

## CprE 281: Digital Logic

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## Multiplexers

CprE 281: Digital Logic
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## Administrative Stuff

- HW 6 is due on Monday


## Administrative Stuff

- HW 7 is out
- It is due on Monday Oct 17 @ 4pm


## 2-1 Multiplexer (Definition)

- Has two inputs: $x_{1}$ and $x_{2}$
- Also has another input line s
- If $\mathbf{s}=0$, then the output is equal to $\mathbf{x}_{1}$
- If $\mathbf{s}=1$, then the output is equal to $\mathbf{x}_{2}$


## Graphical Symbol for a 2-1 Multiplexer



## Truth Table for a 2-1 Multiplexer

| $s$ | $x_{1}$ | $x_{2}$ | $f\left(s, x_{1}, x_{2}\right)$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

[ Figure 2.33a from the textbook]

## Let's Derive the SOP form

| $s$ | $x_{1}$ | $x_{2}$ | $f\left(s, x_{1}, x_{2}\right)$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

## Let's Derive the SOP form

| $s x_{1} x_{2}$ | $f\left(s, x_{1}, x_{2}\right)$ |
| :---: | :---: |
| 000 | 0 |
| 001 | 0 |
| 010 | 1 |
| 011 | 1 |
| 100 | 0 |
| 101 | 1 |
| 110 | 0 |
| 111 | 1 |

## Let's Derive the SOP form

| $s$ | $x_{1}$ | $x_{2}$ |
| :---: | :---: | :---: | \left\lvert\, \(f\left(s, x_{1}, x_{2}\right) ~\left(\left.\begin{array}{ccc}0 \& 0 <br>

\hline 0 \& 0 \& 0\end{array} \right\rvert\,\right.\right.\)

Where should we put the negation signs?

$$
\begin{array}{lll}
s & x_{1} & x_{2} \\
s & x_{1} & x_{2} \\
& & x_{1} \\
x_{2}
\end{array}
$$

## Let's Derive the SOP form



## Let's Derive the SOP form



## Let's simplify this expression

$$
f\left(s, x_{1}, x_{2}\right)=\bar{s} x_{1} \bar{x}_{2}+\bar{s} x_{1} x_{2}+s \bar{x}_{1} x_{2}+s x_{1} x_{2}
$$

## Let's simplify this expression

$$
\begin{aligned}
& f\left(s, x_{1}, x_{2}\right)=\bar{s} x_{1} \bar{x}_{2}+\bar{s} x_{1} x_{2}+s \bar{x}_{1} x_{2}+s x_{1} x_{2} \\
& f\left(s, x_{1}, x_{2}\right)=\bar{s} x_{1}\left(\bar{x}_{2}+x_{2}\right)+s\left(\bar{x}_{1}+x_{1}\right) x_{2}
\end{aligned}
$$

## Let's simplify this expression

$$
\begin{aligned}
& f\left(s, x_{1}, x_{2}\right)=\bar{s} x_{1} \bar{x}_{2}+\bar{s} x_{1} x_{2}+s \bar{x}_{1} x_{2}+s x_{1} x_{2} \\
& f\left(s, x_{1}, x_{2}\right)=\bar{s} x_{1}\left(\bar{x}_{2}+x_{2}\right)+s\left(\bar{x}_{1}+x_{1}\right) x_{2} \\
& f\left(s, x_{1}, x_{2}\right)=\bar{s} x_{1}+s x_{2}
\end{aligned}
$$

## Circuit for 2-1 Multiplexer


(b) Circuit

(c) Graphical symbol

$$
f\left(s, x_{1}, x_{2}\right)=\bar{s} x_{1}+s x_{2}
$$

[ Figure 2.33b-c from the textbook ]

## Analysis of the 2-1 Multiplexer (when the input $\mathbf{s}=0$ )



## Analysis of the 2-1 Multiplexer (when the input $\mathrm{s}=1$ )



## Analysis of the 2-1 Multiplexer (when the input $\mathrm{s}=0$ )



## Analysis of the 2-1 Multiplexer (when the input $\mathrm{s}=1$ )



## More Compact Truth-Table Representation

$\left.\begin{array}{cc||c}\hline s & x_{1} & x_{2} \\ \hline 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 1 & \left.x_{1}, x_{2}\right) \\ 1 & 0 & 0 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 1 & 1\end{array}\right]$

(a)Truth table
[ Figure 2.33 from the textbook ]

## 4-1 Multiplexer (Definition)

- Has four inputs: $\mathbf{w}_{0}, w_{1}, w_{2}, w_{3}$
- Also has two select lines: $\mathbf{s}_{1}$ and $\mathbf{s}_{0}$
- If $s_{1}=0$ and $s_{0}=0$, then the output $f$ is equal to $w_{0}$
- If $s_{1}=0$ and $s_{0}=1$, then the output $f$ is equal to $w_{1}$
- If $s_{1}=1$ and $s_{0}=0$, then the output $f$ is equal to $w_{2}$
- If $s_{1}=1$ and $s_{0}=1$, then the output $f$ is equal to $w_{3}$


## Graphical Symbol and Truth Table


(a) Graphic symbol

(b) Truth table

## The long-form truth table

## The long-form truth table



## 4-1 Multiplexer (SOP circuit)


[ Figure 4.2c from the textbook ]

## Analysis of the 4-1 Multiplexer ( $s_{1}=0$ and $s_{0}=0$ )



## Analysis of the 4-1 Multiplexer ( $s_{1}=0$ and $s_{0}=0$ )



## Analysis of the 4-1 Multiplexer ( $s_{1}=0$ and $s_{0}=0$ )



## Analysis of the 4-1 Multiplexer ( $s_{1}=0$ and $s_{0}=0$ )



## Analysis of the 4-1 Multiplexer ( $s_{1}=0$ and $s_{0}=1$ )



Analysis of the 4-1 Multiplexer ( $s_{1}=1$ and $s_{0}=0$ )


Analysis of the 4-1 Multiplexer
( $s_{1}=1$ and $s_{0}=1$ )


## Analysis of the 4-1 Multiplexer ( $s_{1}=0$ and $s_{0}=0$ )



## Analysis of the 4-1 Multiplexer ( $s_{1}=0$ and $s_{0}=1$ )



## Analysis of the 4-1 Multiplexer ( $s_{1}=1$ and $s_{0}=0$ )



Analysis of the 4-1 Multiplexer ( $s_{1}=1$ and $s_{0}=1$ )


## Using three 2-to-1 multiplexers to build one 4-to-1 multiplexer


[ Figure 4.3 from the textbook]

## Using three 2-to-1 multiplexers to build one 4-to-1 multiplexer



## Using three 2-to-1 multiplexers to build one 4-to-1 multiplexer



## Using three 2-to-1 multiplexers to build one 4-to-1 multiplexer



## Using three 2-to-1 multiplexers to build one 4-to-1 multiplexer



## That is different from the SOP form of the 4-1 multiplexer shown below, which uses less gates



## Analysis of the Hierarchical Implementation ( $\mathrm{s}_{1}=0$ and $\mathrm{s}_{0}=0$ )


[ Figure 4.3 from the textbook ]

## Analysis of the Hierarchical Implementation ( $s_{1}=0$ and $s_{0}=1$ )


[ Figure 4.3 from the textbook ]

Analysis of the Hierarchical Implementation ( $\mathrm{s}_{1}=1$ and $\mathrm{s}_{0}=0$ )

[ Figure 4.3 from the textbook ]

Analysis of the Hierarchical Implementation ( $\mathrm{s}_{1}=1$ and $\mathrm{s}_{0}=1$ )

[ Figure 4.3 from the textbook ]

## 16-1 Multiplexer


[ Figure 4.4 from the textbook ]

## Multiplexers Are Special

## The Three Basic Logic Gates



NOT gate


AND gate


OR gate

## Truth Table for NOT



## Truth Table for AND



| $x_{1}$ | $x_{2}$ | $x_{1} \cdot x_{2}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

## Truth Table for OR



| $x_{1}$ | $x_{2}$ | $x_{1}+x_{2}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

## Building an AND Gate with 4-to-1 Mux



| $x_{1}$ | $x_{2}$ | $x_{1} \cdot x_{2}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |



## Building an AND Gate with 4-to-1 Mux



These two are the same.

## Building an AND Gate with 4-to-1 Mux



These two are the same.
And so are these two.

## Building an OR Gate with 4-to-1 Mux

$$
\begin{array}{cc||c}
x_{1} & & x_{1}+x_{2} \\
x_{2} & & \\
x_{1} & x_{2} & x_{1}+x_{2} \\
\hline 0 & 0 & 0 \\
0 & 1 & 1 \\
1 & 0 & 1 \\
1 & 1 & 1
\end{array}
$$

## Building an OR Gate with 4-to-1 Mux



These two are the same.

## Building an OR Gate with 4-to-1 Mux



These two are the same.
And so are these two.

## Building a NOT Gate with 4-to-1 Mux



## Building a NOT Gate with 4-to-1 Mux



Introduce a dummy variable y .

## Building a NOT Gate with 4-to-1 Mux



## Building a NOT Gate with 4-to-1 Mux



Now set y to either 0 or 1 (both will work). Why?

## Building a NOT Gate with 4-to-1 Mux



| $x$ | $\bar{x}$ |
| :---: | :---: |
| 0 | 1 |
| 1 | 0 |

Two alternative solutions.

## Implications

# Any Boolean function can be implemented using only 4-to-1 multiplexers! 

## Building an AND Gate with 2-to-1 Mux



| $x_{1}$ | $x_{2}$ | $x_{1} \cdot x_{2}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |



## Building an AND Gate with 2-to-1 Mux



| $x_{1}$ | $x_{2}$ | $x_{1} \cdot x_{2}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |



## Building an AND Gate with 2-to-1 Mux


$\left.\begin{array}{c|c||cc}x_{1} & x_{2} & x_{1} \cdot x_{2} & \\ \hline 0 & 0 & 0 \\ 0 & 1 & 0\end{array}\right\} \quad 0$


## Building an OR Gate with 2-to-1 Mux

| $x_{1}$ | $x_{1}+x_{2}$ |  |
| :---: | :---: | :---: |
| $x_{2}$ | $x_{2}$ | $x_{1}+x_{2}$ |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |



## Building an OR Gate with 2-to-1 Mux

| $x_{1}$ |  |  |
| :---: | :---: | :---: |
| $x_{2}$ | $x_{1}+x_{2}$ |  |
| $x_{1}$ | $x_{2}$ | $x_{1}+x_{2}$ |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |



## Building an OR Gate with 2-to-1 Mux



## Building a NOT Gate with 2-to-1 Mux



## Building a NOT Gate with 2-to-1 Mux



| $x$ | $\bar{x}$ |
| :---: | :---: |
| 0 | 1 |
| 1 | 0 |



## Implications

# Any Boolean function can be implemented using only 2-to-1 multiplexers! 

## Synthesis of Logic Circuits Using Multiplexers

## $2 \times 2$ Crossbar switch


[ Figure 4.5a from the textbook ]

## $2 \times 2$ Crossbar switch



## Implementation of a $2 \times 2$ crossbar switch with multiplexers


[ Figure 4.5 b from the textbook ]

## Implementation of a $2 \times 2$ crossbar switch with multiplexers



## Implementation of a $2 \times 2$ crossbar switch with multiplexers



## Implementation of a logic function with a 4x1 multiplexer

| $w_{1}$ | $w_{2}$ | $f$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |


[ Figure 4.6a from the textbook ]

## Implementation of the same logic function with a $2 \times 1$ multiplexer


(b) Modified truth table

(c) Circuit

## The XOR Logic Gate



(b) Truth table

## The XOR Logic Gate


(a) Two switches that control a light
(c) Logic network

(b) Truth table


(d) XOR gate symbol
[ Figure 2.11 from the textbook]

Implementation of the XOR Logic Gate with a 2-to-1 multiplexer and one NOT


## Implementation of the XOR Logic Gate with a 2-to-1 multiplexer and one NOT



## Implementation of the XOR Logic Gate with a 2-to-1 multiplexer and one NOT



These two circuits are equivalent (the wires of the bottom AND gate are flipped)


## In other words, all four of these are equivalent!



## Implementation of another logic function

| $w_{1}$ | $w_{2}$ | $w_{3}$ | $f$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

## Implementation of another logic function

| $w_{1}$ | $w_{2}$ | $w_{3}$ | $f$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

## Implementation of another logic function



## Implementation of another logic function


[Figure 4.7 from the textbook]

Another Example (3-input XOR)

## Implementation of 3-input XOR with 2-to-1 Multiplexers

| $w_{1}$ | $w_{2}$ | $w_{3}$ | $f$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

## Implementation of 3-input XOR with 2-to-1 Multiplexers

$\left.\begin{array}{l|ll|l}w_{1} & w_{2} & w_{3} & f \\ \hline 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0\end{array}\right\} w_{2} \oplus w_{3}$

## Implementation of 3-input XOR with 2-to-1 Multiplexers


(a) Truth table
(b) Circuit


## Implementation of 3-input XOR with 2-to-1 Multiplexers

$\left.\begin{array}{l|l|l|l}w_{1} & w_{2} & w_{3} & f \\ \hline 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0\end{array}\right] \quad \mathbf{W}_{\mathbf{3}}$
(a) Truth table

(b) Circuit

## Implementation of 3-input XOR with a 4-to-1 Multiplexer

| $w_{1}$ | $w_{2}$ | $w_{3}$ | $f$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

## Implementation of 3-input XOR with a 4-to-1 Multiplexer

| $w_{1}$ | $w_{2}$ | $w_{3}$ | $f$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

## Implementation of 3-input XOR with a 4-to-1 Multiplexer

$\left.\begin{array}{ll|l|l}w_{1} & w_{2} & w_{3} & f \\ \hline 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1\end{array}\right\} w_{3}$

## Implementation of 3-input XOR with a 4-to-1 Multiplexer

$\left.\begin{array}{ll|l|l}w_{1} & w_{2} & w_{3} & f \\ \hline 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1\end{array}\right\} w_{3}$
(a) Truth table

(b) Circuit
[ Figure 4.9 from the textbook ]

## Multiplexor Synthesis Using Shannon's Expansion

## Three-input majority function

| $w_{1}$ | $w_{2}$ | $w_{3}$ | $f$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

## Three-input majority function



## Three-input majority function



## Three-input majority function


(b) Truth table

(b) Circuit
[Figure 4.10a from the textbook]

## Three-input majority function

$$
\begin{aligned}
f & =\bar{w}_{1} w_{2} w_{3}+w_{1} \bar{w}_{2} w_{3}+w_{1} w_{2} \bar{w}_{3}+w_{1} w_{2} w_{3} \\
f & =\bar{w}_{1}\left(w_{2} w_{3}\right)+w_{1}\left(\bar{w}_{2} w_{3}+w_{2} \bar{w}_{3}+w_{2} w_{3}\right) \\
& =\bar{w}_{1}\left(w_{2} w_{3}\right)+w_{1}\left(w_{2}+w_{3}\right)
\end{aligned}
$$



## Shannon's Expansion Theorem

Any Boolean function $f\left(w_{1}, \ldots, w_{n}\right)$ can be rewritten in the form:

$$
f\left(w_{1}, w_{2}, \ldots, w_{n}\right)=\bar{w}_{1} \cdot f\left(0, w_{2}, \ldots, w_{n}\right)+w_{1} \cdot f\left(1, w_{2}, \ldots, w_{n}\right)
$$

## Shannon's Expansion Theorem

Any Boolean function $f\left(w_{1}, \ldots, w_{n}\right)$ can be rewritten in the form:

$$
f\left(w_{1}, w_{2}, \ldots, w_{n}\right)=\bar{w}_{1} \cdot f\left(0, w_{2}, \ldots, w_{n}\right)+w_{1} \cdot f\left(1, w_{2}, \ldots, w_{n}\right)
$$

$$
f=\bar{w}_{1} f_{\bar{w}_{1}}+w_{1} f_{w_{1}}
$$

## Shannon's Expansion Theorem

Any Boolean function $f\left(w_{1}, \ldots, w_{n}\right)$ can be rewritten in the form:

$$
f\left(w_{1}, w_{2}, \ldots, w_{n}\right)=\bar{w}_{1} \cdot f\left(0, w_{2}, \ldots, w_{n}\right)+w_{1} \cdot f\left(1, w_{2}, \ldots, w_{n}\right)
$$



## Shannon's Expansion Theorem (Example)

$$
f\left(w_{1}, w_{2}, w_{3}\right)=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3}
$$

## Shannon's Expansion Theorem (Example)

$$
\begin{aligned}
& f\left(w_{1}, w_{2}, w_{3}\right)=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3} \\
& f\left(w_{1}, w_{2}, w_{3}\right)=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3}\left(\bar{w}_{1}+w_{1}\right)
\end{aligned}
$$

## Shannon's Expansion Theorem (Example)

$$
\begin{aligned}
& f\left(w_{1}, w_{2}, w_{3}\right)=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3} \\
& f\left(w_{1}, w_{2}, w_{3}\right)=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3}\left(\bar{w}_{1}+w_{1}\right) \\
& f=\bar{w}_{1}\left(0 \cdot w_{2}+0 \cdot w_{3}+w_{2} w_{3}\right)+w_{1}\left(1 \cdot w_{2}+1 \cdot w_{3}+w_{2} w_{3}\right) \\
& \quad=\bar{w}_{1}\left(w_{2} w_{3}\right)+w_{1}\left(w_{2}+w_{3}\right)
\end{aligned}
$$

## Shannon's Expansion Theorem (In terms of more than one variable)

$$
\begin{aligned}
f\left(w_{1}, \ldots, w_{n}\right)= & \bar{w}_{1} \bar{w}_{2} \cdot f\left(0,0, w_{3}, \ldots, w_{n}\right)+\bar{w}_{1} w_{2} \cdot f\left(0,1, w_{3}, \ldots, w_{n}\right) \\
& +w_{1} \bar{w}_{2} \cdot f\left(1,0, w_{3}, \ldots, w_{n}\right)+w_{1} w_{2} \cdot f\left(1,1, w_{3}, \ldots, w_{n}\right)
\end{aligned}
$$

This form is suitable for implementation with a $4 \times 1$ multiplexer.

## Another Example

# Factor and implement the following function with a $2 \times 1$ multiplexer 

$$
f=\bar{w}_{1} \bar{w}_{3}+w_{1} w_{2}+w_{1} w_{3}
$$

# Factor and implement the following function with a $2 \times 1$ multiplexer 

$$
f=\bar{w}_{1} \bar{w}_{3}+w_{1} w_{2}+w_{1} w_{3}
$$

$$
\begin{aligned}
f & =\bar{w}_{1} f_{\bar{w}_{1}}+w_{1} f_{w_{1}} \\
& =\bar{w}_{1}\left(\bar{w}_{3}\right)+w_{1}\left(w_{2}+w_{3}\right)
\end{aligned}
$$

## Factor and implement the following function with a $2 \times 1$ multiplexer



$$
\begin{aligned}
f & =\bar{w}_{1} f_{\bar{w}_{1}}+w_{1} f_{w_{1}} \\
& =\bar{w}_{1}\left(\bar{w}_{3}\right)+w_{1}\left(w_{2}+w_{3}\right)
\end{aligned}
$$

[ Figure 4.11a from the textbook]

# Factor and implement the following function with a 4x1 multiplexer 

$$
f=\bar{w}_{1} \bar{w}_{3}+w_{1} w_{2}+w_{1} w_{3}
$$

## Factor and implement the following function with a $4 \times 1$ multiplexer

$$
f=\bar{w}_{1} \bar{w}_{3}+w_{1} w_{2}+w_{1} w_{3}
$$

$$
\begin{aligned}
f & =\bar{w}_{1} \bar{w}_{2} f_{\bar{w}_{1} \bar{w}_{2}}+\bar{w}_{1} w_{2} f_{\bar{w}_{1}}+w_{1} \bar{w}_{2} f_{w_{1} \bar{w}_{2}}+w_{1} w_{2} f_{w_{1} w_{2}} \\
& =\bar{w}_{1} \bar{w}_{2}\left(\bar{w}_{3}\right)+\bar{w}_{1} w_{2}\left(\bar{w}_{3}\right)+w_{1} \bar{w}_{2}\left(w_{3}\right)+w_{1} w_{2}(1)
\end{aligned}
$$

## Factor and implement the following function with a 4x1 multiplexer

$$
\begin{aligned}
& f=\bar{w}_{1} \bar{w}_{2} f_{\bar{w}_{1} \bar{w}_{2}}+\bar{w}_{1} w_{2} f_{\bar{w}_{1} w_{2}}+w_{1} \bar{w}_{2} f_{w_{1} \bar{w}_{2}}+w_{1} w_{2} f_{w_{1} w_{2}} \\
& =\bar{w}_{1} \bar{w}_{2}\left(w_{3}\right)+\bar{w}_{1} w_{2}\left(\bar{w}_{3}\right)+w_{1} \bar{w}_{2}\left(w_{3}\right)+w_{1} w_{2}(1)
\end{aligned}
$$

[ Figure 4.11b from the textbook]

## Yet Another Example

# Factor and implement the following function using only $\mathbf{2 x 1}$ multiplexers 

$$
f=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3}
$$

## Factor and implement the following function using only $2 \times 1$ multiplexers

$$
f=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3}
$$

$$
\begin{aligned}
f & =\bar{w}_{1}\left(w_{2} w_{3}\right)+w_{1}\left(w_{2}+w_{3}+w_{2} w_{3}\right) \\
& =\bar{w}_{1}\left(w_{2} w_{3}\right)+w_{1}\left(w_{2}+w_{3}\right)
\end{aligned}
$$

## Factor and implement the following function using only $2 \times 1$ multiplexers

$$
f=w_{1} w_{2}+w_{1} w_{3}+w_{2} w_{3}
$$

$$
\begin{gathered}
f=\bar{w}_{1}\left(w_{2} w_{3}\right)+w_{1}\left(w_{2}+w_{3}+w_{2} w_{3}\right) \\
=\bar{w}_{1}(\underbrace{w_{2} w_{3}})+w_{1}(\underbrace{w_{2}+w_{3}}) \\
\quad g=w_{2} w_{3} \quad h=w_{2}+w_{3}
\end{gathered}
$$

## Factor and implement the following function using only $2 \times 1$ multiplexers



$$
\begin{gathered}
f=\bar{w}_{1}\left(w_{2} w_{3}\right)+w_{1}\left(w_{2}+w_{3}+w_{2} w_{3}\right) \\
=\bar{w}_{1}(\underbrace{w_{2} w_{3}})+w_{1}(\underbrace{w_{2}+w_{3}}) \\
\quad g=w_{2} w_{3} \quad h=w_{2}+w_{3}
\end{gathered}
$$

# Factor and implement the following function using only $\mathbf{2 x 1}$ multiplexers 

$$
g=w_{2} w_{3}
$$

$$
h=w_{2}+w_{3}
$$

## Factor and implement the following function using only $\mathbf{2 x 1}$ multiplexers

$$
\begin{array}{cc}
g=w_{2} w_{3} & h=w_{2}+w_{3} \\
\downarrow=\bar{w}_{2}(0)+w_{2}\left(w_{3}\right) & h=\bar{w}_{2}\left(w_{3}\right)+w_{2}(1)
\end{array}
$$

## Factor and implement the following function using only $\mathbf{2 x 1}$ multiplexers




$$
g=\bar{w}_{2}(0)+w_{2}\left(w_{3}\right)
$$

$$
h=\bar{w}_{2}\left(w_{3}\right)+w_{2}(1)
$$

## Finally, we are ready to draw the circuit



Finally, we are ready to draw the circuit


## Finally, we are ready to draw the circuit


[ Figure 4.12 from the textbook ]

## Questions?

## THE END

