

## **CprE 281: Digital Logic**

**Instructor: Alexander Stoytchev** 

http://www.ece.iastate.edu/~alexs/classes/

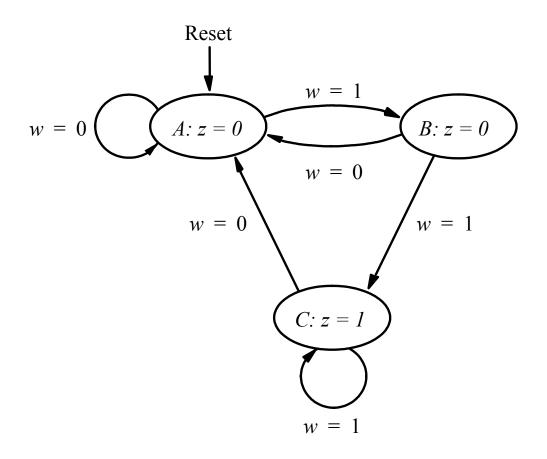
# Synchronous Sequential Circuits Basic Design Steps

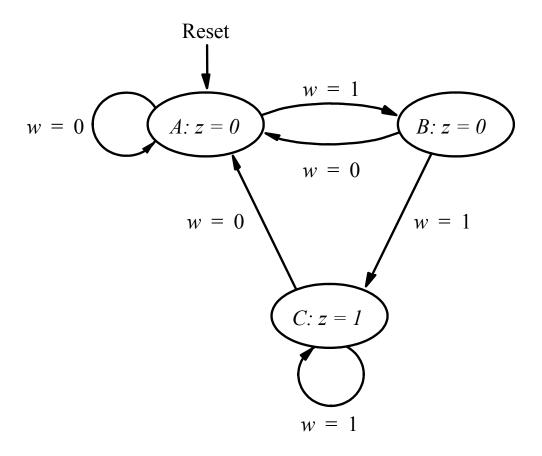
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## First Design Pattern: Moore Machines

#### **Moore Machine:**

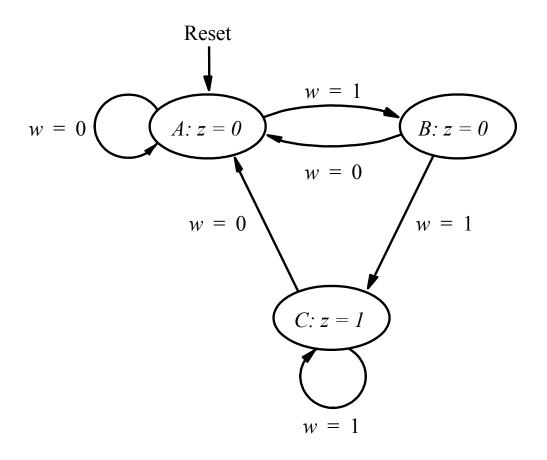
#### A Type of Finite State Machine (FSM)



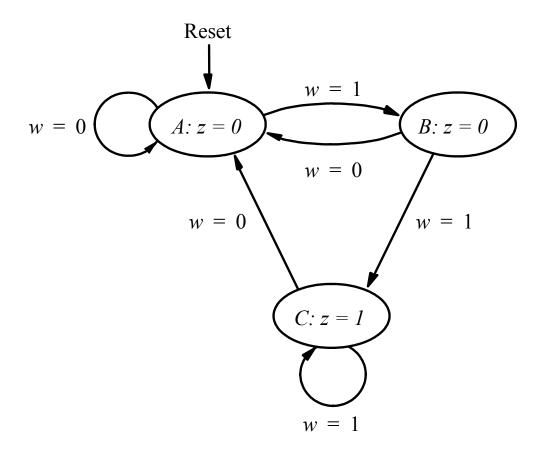


- Finite number of states (nodes).
- Discrete state transitions (edges).
- Only "in" one state at a time.
- One reset state
- Every state has an outgoing state transition for each possible input.

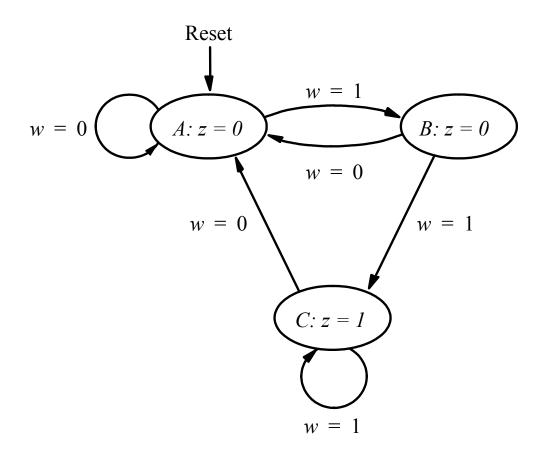
[ Figure 6.3 from the textbook ]



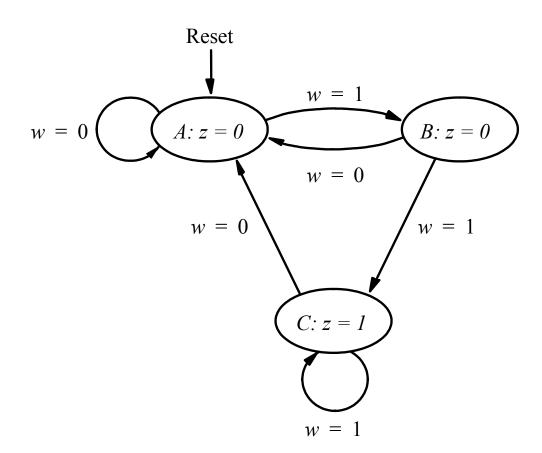
**Key:** Next state depends on both current state and current input



Key: Output only depends on current state.

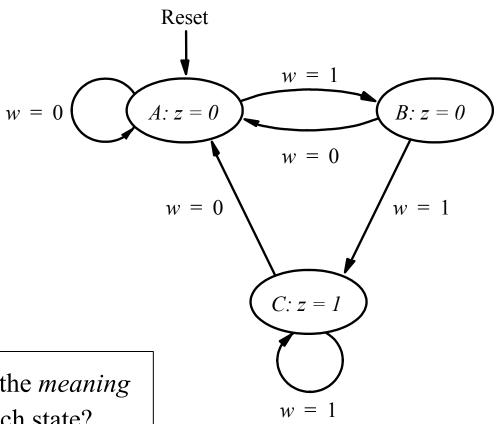


Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
	_		_			_				_	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



In general, we need to start tracing from the beginning to know which state the FSM is in. It may not be clear from a short sequence of outputs.

Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	t <sub>10</sub>
	_	_	_	_	_	_	_	_	_	0	_
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



What is the *meaning* of each state?

Clockcycle:	$t_0$	$t_1$	$t_2$	t <sub>3</sub>	t <sub>4</sub>	t <sub>5</sub>	$t_6$	t <sub>7</sub>	$t_8$	t <sub>9</sub>	$t_{10}$
											1
											0

#### What is a State?

It is not really a memory of every past input.

(We might run out of space to remember it all!)

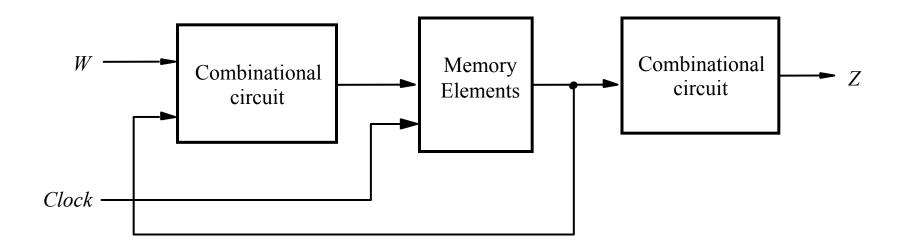
Rather, it is a characterization or snapshot of the pattern of inputs which have come before.

## Moore Machine Implementation

The state diagram is just an illustration to help us describe and reason about how the FSM will behave in each of its states.

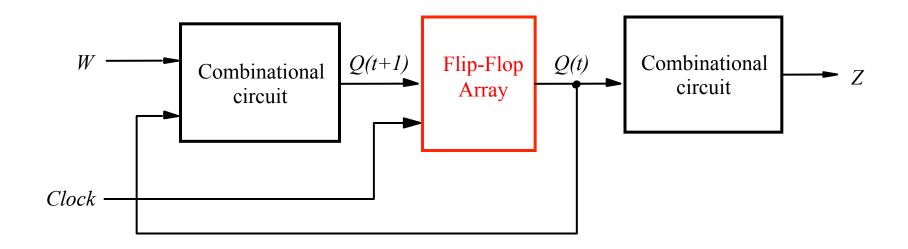
So, how do we turn it into a circuit?

## **Moore Machine Implementation**



Note: The *W* and *Z* lines need not be wires. They can be buses.

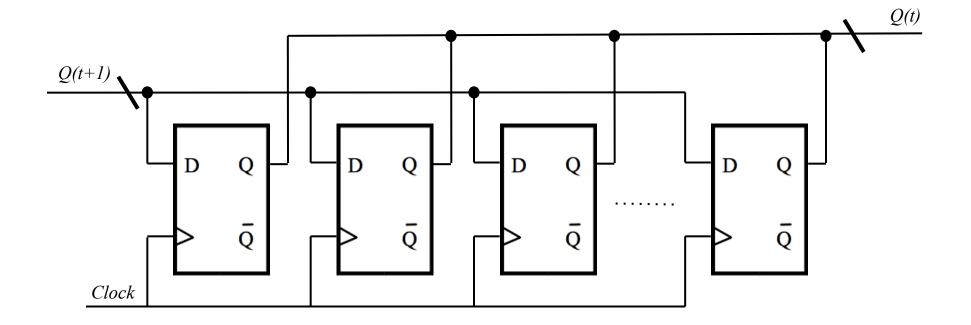
### **State Storage**



Any usable "memory" of the preceding input sequence is encoded in the flip-flop array.

#### **FSM States**

The Flip-Flop array stores an encoding of the current state.



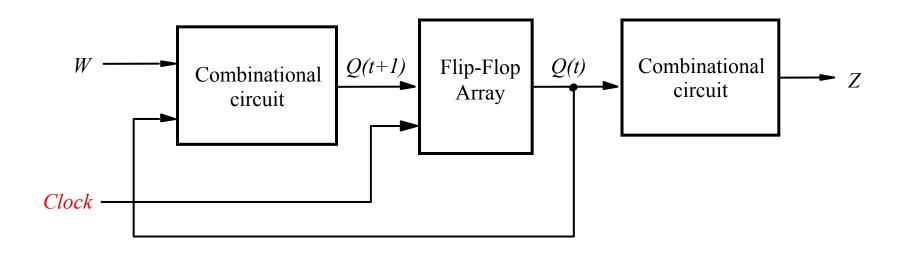
## **State Encoding**

Each of the states in our design is identified by a distinct code.

If we use 3 flip-flops, then the FSM can have up to  $2^3 = 8$  distinct states.

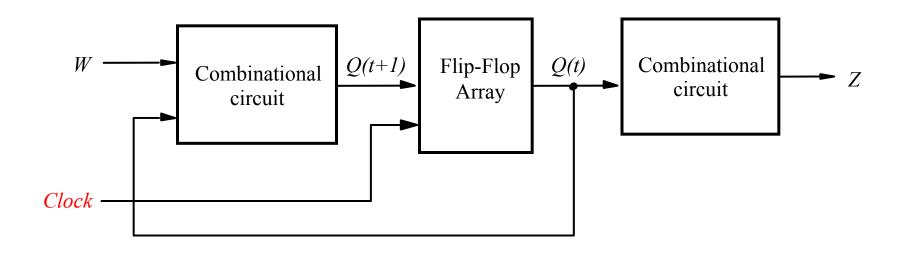
So, when the flip-flop array contains the code 011, we say that the machine is in state 011.

## Synchronous Design



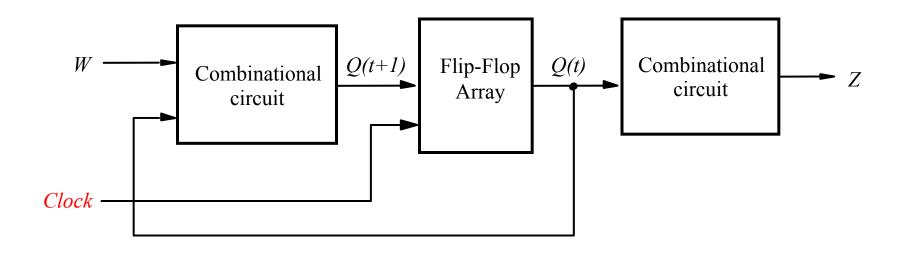
Every active clock edge causes a state transition.

### Synchronous Design



We expect the input signals to be stable before the *active* clock edge occurs.

## Synchronous Design



There is a whole other class of sequential circuits which are asynchronous, but we will not study them in this course.

#### Sequential Circuits: Key Ideas

The current output depends on something about the preceding sequence of inputs (and maybe the current output).

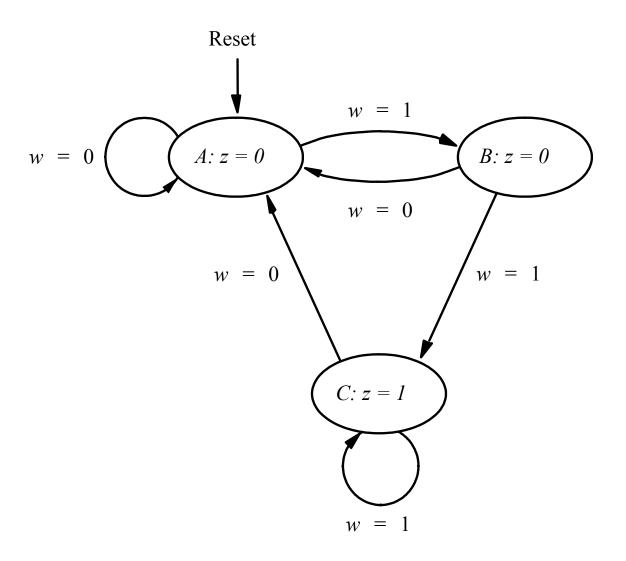
Using *memory elements* (i.e., flip-flops), we design the circuit to remember some *relevant* information about the prior inputs.

#### In general, next output depends on both:

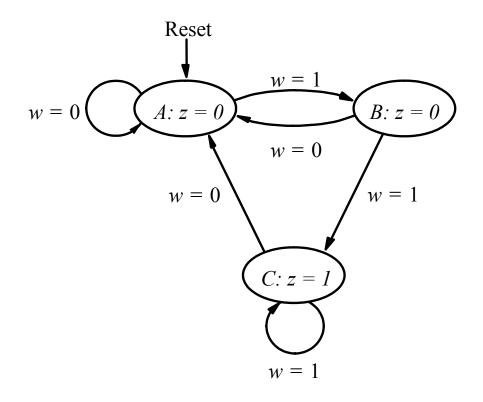
some pattern in the past inputs

• current input

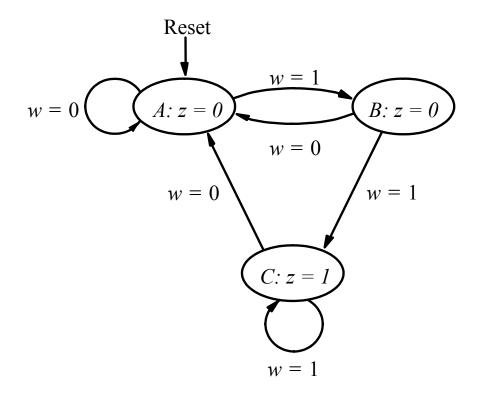
### **Example**



We need to find both the *next state logic* and the *output logic* implied by this machine.



Present	Next state	Output		
state	w = 0 $w = 1$	Z		
A				
В				
С				



Present	Next	Output		
state	w = 0	w = 1	$\overline{z}$	
A	A	В	0	
В	A	C	0	
C	A	C	1	

Figure 6.4 from the textbook ]

#### How to represent the States?

One way is to encode each state with a 2-bit binary number

A ~ 00

B~01

C ~ 10

#### How to represent the states?

One way is to encode each state with a 2-bit binary number

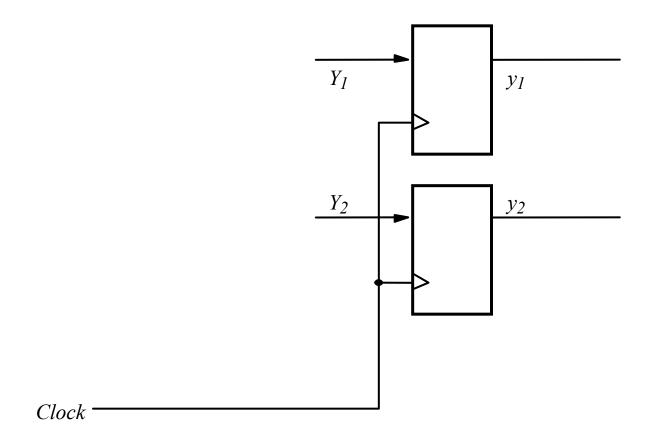
A ~ 00

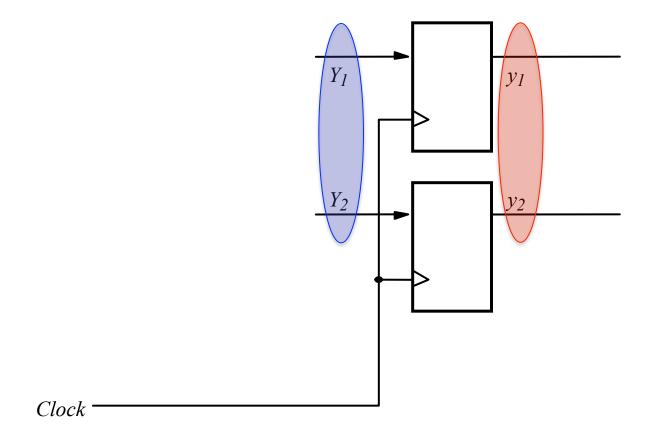
B ~ 01

C ~ 10

How many flip-flops do we need?

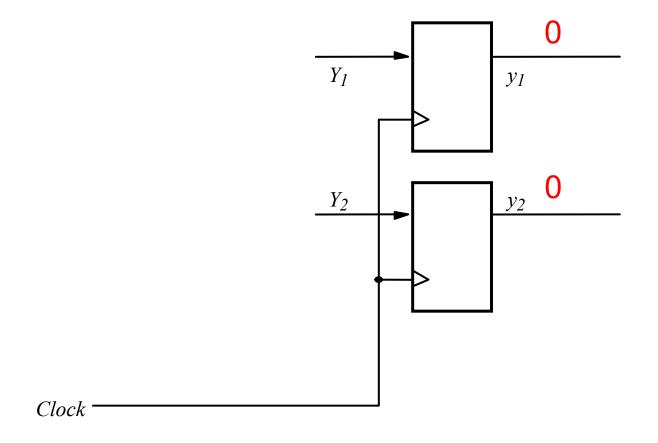
## Let's use two flip flops to hold the machine's state



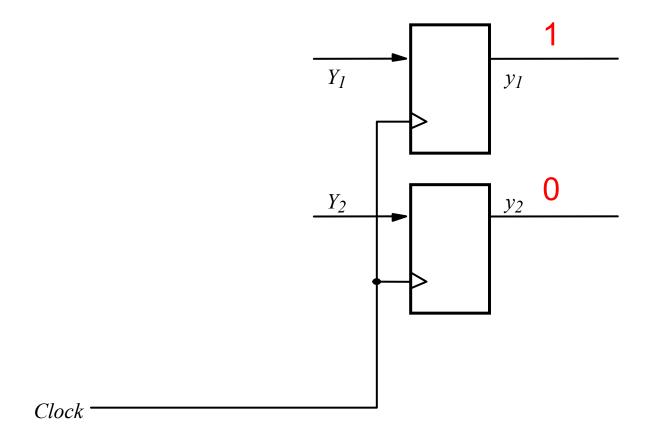


We will call  $y_1$  and  $y_2$  the present state variables.

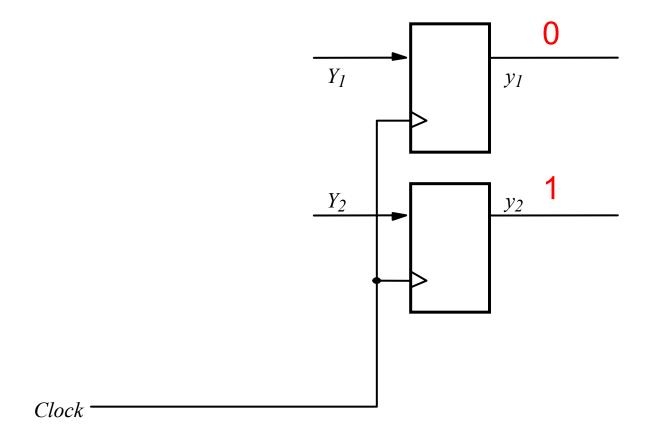
We will call  $Y_1$  and  $Y_2$  the next state variables.



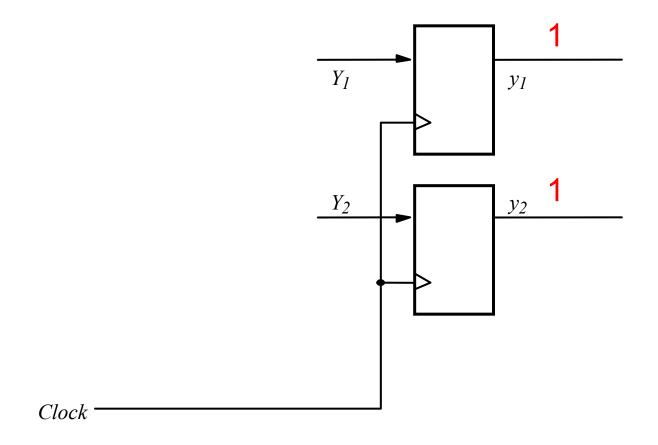
Two zeros on the output JOINTLY represent state A.



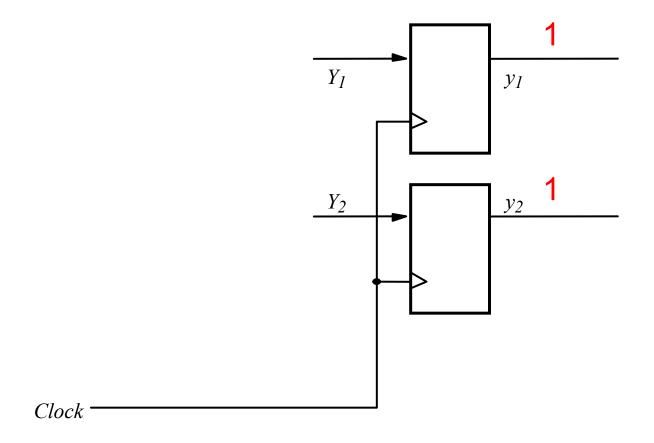
This flip-flop output pattern represents state B.



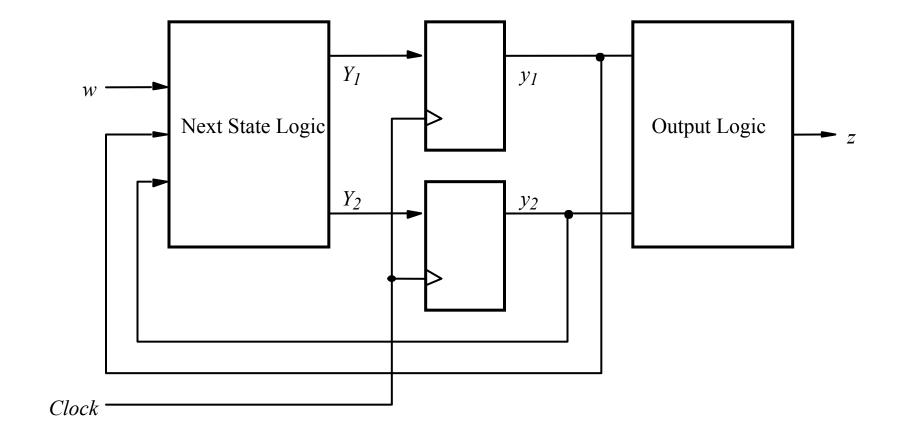
This flip-flop output pattern represents state C.

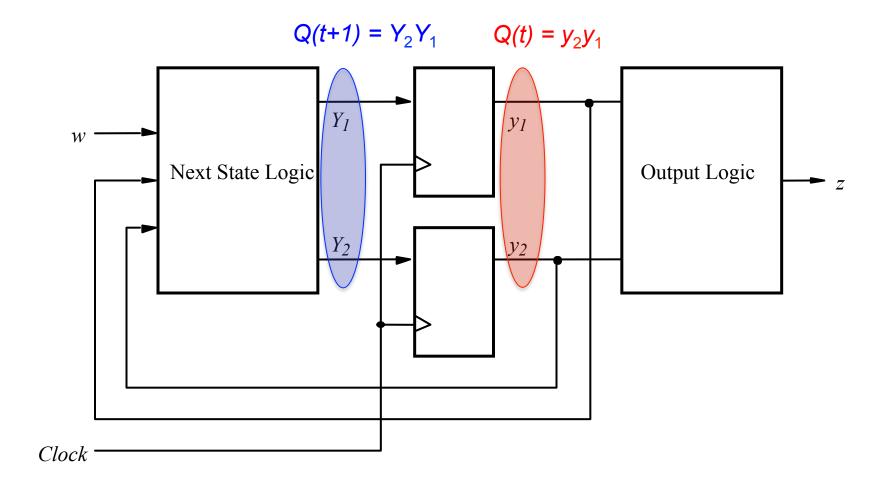


What does this flip-flop output pattern represent?



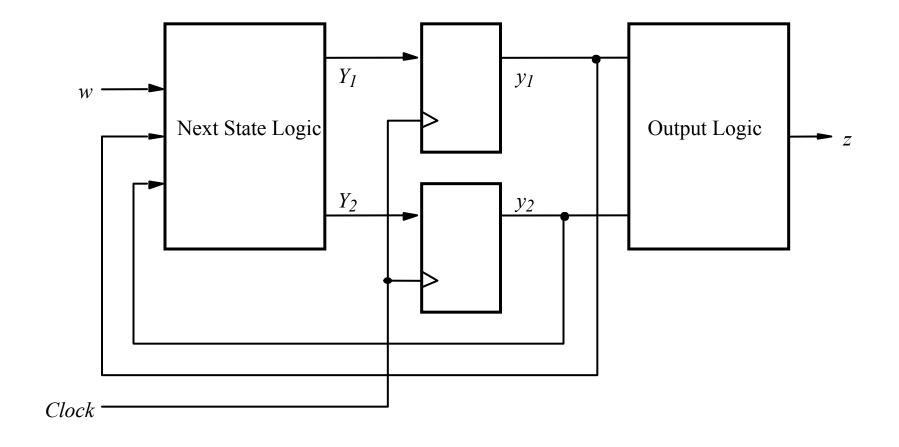
This would be state D, but we don't have one in this example. So this is an impossible state.



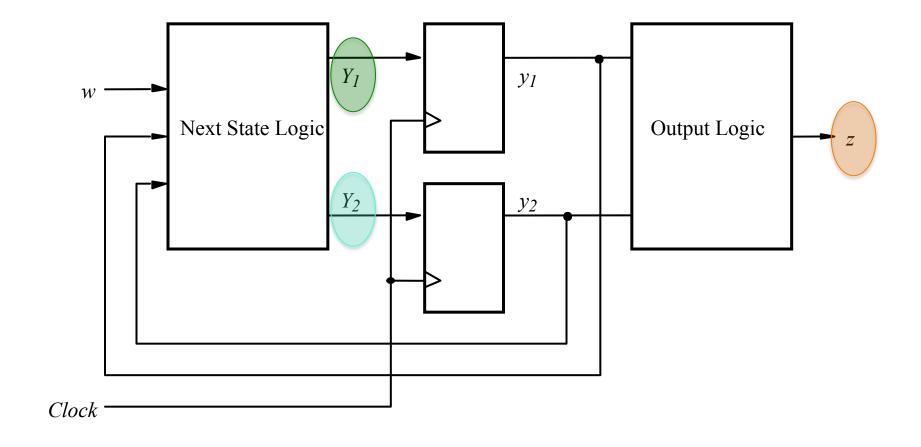


We will call  $y_1$  and  $y_2$  the present state variables.

We will call  $Y_1$  and  $Y_2$  the next state variables.



We need to find logic expressions for  $Y_1(w, y_1, y_2)$ ,  $Y_2(w, y_1, y_2)$ , and  $z(y_1, y_2)$ .



We need to find logic expressions for  $Y_1(w, y_1, y_2)$ ,  $Y_2(w, y_1, y_2)$ , and  $z(y_1, y_2)$ .

Present	Next	Output	
state	w = 0	w = 1	Z
A	A	В	0
В	A	C	0
C	A	C	1

## Suppose we encoded our states in the same order in which they were labeled:

A ~ 00

B ~ 01

C ~ 10

Present	Next state		Output
state	w = 0	w = 1	$\overline{z}$
A	A	В	0
В	A	C	0
C	A	C	1

	Present	Next state	
	state	w = 0 $w = 1$	Output
			Z
Α	00		
В	01		
C	10		
	11		

The finite state machine will never reach a state encoded as 11.

[ Figure 6.6 from the textbook ]

Present	Next	Output	
state	w = 0	w = 1	Z
A	A	В	0
В	A	C	0
C	A	C	1

	Present	Next state		
	state	w = 0 $w = 1$		Output
	<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$Y_2Y_1$	Z
A	00	00	01	0
В	01	00	10	0
C	10	00	10	1
,	11	dd	dd	d

We arbitrarily chose these as our state encodings.
We could have used others.

[ Figure 6.6 from the textbook ]

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

Present	Next s		
state	w = 0	w = 1	Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w	$y_2$	$y_I$	$Y_2$	$Y_I$
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

$y_2$	$y_I$	Z
0	0	
0	1	
1	0	
1	1	

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

Present	Next s		
state	w = 0	Output	
<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$ $Y_2Y_1$		Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w	$y_2$	$y_I$	$Y_2$	$Y_I$
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

			D.
Present	Next s		
state	w = 0	w = 1	Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w	$y_2$	$y_I$	$Y_2$	$Y_I$
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

$y_2$	$y_1$	Z
0	0	0
0	1	0
1	0	1
1	1	d

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

Present	Next s	tate	
state	w = 0	w = 1	Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w		$y_2$	$y_I$	$Y_2$	$Y_I$
0		0	0	0	
0		0	1	0	
0		1	0	0	
0		1	1	d	
1		0	0		
1		0	1		
1	Ī	1	0		
1		1	1		

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

Present			
state	w = 0	Output	
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$ $Y_2Y_1$		Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w	$y_2$	$y_I$	$Y_2$	$Y_{I}$
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	d	
1	0	0	0	
1	0	1	1	
1	1	0	1	
1	1	1	d	

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

			D.
Present	Next s		
state	w = 0	w = 1	Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w	$y_2$	$y_I$	$Y_2$	$Y_I$
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	d	
1	0	0	0	
1	0	1	1	
1	1	0	1	
1	1	1	d	

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

Present	Next s	tate	
state	w = 0	w = 1	Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$y_2y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w		$y_2$	$y_I$	$Y_2$	$Y_I$
0		0	0	0	0
0		0	1	0	0
0		1	0	0	0
0		1	1	d	d
1		0	0	0	
1		0	1	1	
1		1	0	1	
1	Ī	1	1	d	

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

Present	Next s		
state	w = 0	w = 1	Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$y_2y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

w	$y_2$	$y_I$	$Y_2$	$Y_I$
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	d	d
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	d	d

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

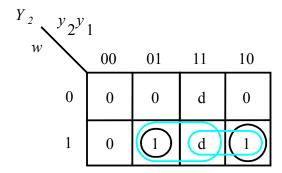
Present	Next s	xt state	
state	w = 0	w = 1	Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

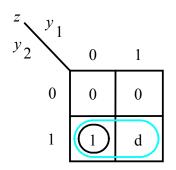
w	$y_2$	$y_I$	$Y_2$	$Y_I$
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	d	d
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	d	d

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

Note that the textbook draws these K-Maps differently from all previous K-maps (the least significant bits index the columns, instead of the most significant bits).

$Y_1$ $y_2$	1			
w	00	01	11	10
0	0	0	d	0
1	1	0	d	0



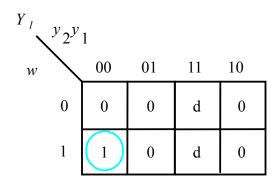


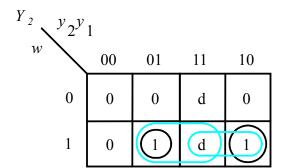
$$Q(t) = y_2 y_1 \text{ and } Q(t+1) = Y_2 Y_1$$

w	$y_2$	$y_I$	$Y_2$	$Y_I$
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	d	d
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	d	d

$y_2$	$y_I$	Z
0	0	0
0	1	0
1	0	1
1	1	d

## Don't care conditions simplify the combinatorial logic





Ignoring don't cares

$$Y_1 = w\overline{y_1}\overline{y_2}$$

$$Y_2 = wy_1\overline{y}_2 + \overline{w}y_1y_2$$

$$z = \overline{y}_1 y_2$$

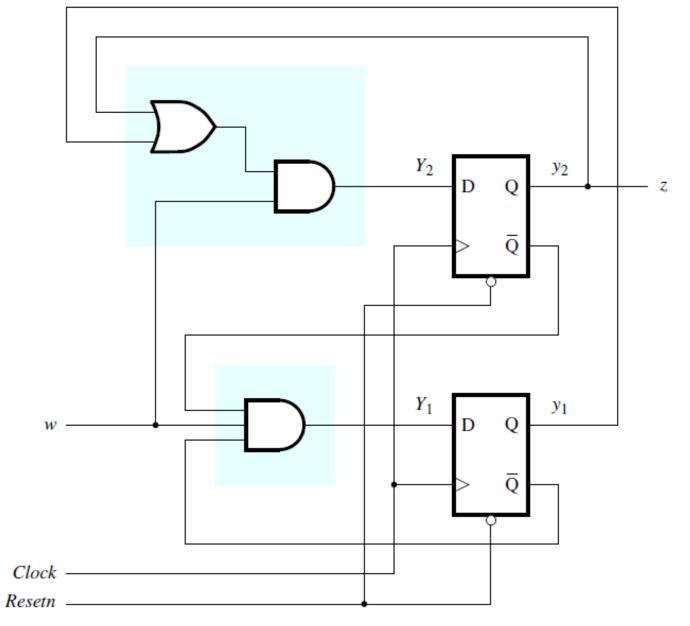
Using don't cares

$$Y_1 = w\overline{y_1}\overline{y_2}$$

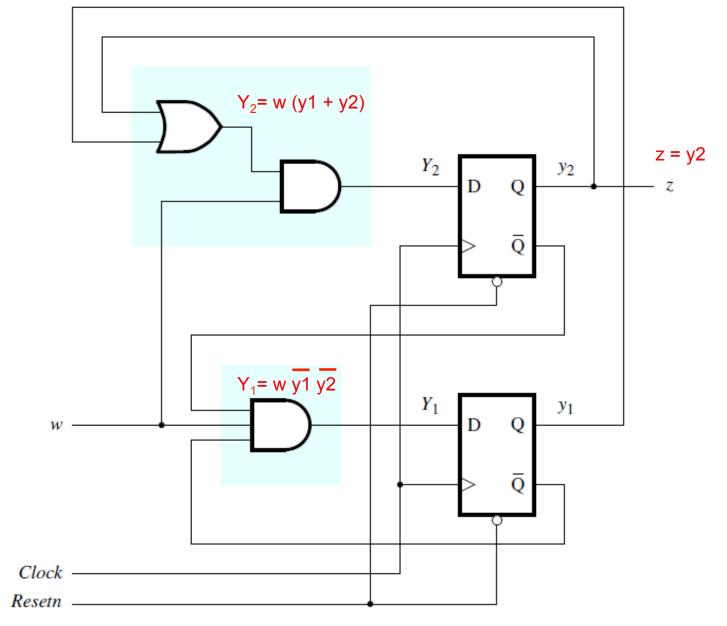
$$Y_2 = wy_1 + wy_2$$
$$= w(y_1 + y_2)$$

$$z = y_2$$

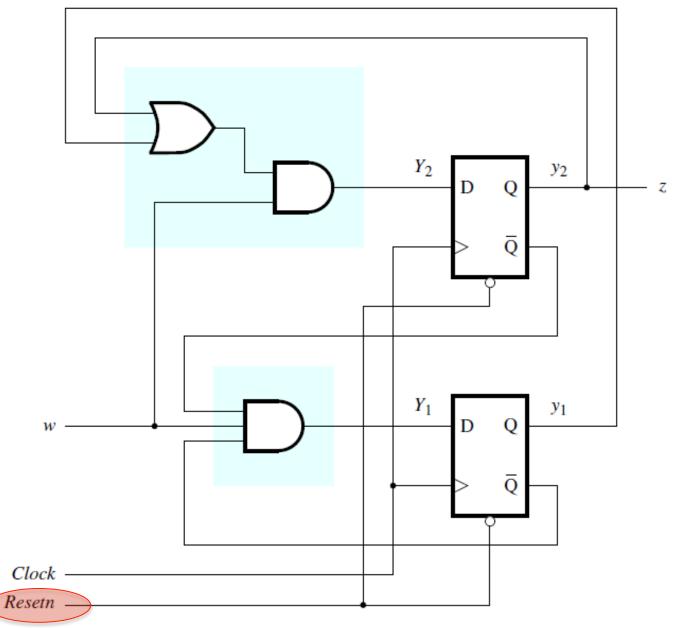
[ Figure 6.7 from the textbook ]



[ Figure 6.8 from the textbook ]

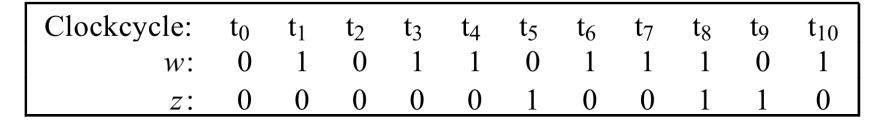


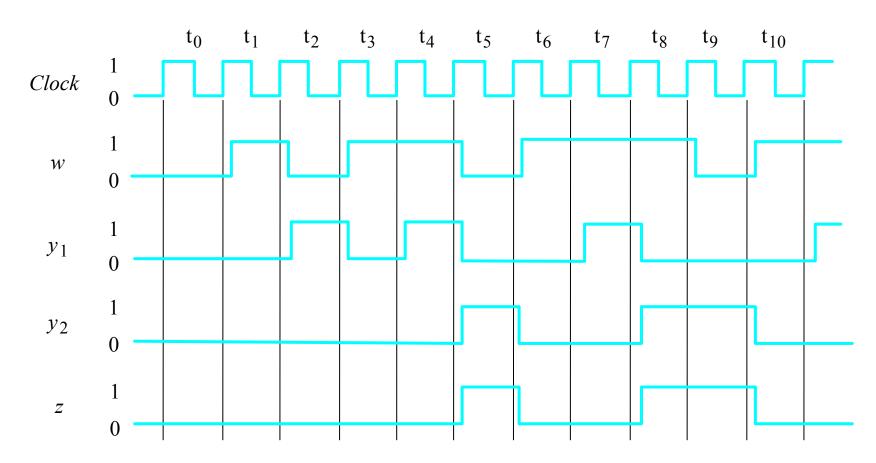
[ Figure 6.8 from the textbook ]



Lastly, we add a reset signal, which forces the machine back to its start state, which is state 00 in this case.

[ Figure 6.8 from the textbook ]





[ Figure 6.9 from the textbook ]

## Summary: Designing a Moore Machine

- Obtain the circuit specification.
- Derive a state diagram.
- Derive the state table.
- Decide on a state encoding.
- Encode the state table.
- Derive the output logic and next-state logic.
- Add a reset signal.

**Questions?** 

## THE END