

CprE 281: Digital Logic

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http://www.ece.iastate.edu/~alexs/classes/

Examples of Solved Problems

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Administrative Stuff

No homework is due next week.

Administrative Stuff

- Midterm Exam #1
- When: Friday Sep 23.
- Where: This classroom
- What: Chapter 1 and Chapter 2 plus number systems
- The exam will be open book and open notes (you can bring up to 3 pages of handwritten notes).

Topics for the Midterm Exam

- Binary Numbers
- Octal Numbers
- Hexadecimal Numbers
- Conversion between the different number systems
- Truth Tables
- Boolean Algebra
- Logic Gates
- Circuit Synthesis with AND, OR, NOT
- Circuit Synthesis with NAND, NOR
- Converting an AND/OR/NOT circuit to NAND circuit
- Converting an AND/OR/NOT circuit to NOR circuit
- SOP and POS expressions

Topics for the Midterm Exam

- Mapping a Circuit to Verilog code
- Mapping Verilog code to a circuit
- Multiplexers
- Venn Diagrams
- K-maps for 2, 3, and 4 variables
- Minimization of Boolean expressions using theorems
- Minimization of Boolean expressions with K-maps
- Incompletely specified functions (with don't cares)
- Functions with multiple outputs

Example 1

Determine if the following equation is valid

$$\bar{x}_1\bar{x}_3 + x_2x_3 + x_1\bar{x}_2 = \bar{x}_1x_2 + x_1x_3 + \bar{x}_2\bar{x}_3$$

$$\overline{x}_{1}\overline{x}_{3} + x_{2}x_{3} + x_{1}\overline{x}_{2} = \overline{x}_{1}x_{2} + x_{1}x_{3} + \overline{x}_{2}\overline{x}_{3}$$

$$\bar{x}_1 \bar{x}_3 + x_2 x_3 + x_1 \bar{x}_2 = \bar{x}_1 x_2 + x_1 x_3 + \bar{x}_2 \bar{x}_3$$

LHS RHS

Left-Hand Side (LHS)

| Row number | x_1 | x_2 | x_3 | $\frac{1}{x_1}$ | $x_2 x_3$ | $x_1 x_2$ | \int |
|--|------------------|------------------|---|-----------------|------------|-----------|--------|
| 0 1 2 3 | 0 0 0 0 | 0 0 1 1 | 0 1 0 1 | | | | |
| $egin{array}{c} 4 \ 5 \ 6 \ 7 \end{array}$ | 1 1 1 1 | 0 0 1 1 | $egin{array}{c} 0 \\ 1 \\ 0 \\ 1 \end{array}$ | | | | |

Left-Hand Side (LHS)

| Row number | x_1 | x_2 | x_3 | $\frac{-}{x_1}\frac{-}{x_3}$ | x_2x_3 | $x_1 \overline{x_2}$ | f |
|---------------|-------|-------|-------|------------------------------|----------|----------------------|---|
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | |
| 1 | 0 | 0 | 1 | 0 | 0 | 0 | |
| 2 | 0 | 1 | 0 | 1 | 0 | 0 | |
| 3 | 0 | 1 | 1 | 0 | 1 | 0 | |
| 4 | 1 | 0 | 0 | 0 | 0 | 1 | |
| 5 | 1 | 0 | 1 | 0 | 0 | 1 | |
| 6 | 1 | 1 | 0 | 0 | 0 | 0 | |
| 7 | 1 | 1 | 1 | 0 | 1 | $\overset{\circ}{0}$ | |

Left-Hand Side (LHS)

| Row number | x_1 | x_2 | x_3 | $\frac{1}{x_1}\frac{1}{x_3}$ | x_2x_3 | $x_1 \overline{x_2}$ | f |
|--|---|-------------|--|---|--|---|---|
| $egin{array}{c} 0 \ 1 \ 2 \end{array}$ | $\begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$ | 0 0 1 | 0 1 0 | 1 0 1 | 0 0 | 0 | 1 0 1 |
| $\frac{3}{4}$ | $\begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}$ | 1 0 | 1 0 | $\begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$ | 1 0 | $egin{array}{c} 0 \ 0 \ 1 \end{array}$ | $\begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$ |
| 5 6 7 | $egin{array}{c c} 1 \\ 1 \\ 1 \end{array}$ | 0 1 1 | $egin{array}{c} 1 \ 0 \ 1 \end{array}$ | $\begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$ | $egin{array}{c} 0 \ 0 \ 1 \end{array}$ | $\begin{matrix} 1 \\ 0 \\ 0 \end{matrix}$ | $\begin{array}{ c c }\hline 1\\0\\1\\\end{array}$ |

Right-Hand Side (RHS)

| Row number | x_1 | x_2 | x_3 | $\frac{-}{x_1x_2}$ | x_1x_3 | $-\frac{1}{x_2} - \frac{1}{x_3}$ | f |
|---------------|-------|-------|-------|--------------------|----------|----------------------------------|---|
| 0 | 0 | 0 | 0 | | | | |
| 1 | 0 | 0 | 1 | | | | |
| 2 | 0 | 1 | 0 | | | | |
| 3 | 0 | 1 | 1 | | | | |
| 4 | 1 | 0 | 0 | | | | |
| 5 | 1 | 0 | 1 | | | | |
| 6 | 1 | 1 | 0 | | | | |
| 7 | 1 | 1 | 1 | | | | |

Right-Hand Side (RHS)

| Row number | x_1 | x_2 | x_3 | $\frac{-}{x_1x_2}$ | x_1x_3 | $\frac{-}{x_2} \frac{-}{x_3}$ | f |
|---------------|-------|-------|-------|--------------------|----------|-------------------------------|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 1 | |
| 1 | 0 | 0 | 1 | 0 | 0 | 0 | |
| 2 | 0 | 1 | 0 | 1 | 0 | 0 | |
| 3 | 0 | 1 | 1 | 1 | 0 | 0 | |
| 4 | 1 | 0 | 0 | 0 | 0 | 1 | |
| 5 | 1 | 0 | 1 | 0 | 1 | 0 | |
| 6 | 1 | 1 | 0 | 0 | 0 | 0 | |
| 7 | 1 | 1 | 1 | 0 | 1 | Ö | |

Right-Hand Side (RHS)

| Row number | x_1 | x_2 | x_3 | $\frac{1}{x_1x_2}$ | x_1x_3 | $\frac{}{x_2}\frac{}{x_3}$ | f |
|---------------|-------|-------|-------|--------------------|----------|----------------------------|--|
| 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |
| 1 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 2 | 0 | 1 | 0 | 1 | 0 | 0 | 1 |
| 3 | 0 | 1 | 1 | 1 | 0 | 0 | 1 |
| 4 | 1 | 0 | 0 | 0 | 0 | 1 | 1 |
| 5 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |
| 6 | 1 | 1 | 0 | 0 | 0 | 0 | |
| 7 | 1 | 1 | 1 | 0 | 1 | $\overset{\circ}{0}$ | $\begin{bmatrix} & \cdot \\ & 1 \end{bmatrix}$ |

$$\overline{x}_1\overline{x}_3 + x_2x_3 + x_1\overline{x}_2 \stackrel{?}{=} \overline{x}_1x_2 + x_1x_3 + \overline{x}_2\overline{x}_3$$
LHS
RHS

| f | f |
|---|---|
| 1 | 1 |
| 0 | 0 |
| 1 | 1 |
| 1 | 1 |
| 1 | 1 |
| 1 | 1 |
| 0 | 0 |
| 1 | 1 |

Example 2

Design the minimum-cost product-of-sums expression for the function

$$f(x_1, x_2, x_3) = \sum m(0, 2, 4, 5, 6, 7)$$

Minterms and Maxterms (with three variables)

| Row number | x_1 | x_2 | x_3 | Minterm | Maxterm |
|--|---------------------------------|--------------------------------------|---------------------------------|---|---|
| $egin{array}{c} 0 \\ 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \\ 7 \\ \end{array}$ | 0 0 0 0 1 1 1 | 0 0 1 1 0 0 1 1 | 0 1 0 1 0 1 0 | $m_0 = \overline{x}_1 \overline{x}_2 \overline{x}_3$ $m_1 = \overline{x}_1 \overline{x}_2 x_3$ $m_2 = \overline{x}_1 x_2 \overline{x}_3$ $m_3 = \overline{x}_1 x_2 x_3$ $m_4 = x_1 \overline{x}_2 \overline{x}_3$ $m_5 = x_1 \overline{x}_2 x_3$ $m_6 = x_1 x_2 \overline{x}_3$ $m_7 = x_1 x_2 x_3$ | $M_0 = x_1 + x_2 + x_3$ $M_1 = x_1 + x_2 + \overline{x_3}$ $M_2 = x_1 + \overline{x_2} + x_3$ $M_3 = x_1 + \overline{x_2} + \overline{x_3}$ $M_4 = \overline{x_1} + x_2 + x_3$ $M_5 = \overline{x_1} + x_2 + \overline{x_3}$ $M_6 = \overline{x_1} + \overline{x_2} + x_3$ $M_7 = \overline{x_1} + \overline{x_2} + \overline{x_3}$ |

Minterms and Maxterms (with three variables)

| Row number | $ x_1 $ | x_2 | x_3 | Minterm | Maxterm |
|---------------|-----------|-------|-------|--|--|
| 0 | 0 | 0 | 0 | $m_0 = \overline{x}_1 \overline{x}_2 \overline{x}_3$ | $M_0 = x_1 + x_2 + x_3$ |
| 1 | 0 | 0 | 1 | $m_1 = \overline{x}_1 \overline{x}_2 x_3$ | $M_1 = x_1 + x_2 + \overline{x}_3$ |
| 2 | 0 | 1 | 0 | $m_2 = \overline{x}_1 x_2 \overline{x}_3$ | $M_2 = x_1 + \overline{x}_2 + x_3$ |
| 3 | 0 | 1 | 1 | $m_3 = \overline{x}_1 x_2 x_3$ | $M_3 = x_1 + \overline{x}_2 + \overline{x}_3$ |
| 4 | 1 | 0 | 0 | $m_4 = x_1 \overline{x}_2 \overline{x}_3$ | $M_4 = \overline{x}_1 + x_2 + x_3$ |
| 5 | 1 | 0 | 1 | $m_5 = x_1 \overline{x}_2 x_3$ | $M_5 = \overline{x}_1 + x_2 + \overline{x}_3$ |
| 6 | 1 | 1 | 0 | $m_6 = x_1 x_2 \overline{x}_3$ | $M_6 = \overline{x}_1 + \overline{x}_2 + x_3$ |
| 7 | 1 | 1 | 1 | $m_7 = x_1 x_2 x_3$ | $M_7 = \overline{x}_1 + \overline{x}_2 + \overline{x}_3$ |

The function is 1 for these rows

Minterms and Maxterms (with three variables)

| Row number | x_1 | x_2 | x_3 | Minterm | Maxterm |
|--|----------------------------|----------------------------|----------------------------|---|--|
| $egin{array}{c} 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ \end{array}$ | 0 0 0 0 1 1 | 0 0 1 1 0 0 | 0 1 0 1 0 1 | $m_0 = \overline{x}_1 \overline{x}_2 \overline{x}_3$ $m_1 = \overline{x}_1 \overline{x}_2 x_3$ $m_2 = \overline{x}_1 x_2 \overline{x}_3$ $m_3 = \overline{x}_1 x_2 x_3$ $m_4 = x_1 \overline{x}_2 \overline{x}_3$ $m_5 = x_1 \overline{x}_2 x_3$ $m_6 = x_1 x_2 \overline{x}_3$ | $M_0 = x_1 + x_2 + x_3$ $M_1 = x_1 + x_2 + \overline{x}_3$ $M_2 = x_1 + \overline{x}_2 + x_3$ $M_3 = x_1 + \overline{x}_2 + \overline{x}_3$ $M_4 = \overline{x}_1 + x_2 + x_3$ $M_5 = \overline{x}_1 + x_2 + \overline{x}_3$ $M_6 = \overline{x}_1 + \overline{x}_2 + x_3$ |
| 7 | 1 | 1 | 1 | $m_7 = x_1 x_2 x_3$ | $M_7 = \overline{x}_1 + \overline{x}_2 + \overline{x}_3$ |

The function is 1 for these rows

The function is 0 for these rows

Two different ways to specify the same function f of three variables

$$f(x_1, x_2, x_3) = \sum m(0, 2, 4, 5, 6, 7)$$

$$f(x_1, x_2, x_3) = \prod M(1, 3)$$

The POS Expression

$$M_1 = x_1 + x_2 + \overline{x}_3$$

$$M_3 = x_1 + \overline{x}_2 + \overline{x}_3$$

$$f(x_1, x_2, x_3) = \Pi M(1, 3)$$

$$= M_1 \cdot M_3$$

$$= (x_1 + x_2 + \overline{x_3}) \cdot (x_1 + \overline{x_2} + \overline{x_3})$$

The Minimum POS Expression

$$f(x_1, x_2, x_3) = (x_1 + x_2 + \overline{x_3}) \cdot (x_1 + \overline{x_2} + \overline{x_3})$$

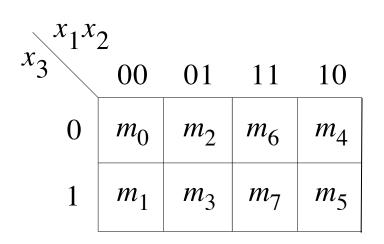
$$= (x_1 + \overline{x_3} + x_2) \cdot (x_1 + \overline{x_3} + \overline{x_2})$$

$$= (x_1 + \overline{x_3})$$

Hint: Use the following Boolean Algebra theorem

14b.
$$(x + y) \cdot (x + \overline{y}) = x$$

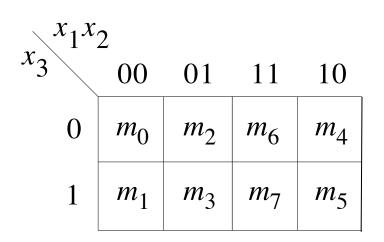
| x_1 | x_2 | x_3 | |
|-------|-------|-------|-------|
| 0 | 0 | 0 | m_0 |
| 0 | 0 | 1 | m_1 |
| 0 | 1 | 0 | m_2 |
| 0 | 1 | 1 | m_3 |
| 1 | 0 | 0 | m_4 |
| 1 | 0 | 1 | m_5 |
| 1 | 1 | 0 | m_6 |
| 1 | 1 | 1 | m_7 |
| | | | Į. |



(b) Karnaugh map

(a) Truth table

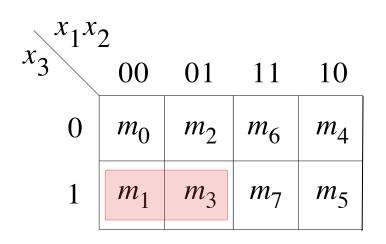
| x_1 | x_2 | x_3 | |
|-------|-------|-------|-------|
| 0 | 0 | 0 | m_0 |
| 0 | 0 | 1 | m_1 |
| 0 | 1 | 0 | m_2 |
| 0 | 1 | 1 | m_3 |
| 1 | 0 | 0 | m_4 |
| 1 | 0 | 1 | m_5 |
| 1 | 1 | 0 | m_6 |
| 1 | 1 | 1 | m_7 |
| | | | Į. |



(b) Karnaugh map

(a) Truth table

| x_1 | x_2 | x_3 | |
|-------|-------|-------|-------|
| 0 | 0 | 0 | m_0 |
| 0 | 0 | 1 | m_1 |
| 0 | 1 | 0 | m_2 |
| 0 | 1 | 1 | m_3 |
| 1 | 0 | 0 | m_4 |
| 1 | 0 | 1 | m_5 |
| 1 | 1 | 0 | m_6 |
| 1 | 1 | 1 | m_7 |
| | | | I |

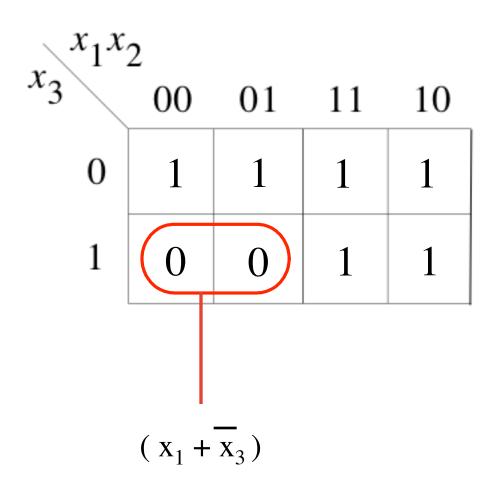


(b) Karnaugh map

(a) Truth table

| $x_1^{x_1}$ | | | | | | |
|-------------|-------|-------|-------|-------|--|--|
| *3 | 00 | 01 | 11 | 10 | | |
| 0 | m_0 | m_2 | m_6 | m_4 | | |
| 1 | m_1 | m_3 | m_7 | m_5 | | |

| $x_1^{x_1}$ | | | | | | |
|-------------|----|----|----|----|--|--|
| *3 | 00 | 01 | 11 | 10 | | |
| 0 | 1 | 1 | 1 | 1 | | |
| 1 | 0 | 0 | 1 | 1 | | |



Example 3

Problem: A circuit that controls a given digital system has three inputs: x_1 , x_2 , and x_3 . It has to recognize three different conditions:

- Condition A is true if x_3 is true and either x_1 is true or x_2 is false
- Condition B is true if x_1 is true and either x_2 or x_3 is false
- Condition C is true if x_2 is true and either x_1 is true or x_3 is false

The control circuit must produce an output of 1 if at least two of the conditions A, B, and C are true. Design the simplest circuit that can be used for this purpose.

Condition A

Condition A is true if x_3 is true and either x_1 is true or x_2 is false

Condition A

Condition A is true if x_3 is true and either x_1 is true or x_2 is false

$$A = x_3(x_1 + \overline{x}_2) = x_3x_1 + x_3\overline{x}_2$$

Condition B

Condition B is true if x_1 is true and either x_2 or x_3 is false

Condition B

Condition B is true if x_1 is true and either x_2 or x_3 is false

$$B = x_1(\bar{x}_2 + \bar{x}_3) = x_1\bar{x}_2 + x_1\bar{x}_3$$

Condition C

Condition C is true if x_2 is true and either x_1 is true or x_3 is false

Condition C

Condition C is true if x_2 is true and either x_1 is true or x_3 is false

$$C = x_2(x_1 + \overline{x}_3) = x_2x_1 + x_2\overline{x}_3$$

The output of the circuit can be expressed as f = AB + AC + BC

$$AB = (x_3x_1 + x_3\bar{x}_2)(x_1\bar{x}_2 + x_1\bar{x}_3)$$

$$= x_3x_1x_1\bar{x}_2 + x_3x_1x_1\bar{x}_3 + x_3\bar{x}_2x_1\bar{x}_2 + x_3\bar{x}_2x_1\bar{x}_3$$

$$= x_3x_1\bar{x}_2 + 0 + x_3\bar{x}_2x_1 + 0$$

$$= x_1\bar{x}_2x_3$$

The output of the circuit can be expressed as f = AB + AC + BC

$$AC = (x_3x_1 + x_3\overline{x}_2)(x_2x_1 + x_2\overline{x}_3)$$

$$= x_3x_1x_2x_1 + x_3x_1x_2\overline{x}_3 + x_3\overline{x}_2x_2x_1 + x_3\overline{x}_2x_2\overline{x}_3$$

$$= x_3x_1x_2 + 0 + 0 + 0$$

$$= x_1x_2x_3$$

The output of the circuit can be expressed as f = AB + AC + BC

$$BC = (x_1 \bar{x}_2 + x_1 \bar{x}_3)(x_2 x_1 + x_2 \bar{x}_3)$$

$$= x_1 \bar{x}_2 x_2 x_1 + x_1 \bar{x}_2 x_2 \bar{x}_3 + x_1 \bar{x}_3 x_2 x_1 + x_1 \bar{x}_3 x_2 \bar{x}_3$$

$$= 0 + 0 + x_1 \bar{x}_3 x_2 + x_1 \bar{x}_3 x_2$$

$$= x_1 x_2 \bar{x}_3$$

Finally, we get

$$f = x_1 \overline{x}_2 x_3 + x_1 x_2 x_3 + x_1 x_2 \overline{x}_3$$

$$= x_1 (\overline{x}_2 + x_2) x_3 + x_1 x_2 (x_3 + \overline{x}_3)$$

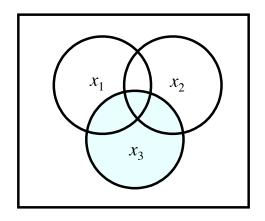
$$= x_1 x_3 + x_1 x_2$$

$$= x_1 (x_3 + x_2)$$

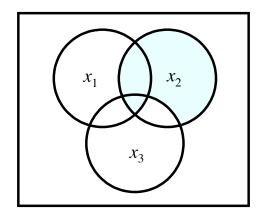
Example 4

Solve the previous problem using Venn diagrams.

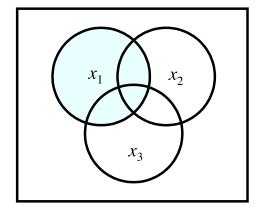
Venn Diagrams (find the areas that are shaded at least two times)



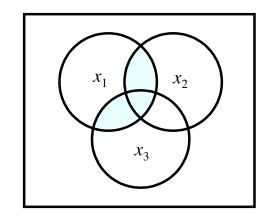
(a) Function A



(c) Function C



(b) Function B

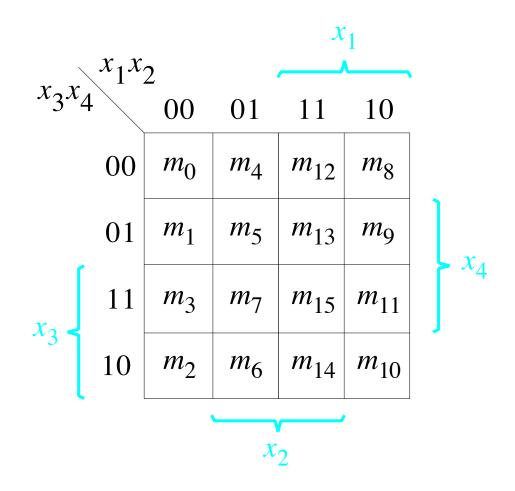


(d) Function f

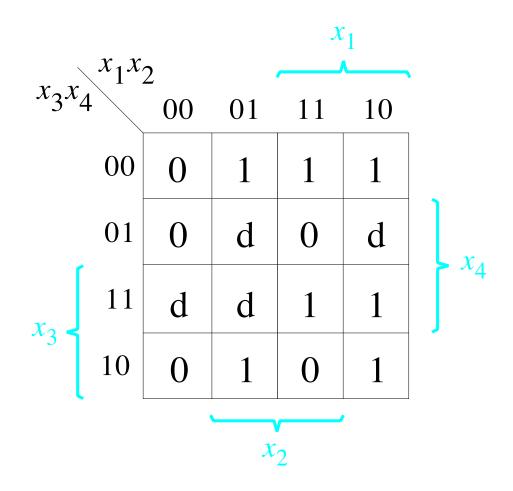
Example 5

Design the minimum-cost SOP and POS expression for the function

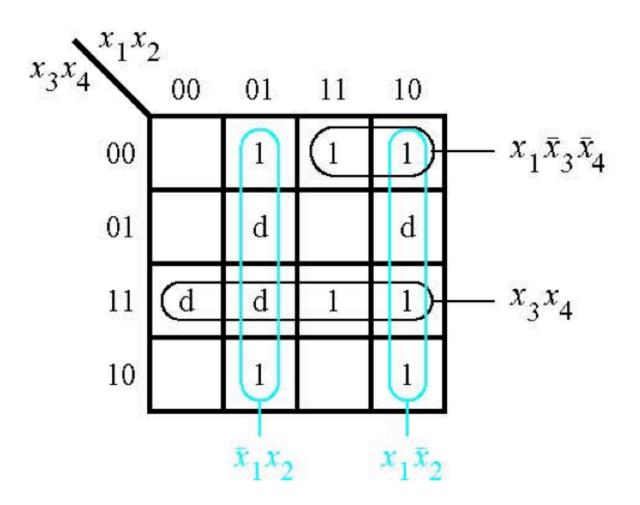
Let's Use a K-Map



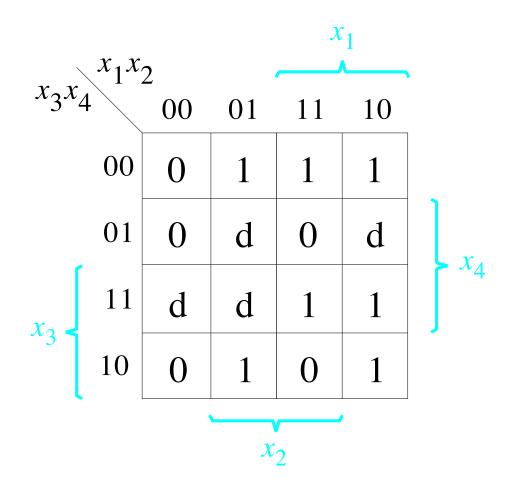
Let's Use a K-Map



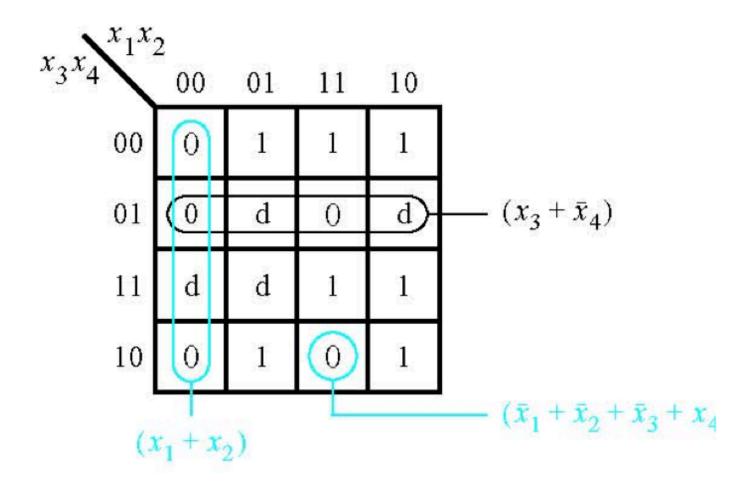
The SOP Expression



What about the POS Expression?



The POS Expression



Example 6

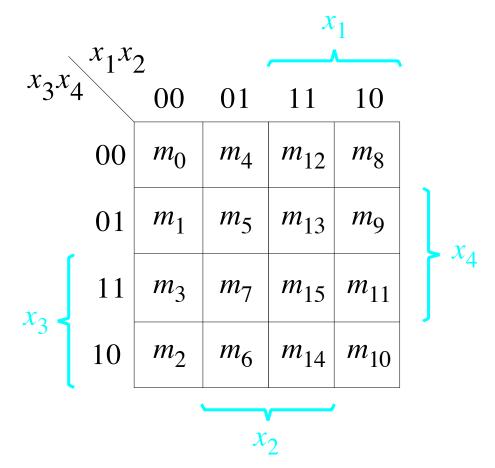
Use K-maps to find the minimum-cost SOP and POS expression for the function

$$f(x_1, \ldots, x_4) = \overline{x}_1 \overline{x}_3 \overline{x}_4 + x_3 x_4 + \overline{x}_1 \overline{x}_2 x_4 + x_1 x_2 \overline{x}_3 x_4$$

assuming that there are also don't-cares defined as $D = \sum (9, 12, 14)$.

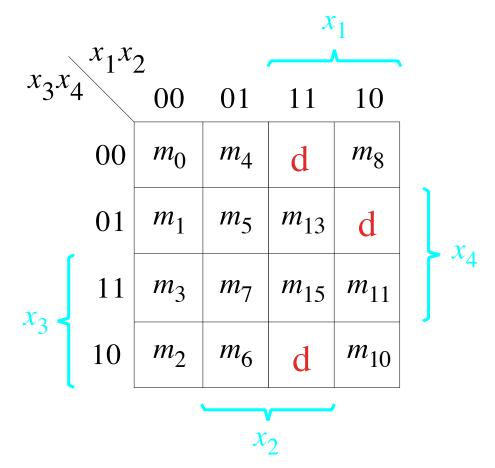
Let's map the expression to the K-Map

$$f(x_1, \dots, x_4) = \overline{x}_1 \overline{x}_3 \overline{x}_4 + x_3 x_4 + \overline{x}_1 \overline{x}_2 x_4 + x_1 x_2 \overline{x}_3 x_4$$
$$D = \sum (9, 12, 14).$$



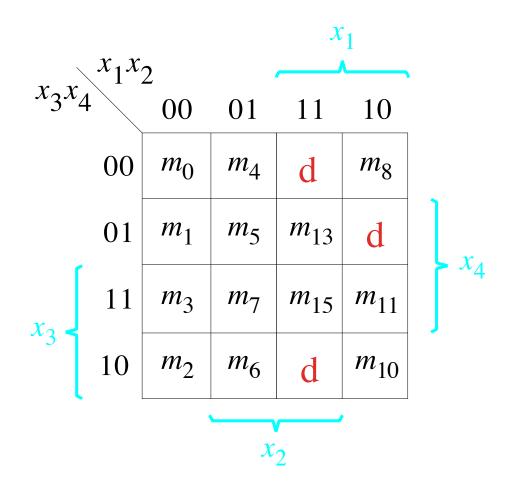
Let's map the expression to the K-Map

$$f(x_1, \dots, x_4) = \overline{x}_1 \overline{x}_3 \overline{x}_4 + x_3 x_4 + \overline{x}_1 \overline{x}_2 x_4 + x_1 x_2 \overline{x}_3 x_4$$
$$D = \sum (9, 12, 14).$$



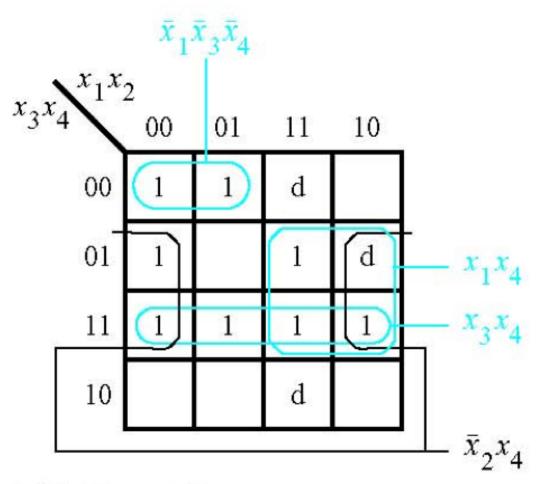
Let's map the expression to the K-Map

$$f(x_1, \ldots, x_4) = \overline{x}_1 \overline{x}_3 \overline{x}_4 + x_3 x_4 + \overline{x}_1 \overline{x}_2 x_4 + x_1 x_2 \overline{x}_3 x_4$$



The SOP Expression

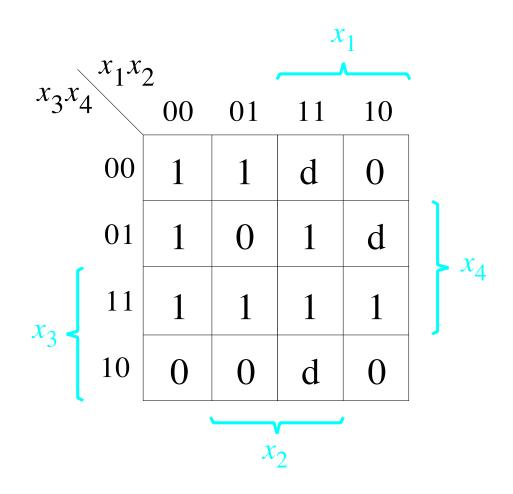
$$f(x_1, \ldots, x_4) = \overline{x}_1 \overline{x}_3 \overline{x}_4 + x_3 x_4 + \overline{x}_1 \overline{x}_2 x_4 + x_1 x_2 \overline{x}_3 x_4$$



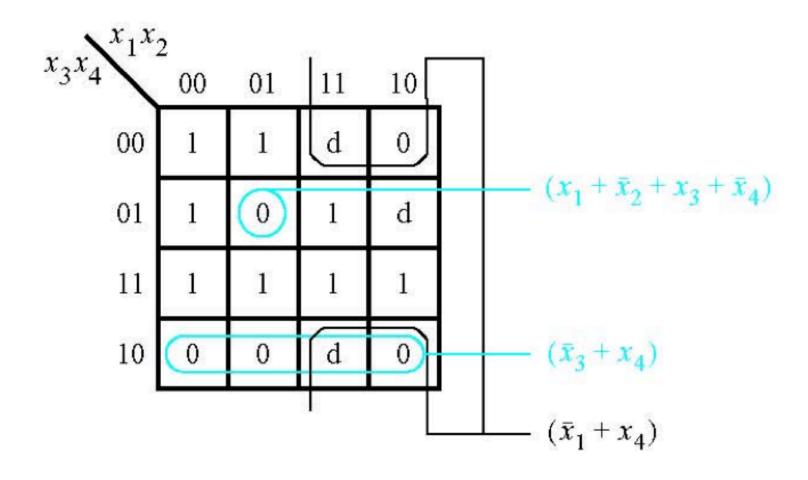
$$f = x_3 x_4 + \bar{x}_1 \bar{x}_3 \bar{x}_4 + \bar{x}_2 x_4 + x_1 x_4$$

[Figure 2.68a from the textbook]

What about the POS Expression?



The POS Expression



$$f = (\bar{x}_3 + x_4)(\bar{x}_1 + x_4)(x_1 + \bar{x}_2 + x_3 + \bar{x}_4)$$

Example 7

Derive the minimum-cost SOP expression for

$$f = s_3(\overline{s}_1 + \overline{s}_2) + s_1s_2$$

First, expand the expression using property 12a

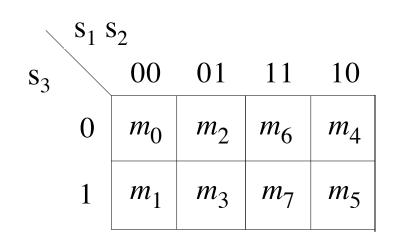
$$f = s_3(\bar{s}_1 + \bar{s}_2) + s_1s_2$$

$$f = \bar{s}_1 s_3 + \bar{s}_2 s_3 + s_1 s_2$$

Construct the K-Map for this expression

$$f = \bar{s}_1 s_3 + \bar{s}_2 s_3 + s_1 s_2$$

| s_1 | S_2 | s_3 | |
|-------|-------|-------|-------|
| 0 | 0 | 0 | m_0 |
| 0 | 0 | 1 | m_1 |
| 0 | 1 | 0 | m_2 |
| 0 | 1 | 1 | m_3 |
| 1 | 0 | 0 | m_4 |
| 1 | 0 | 1 | m_5 |
| 1 | 1 | 0 | m_6 |
| 1 | 1 | 1 | m_7 |
| | | | |

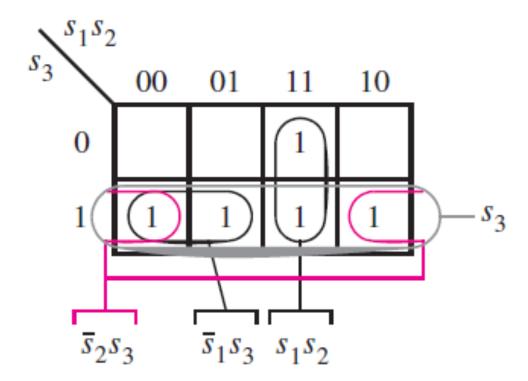


(b) Karnaugh map

(a) Truth table

Construct the K-Map for this expression

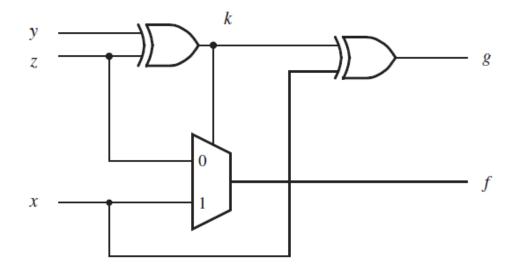
$$f = \bar{s}_1 s_3 + \bar{s}_2 s_3 + s_1 s_2$$



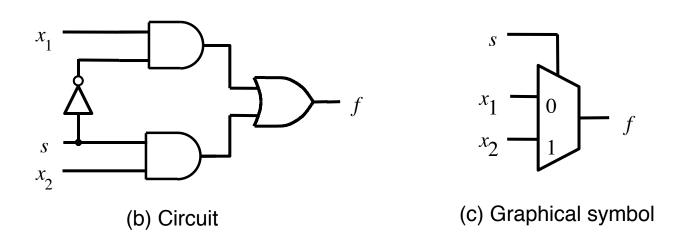
Example 8

Write the Verilog code for the following circuit ...

Logic Circuit

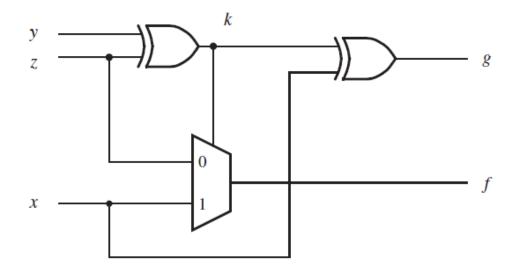


Circuit for 2-1 Multiplexer



$$f(s, x_1, x_2) = \overline{s} x_1 + s x_2$$

Logic Circuit vs Verilog Code



```
module f_g(x, y, z, f, g);

input x, y, z;

output f, g;

wire k;

assign k = y \ z;

assign g = k \ x;

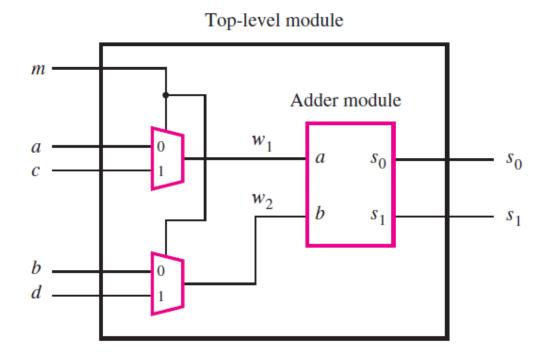
assign f = (\sim k \& z) \mid (k \& x);
```

endmodule

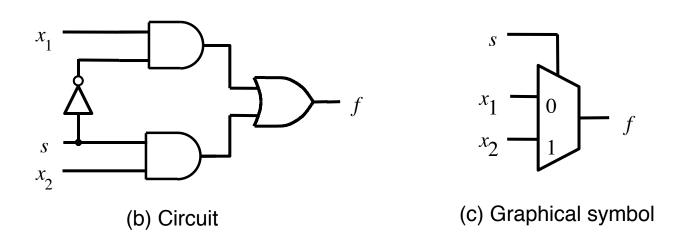
Example 9

Write the Verilog code for the following circuit ...

The Logic Circuit for this Example

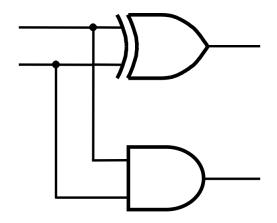


Circuit for 2-1 Multiplexer



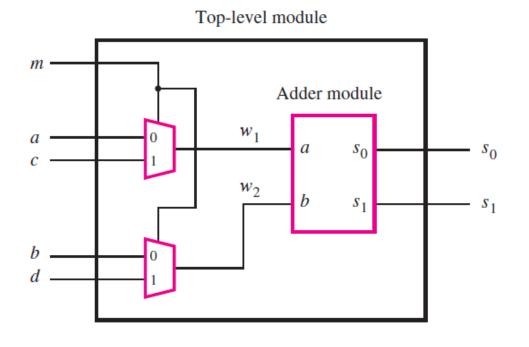
$$f(s, x_1, x_2) = \overline{s} x_1 + s x_2$$

Addition of Binary Numbers



| <i>s</i> ₁ | s_0 |
|-----------------------|-------------|
| 0 | 0 |
| 0 | 1 |
| 0 | 1 |
| 1 | 0 |
| | 0 0 0 |

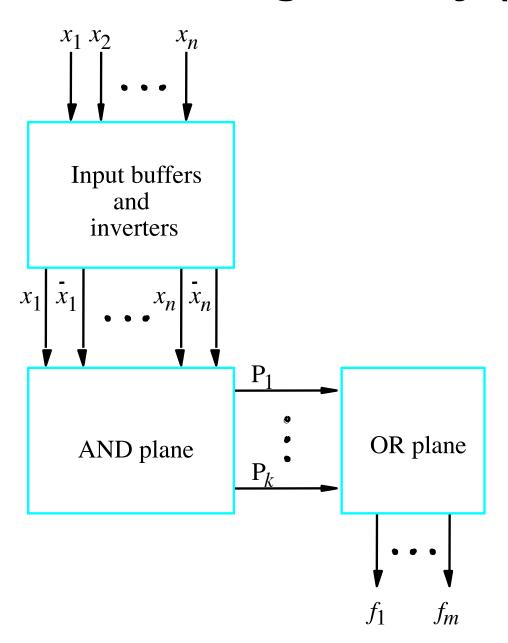
Logic Circuit vs Verilog Code



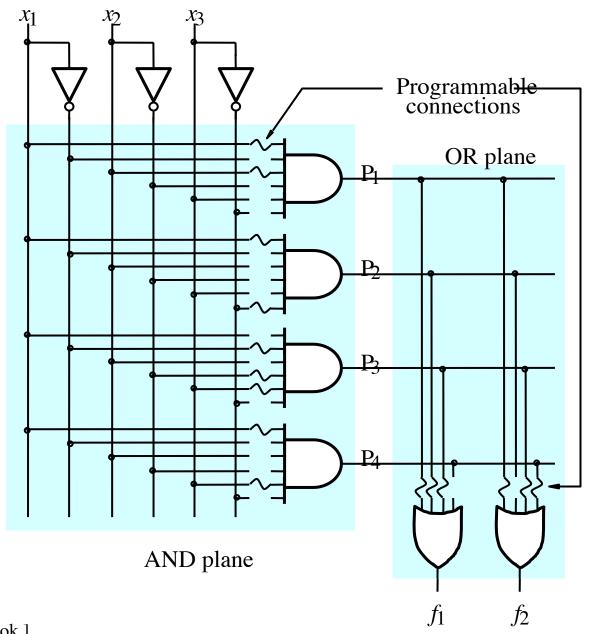
```
module shared (a, b, c, d, m, s1, s0);
   input a, b, c, d, m;
  output s1, s0;
  wire w1, w2;
  mux2to1 U1 (a, c, m, w1);
  mux2to1 U2 (b, d, m, w2);
   adder U3 (w1, w2, s1, s0);
endmodule
module mux2to1 (x1, x2, s, f);
   input x1, x2, s;
  output f;
   assign f = (\sim s \& x1) | (s \& x2);
endmodule
module adder (a, b, s1, s0);
  input a, b;
  output s1, s0;
   assign s1 = a \& b;
  assign s0 = a ^ b;
endmodule
```

Some material form Appendix B

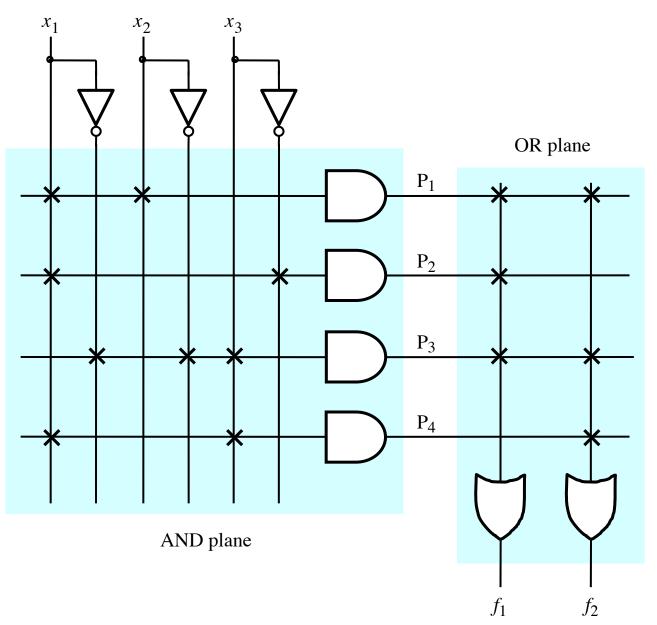
Programmable Logic Array (PLA)



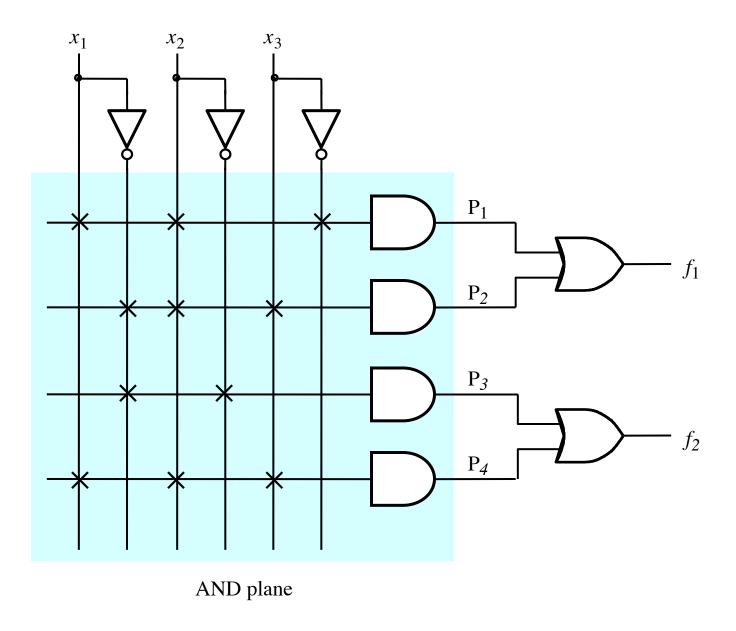
Gate-Level Diagram of a PLA



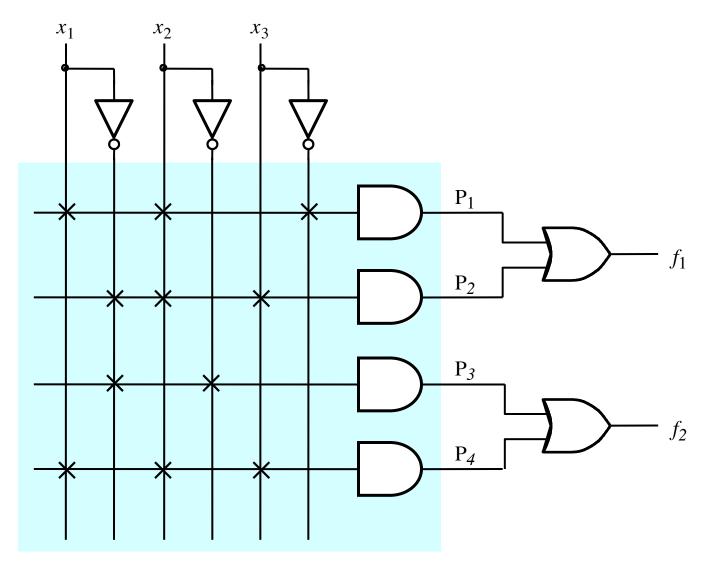
Customary Schematic for PLA



Programmable Array Logic (PAL)



Programmable Array Logic (PAL)



AND plane

Only the AND plane is programmable. The OR plane is fixed.

Questions?

THE END