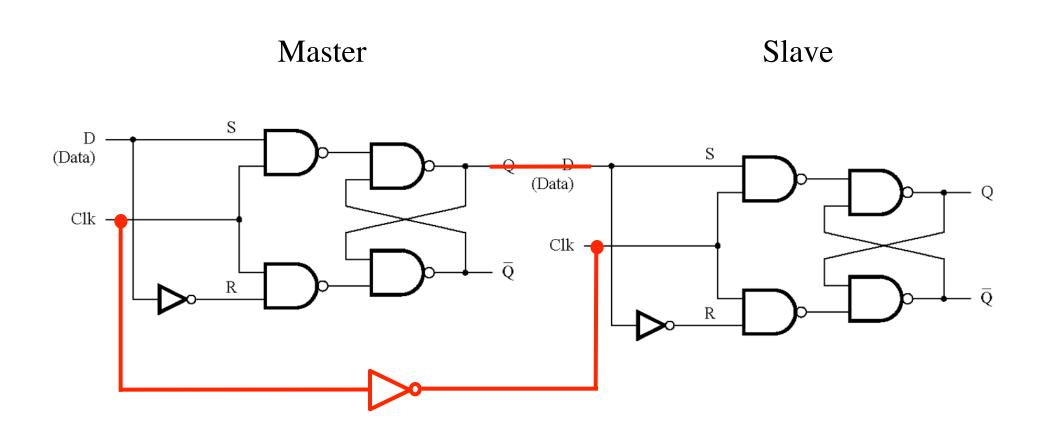
Q

Clk

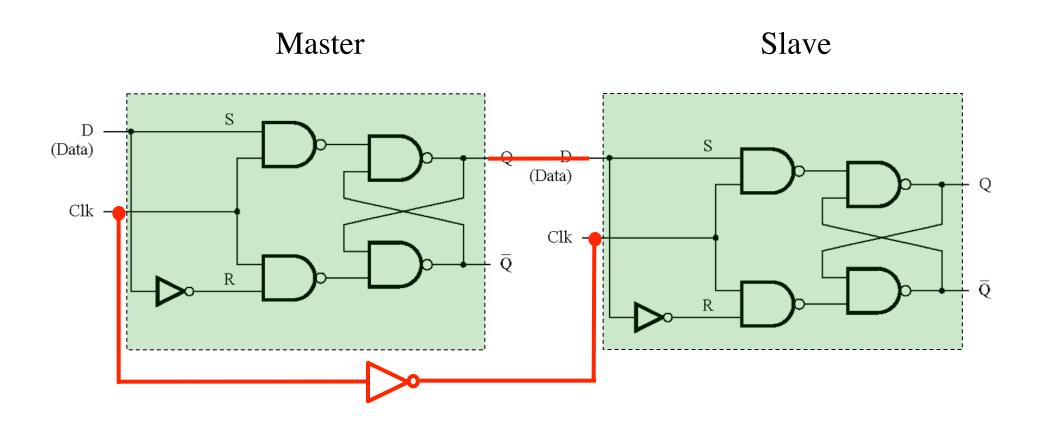
R

R

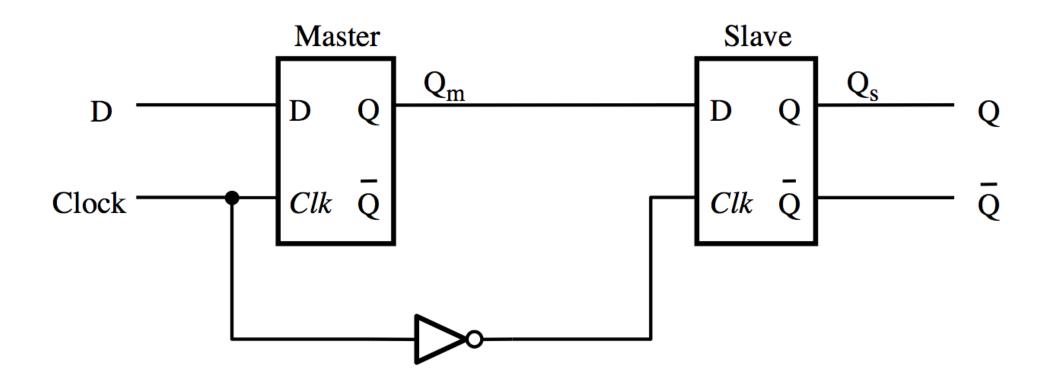
Constructing a Master-Slave D Flip-Flop From Two D Latches



Constructing a Master-Slave D Flip-Flop From Two D Latches



Constructing a Master-Slave D Flip-Flop From Two D Latches



Constructing a Master-Slave D Flip-Flop From one D Latch and one Gated SR Latch

(This version uses one less NOT gate)

Master Slave

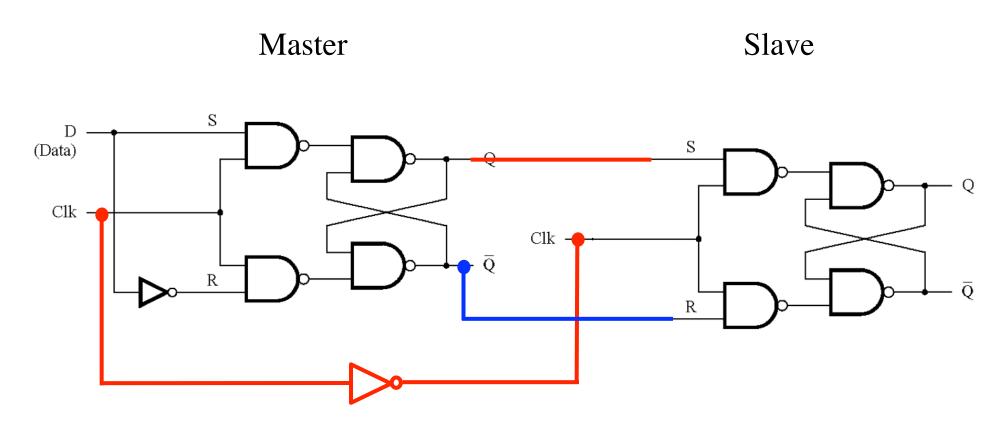
(Data)

Clk

R

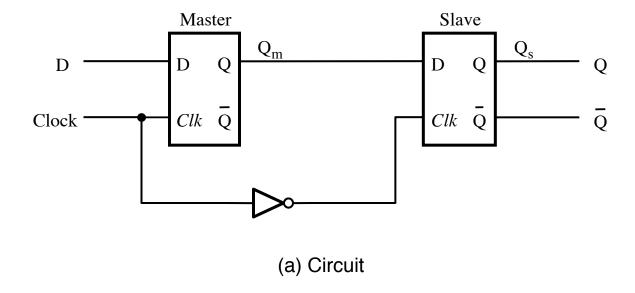
Constructing a Master-Slave D Flip-Flop From one D Latch and one Gated SR Latch

(This version uses one less NOT gate)

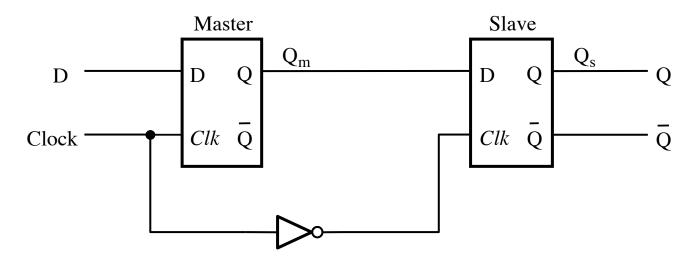


Edge-Triggered D Flip-Flops

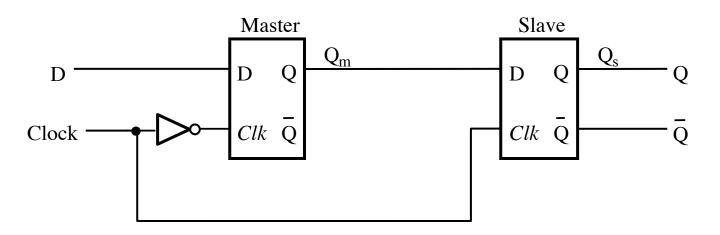
Master-Slave D Flip-Flop



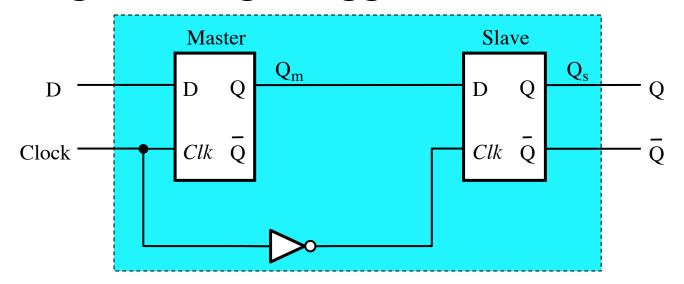
Negative-Edge-Triggered Master-Slave D Flip-Flop



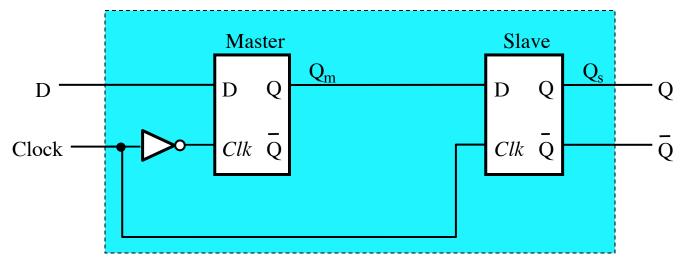
Positive-Edge-Triggered Master-Slave D Flip-Flop



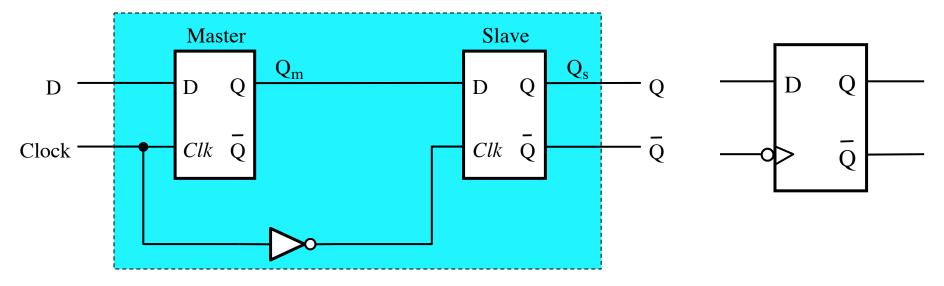
Negative-Edge-Triggered Master-Slave D Flip-Flop



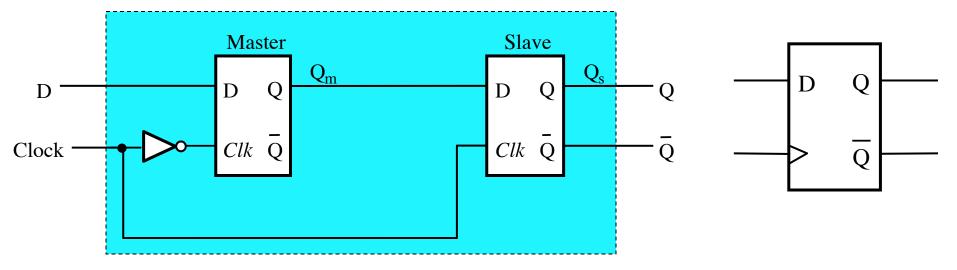
Positive-Edge-Triggered Master-Slave D Flip-Flop

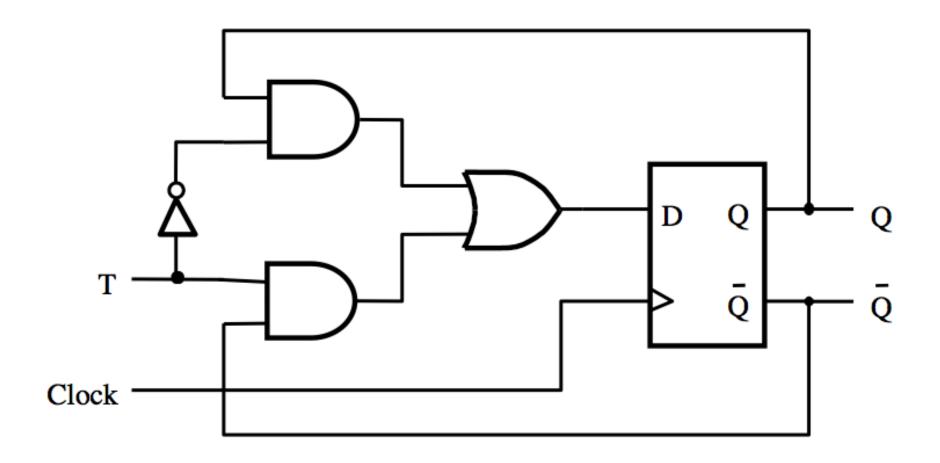


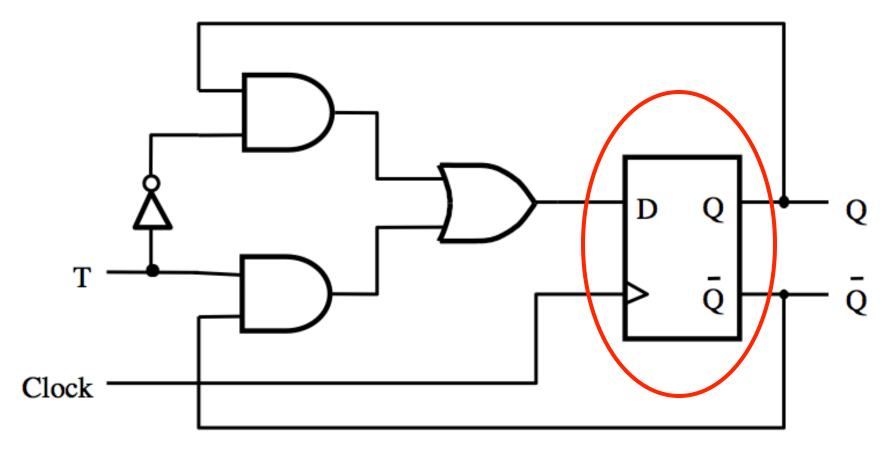
Negative-Edge-Triggered Master-Slave D Flip-Flop



Positive-Edge-Triggered Master-Slave D Flip-Flop

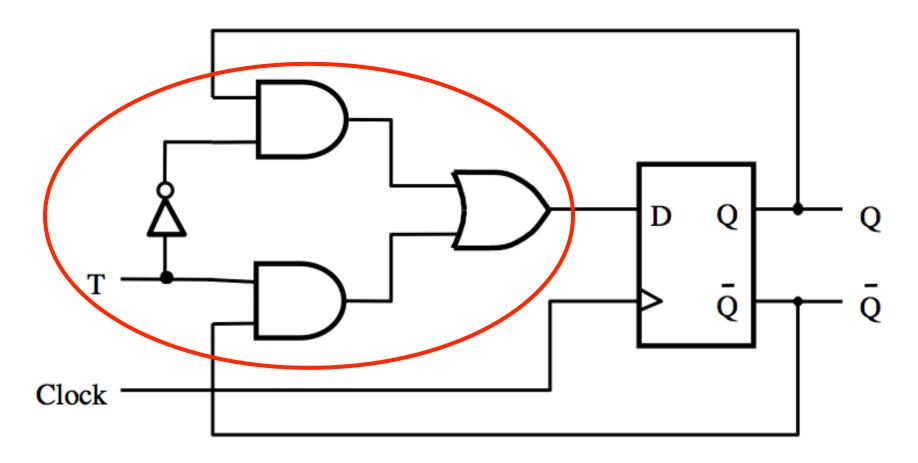


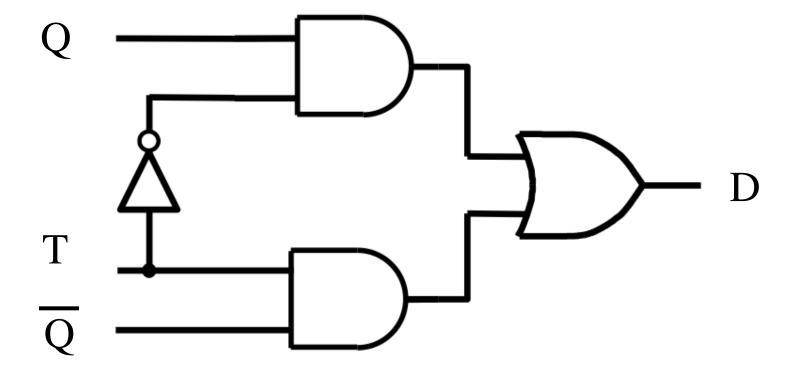


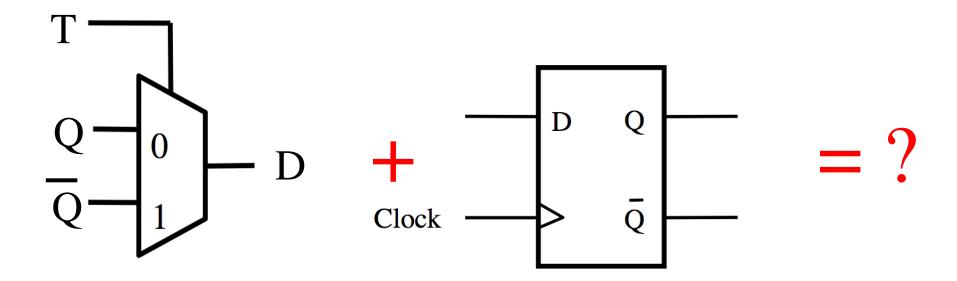


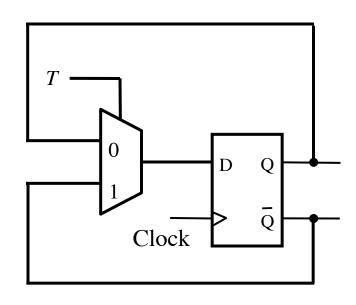
Positive-edge-triggered D Flip-Flop

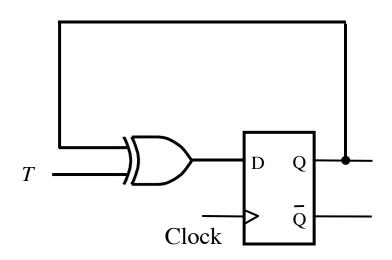
[Figure 5.15a from the textbook]



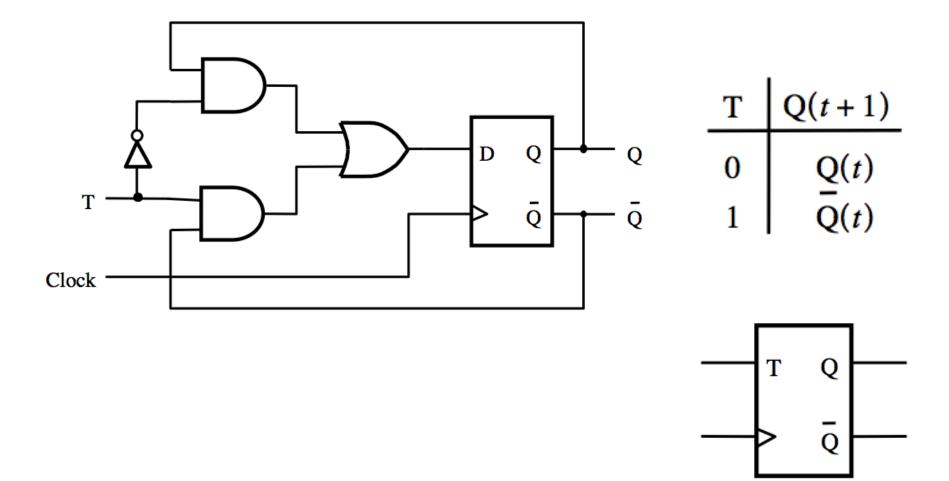








T Flip-Flop (circuit, truth table and graphical symbol)



[Figure 5.15a-c from the textbook]

T Flip-Flop (How it Works)

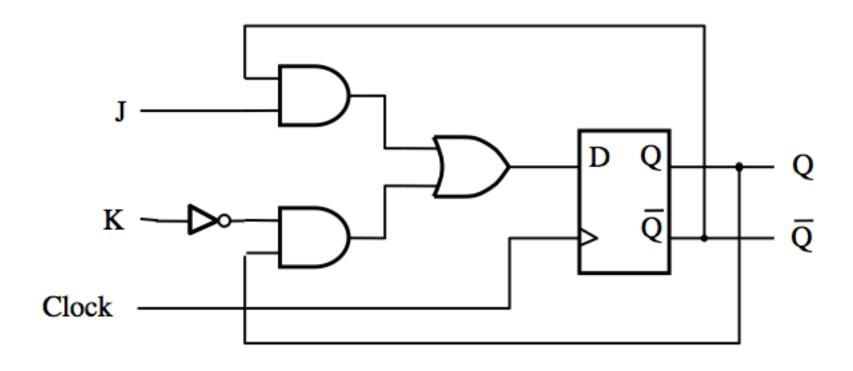
If T=0 then it stays in its current state

If T=1 then it reverses its current state

In other words the circuit "toggles" its state when T=1. This is why it is called T flip-flop.

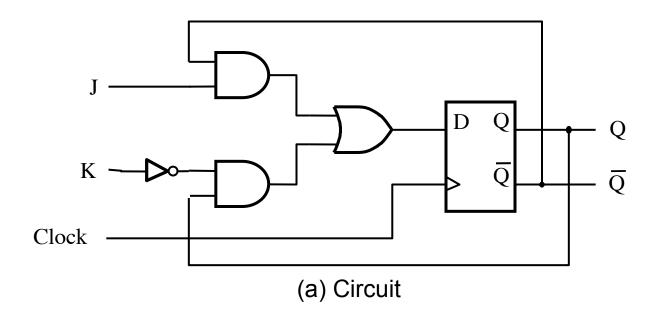
JK Flip-Flop

JK Flip-Flop



$$D = \overline{JQ} + \overline{KQ}$$

JK Flip-Flop



J K	Q(t+1)	
0 0	Q(t)	JO
0 1	0	
1 0	1	$ \begin{bmatrix} \kappa & \overline{\alpha} \end{bmatrix}$
1 1	$\overline{\mathbf{Q}}\left(\mathbf{t}\right)$	IX Q

(b) Truth table

(c) Graphical symbol

JK Flip-Flop (How it Works)

A versatile circuit that can be used both as a SR flip-flop and as a T flip flop

If J=0 and S =0 it stays in the same state

Just like SR It can be set and reset J=S and K=R

If J=K=1 then it behaves as a T flip-flop

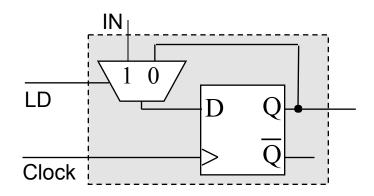
Registers

Register (Definition)

An n-bit structure consisting of flip-flops

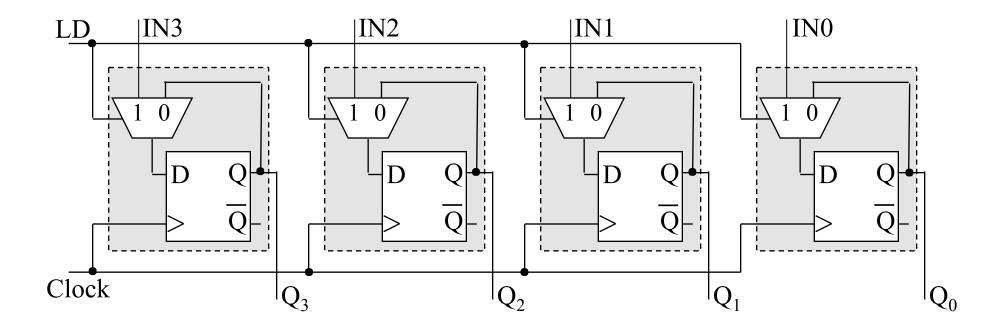
Parallel-Access Register

1-bit Parallel-access register



- At the input of the D flip-flop, a 2-to-1 Multiplexer is used to select whether to load a new input value or to retain the old value
- If signal LD = 1 then load the new value
- If signal LD = 0 then retain the old value

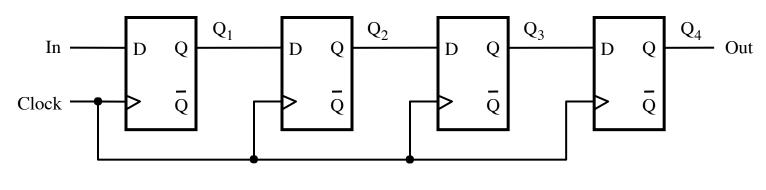
4-bit Parallel-access register



Notice that all flip-flops are on the same clock cycle.

Shift Register

A simple shift register



(a) Circuit

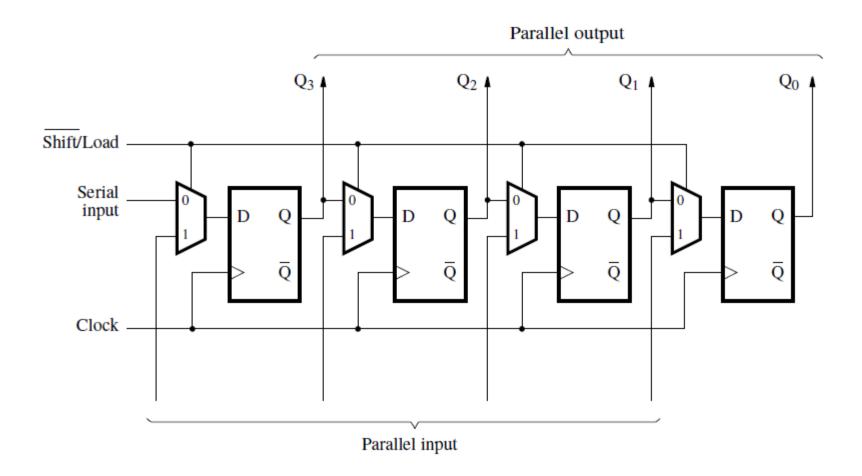
In
$$Q_1$$
 Q_2 Q_3 $Q_4 = Out$
 t_0 1 0 0 0 0
 t_1 0 1 0 0 0
 t_2 1 0 1 0 0
 t_3 1 1 0 1 0
 t_4 1 1 1 0 1
 t_5 0 1 1 1 0
 t_6 0 0 1 1 1
 t_7 0 0 0 0 1 1

(b) A sample sequence

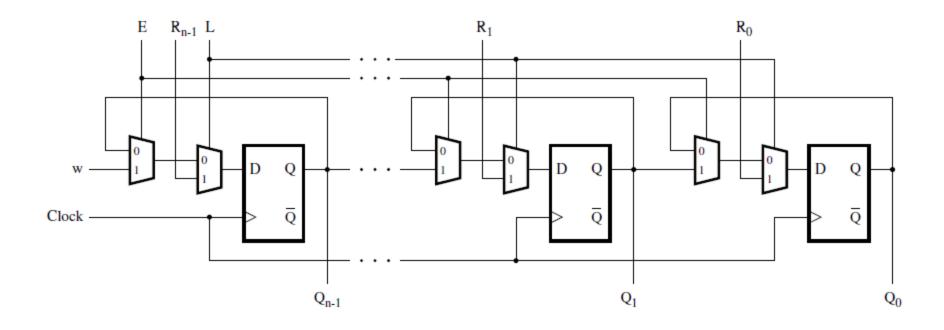
[Figure 5.17 from the textbook]

Parallel-Access Shift Register

Parallel-access shift register



A shift register with parallel load and enable control inputs



Register File

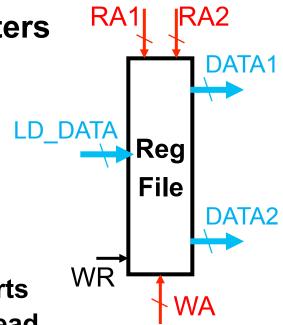
Register File

Register file is a unit containing r registers

r can be 4, 8, 16, 32, etc.

Each register has n bits

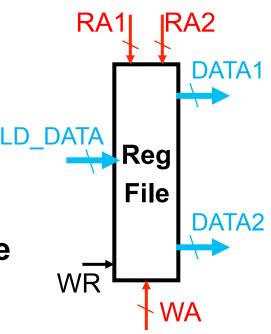
- n can be 4, 8, 16, 32, etc.
- n defines the data path width
- Output ports (DATA1 and DATA2) are used for reading the register file
 - Any register can be read from any of the ports
 - Each port needs a log₂r bits to specify the read address (RA1 and RA2)
- Input port (LD_DATA) is used for writing data to the register file
 - Write address is also specified by log₂r bits (WA)
 - Writing is enabled by a 1-bit signal (WR)



Register File: Exercise

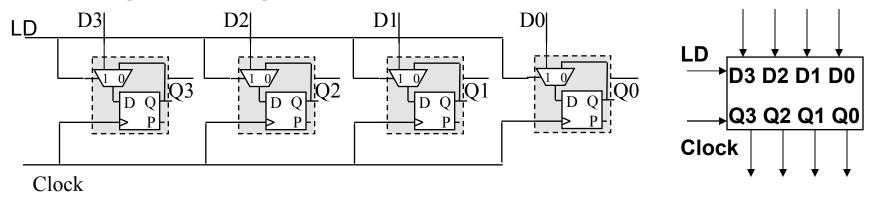
- Suppose that a register file
 - contains 32 registers
 - width of data path is 16 bits(i.e., each register has 16 bits)

- How many bits are there for each of the signals?
 - RA1 5
 - RA2 5
 - DATA1 16
 - DATA216
 - WA 5
 - LD_DATA 16
 - WR 1

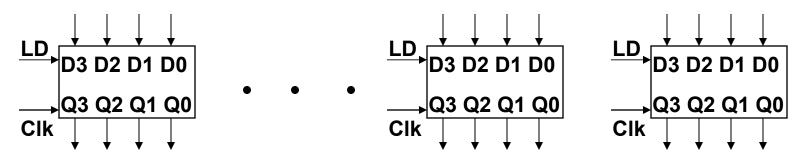


Register file design

- We will design an eight-register file with 4-bit wide registers
- A single 4-bit register and its abstraction are shown below



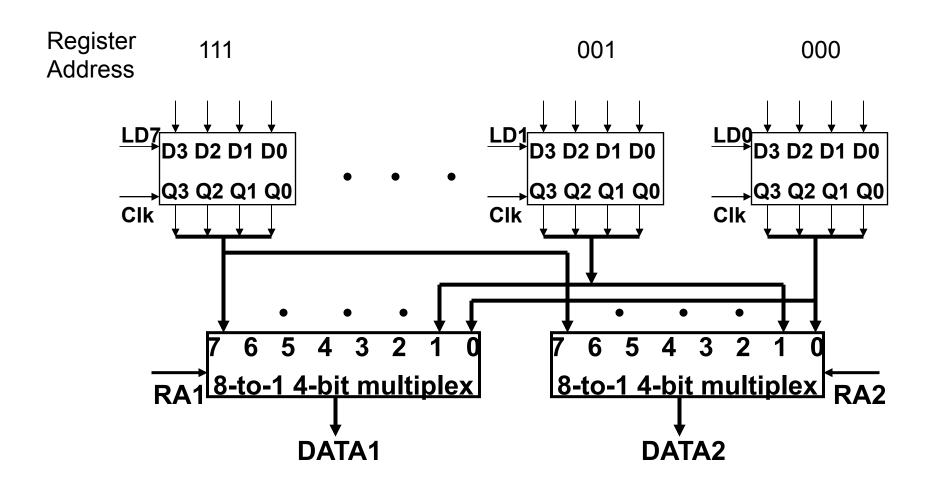
We have to use eight such registers to make an eight register file



How many bits are required to specify a register address?

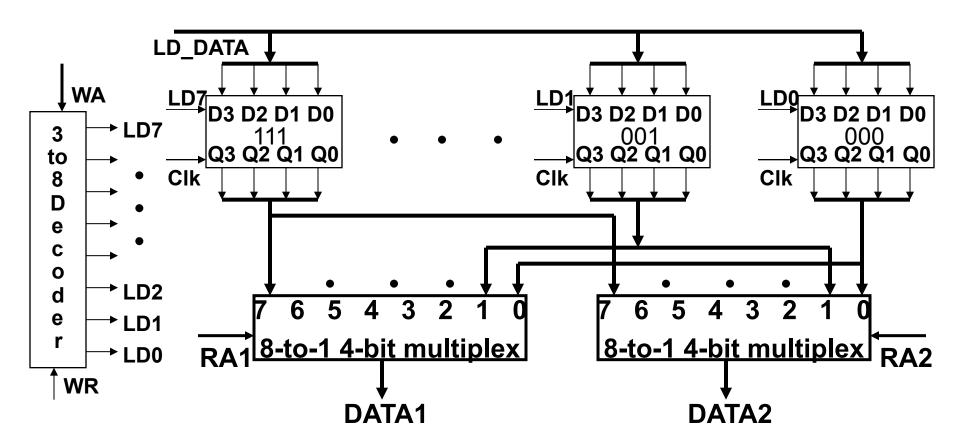
Reading Circuit

- A 3-bit register address, RA, specifies which register is to be read
- For each output port, we need one 8-to-1 4-bit multiplier



Adding write control to register file

- To write to any register, we need the register's address (WA) and a write register signal (WR)
- A 3-bit write address is decoded if write register signal is present
- One of the eight registers gets a LD signal from the decoder



Questions?

THE END