

## CprE 281: Digital Logic

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http://www.ece.iastate.edu/~alexs/classes/

## Registers and Counters

CprE 281: Digital Logic
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## Administrative Stuff

- The second midterm is this Friday.
- Homework 8 is due today.
- Homework 9 is out. It is due on Mon Nov 6.
- No HW due next Monday


## Administrative Stuff

- Midterm Exam \#2
- When: Friday October 27 @ 4pm.
- Where: This classroom
- What: Chapters 1, 2, 3, 4 and 5.1-5.8
- The exam will be open book and open notes (you can bring up to 3 pages of handwritten notes).


## Registers

## Register (Definition)

An n-bit structure consisting of flip-flops

## Parallel-Access Register

## 1-Bit Parallel-Access Register



## 1-Bit Parallel-Access Register



The 2-to-1 multiplexer is used to select whether to load a new value into the $D$ flip-flop or to retain the old value.

The output of this circuit is the Q output of the flip-flop.

## 1-Bit Parallel-Access Register



If Load $=\mathbf{0}$, then retain the old value.

If Load = 1, then load the new value from In.

## 2-Bit Parallel-Access Register



## 2-Bit Parallel-Access Register



## 3-Bit Parallel-Access Register



Notice that all flip-flops are on the same clock cycle.

## 3-Bit Parallel-Access Register

Parallel Output


## 4-Bit Parallel-Access Register



## 4-Bit Parallel-Access Register



## Shift Register

## A simple shift register


[ Figure 5.17a from the textbook]

## A simple shift register



Positive-edge-triggered
D Flip-Flop

## A simple shift register



## A simple shift register



## A simple shift register



## A simple shift register



## A simple shift register



## A simple shift register



## A simple shift register



## A simple shift register



## A simple shift register



## A simple shift register



Clock


## A simple shift register



Clock


## A simple shift register



Clock


## A simple shift register



Clock


## A simple shift register


(a) Circuit

|  | In | $\mathrm{Q}_{1}$ | $\mathrm{Q}_{2}$ | $\mathrm{Q}_{3}$ | $\mathrm{Q}_{4}=$ Out |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $t_{0}$ | 1 | 0 | 0 | 0 | 0 |
| $t_{1}$ | 0 | 1 | 0 | 0 | 0 |
| $t_{2}$ | 1 | 0 | 1 | 0 | 0 |
| $t_{3}$ | 1 | 1 | 0 | 1 | 0 |
| $t_{4}$ | 1 | 1 | 1 | 0 | 1 |
| $t_{5}$ | 0 | 1 | 1 | 1 | 0 |
| $t_{6}$ | 0 | 0 | 1 | 1 | 1 |
| $t_{7}$ | 0 | 0 | 0 | 1 | 1 |

(b) A sample sequence
[ Figure 5.17 from the textbook]

## Parallel-Access Shift Register

## Parallel-access shift register


[ Figure 5.18 from the textbook ]

## Parallel-access shift register



When Load=0, this behaves like a shift register.
[ Figure 5.18 from the textbook ]

## Parallel-access shift register



When Load=1, this behaves like a parallel-access register.
[ Figure 5.18 from the textbook ]

## Shift Register With Parallel Load and Enable

## A shift register with parallel load and enable control inputs


[ Figure 5.59 from the textbook ]

## A shift register with parallel load and enable control inputs



The directions of the input and output lines are switched relative to the previous slides.
[ Figure 5.59 from the textbook ]

## A shift register with parallel load and enable control inputs


[ Figure 5.59 from the textbook ]

## A shift register with parallel load and enable control inputs


[ Figure 5.59 from the textbook ]

## A shift register with parallel load and enable control inputs


[ Figure 5.59 from the textbook ]

## A shift register with parallel load and enable control inputs


[ Figure 5.59 from the textbook ]

## A shift register with parallel load and enable control inputs


[ Figure 5.59 from the textbook ]

# Multiplexer Tricks <br> (select one of two 2-bit numbers) 

## Select Either $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$



## Select Either $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$



## Select Either $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$



# Multiplexer Tricks (select one of four 2-bit numbers) 

Select $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$ or $C=C_{1} C_{0}$ or $D=D_{1} D_{0}$


Select $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$ or $C=C_{1} C_{0}$ or $D=D_{1} D_{0}$


Select $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$ or $C=C_{1} C_{0}$ or $D=D_{1} D_{0}$


Select $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$ or $C=C_{1} C_{0}$ or $D=D_{1} D_{0}$


Select $A=A_{1} A_{0}$ or $B=B_{1} B_{0}$ or $C=C_{1} C_{0}$ or $D=D_{1} D_{0}$


## Register File

Complete the following circuit diagram to implement a register file with four 2-bit registers, one write port, one read port, and one write enable line.


$\begin{array}{ll}-w_{0} & y_{0}- \\ w_{1} & y_{1} \\ E n & y_{2}- \\ y_{3}\end{array}-$


| $=0$ | 1 | 1 |
| :---: | :---: | :---: |
| 0 | 0 | 0 |



Register 1


Register 2


Register 3
$\begin{array}{ll}-w_{0} & y_{0} \\ w_{1} & y_{1} \\ -E n & y_{2} \\ E n & y_{3}\end{array}-$




Register B


Register C


Register D


$\stackrel{1111}{=598}$


Register B


Register C


Register D

$\stackrel{\mid 1111}{=528}$
$\stackrel{1111}{=528}$







## Another Register File

## Register File

- Register file is a unit containing registers
- r can be 4, 8, 16, 32, etc.
- Each register has $\mathbf{n}$ bits
- $n$ can be $4,8,16,32$, etc.
- n defines the data path width
- Output ports (DATA1 and DATA2) are used for reading the register file
- Any register can be read from any of the ports
- Each port needs a $\log _{2} r$ bits to specify the read
 address (RA1 and RA2)
- Input port (LD_DATA) is used for writing data to the register file
- Write address is also specified by $\log _{2} r$ bits (WA)
- Writing is enabled by a 1-bit signal (WR)


## Register File: Exercise

- Suppose that a register file
- contains 32 registers
- width of data path is 16 bits (i.e., each register has 16 bits)
- How many bits are there for each of the signals?
- RA1

5


- RA2
- DATA1
- DATA2
- WA
- LD_DATA

5

- WR

16
WR

## Register file design

- We will design an eight-register file with 4-bit wide registers
- A single 4-bit register and its abstraction are shown below

- We have to use eight such registers to make an eight register file

- How many bits are required to specify a register address?


## Reading Circuit

- A 3-bit register address, RA, specifies which register is to be read
- For each output port, we need one 8-to-1 4-bit multiplier

Register
Address


## Adding write control to register file

- To write to any register, we need the register's address (WA) and a write register signal (WR)
- A 3-bit write address is decoded if write register signal is present
- One of the eight registers gets a LD signal from the decoder



## Register File (More Examples)

## Register File



Gray lines are 1-bit signals
Black lines are 10 -bit signals




## Counters

## T Flip-Flop (circuit and graphical symbol)


[ Figure 5.15a,c from the textbook]

## The output of the T Flip-Flop divides the frequency of the clock by 2



## The output of the T Flip-Flop divides the frequency of the clock by 2



## A three-bit down-counter


[ Figure 5.20 from the textbook ]

## A three-bit down-counter



The first flip-flop changes
on the positive edge of the clock
[ Figure 5.20 from the textbook ]

## A three-bit down-counter



The first flip-flop changes on the positive edge of the clock

The second flip-flop changes on the positive edge of $\mathrm{Q}_{0}$

## A three-bit down-counter



The first flip-flop changes on the positive edge of the clock

The second flip-flop changes The third flip-flop changes on the positive edge of $\mathrm{Q}_{0}$ on the positive edge of $\mathrm{Q}_{1}$

## A three-bit down-counter


(b) Timing diagram
[ Figure 5.20 from the textbook]

## A three-bit down-counter



## A three-bit down-counter



## A three-bit down-counter



## A three-bit down-counter


(b) Timing diagram

## A three-bit down-counter



## A three-bit up-counter


[ Figure 5.19 from the textbook ]

## A three-bit up-counter



The first flip-flop changes
on the positive edge of the clock
[ Figure 5.19 from the textbook ]

## A three-bit up-counter



The first flip-flop changes on the positive edge of the clock

The second flip-flop changes on the positive edge of $\overline{\mathrm{Q}}_{0}$

## A three-bit up-counter



The first flip-flop changes on the positive edge of the clock

The second flip-flop changes The third flip-flop changes on the positive edge of $\overline{\mathrm{Q}}_{0} \quad$ on the positive edge of $\overline{\mathrm{Q}}_{1}$

## A three-bit up-counter


(a) Circuit

[ Figure 5.19 from the textbook ]

## A three-bit up-counter


[ Figure 5.19 from the textbook ]

## A three-bit up-counter


(a) Circuit

(b) Timing diagram
[ Figure 5.19 from the textbook ]

## A three-bit up-counter


(a) Circuit

(b) Timing diagram

## A three-bit up-counter


(a) Circuit

(b) Timing diagram

## A three-bit up-counter


(a) Circuit

(b) Timing diagram

## A three-bit up-counter


(a) Circuit

(b) Timing diagram

## A three-bit up-counter


(b) Timing diagram

## A three-bit up-counter


[ Figure 5.19 from the textbook ]

## Synchronous Counters

## A four-bit synchronous up-counter


[ Figure 5.21 from the textbook]

## A four-bit synchronous up-counter



The propagation delay through all AND gates combined must not exceed the clock period minus the setup time for the flip-flops

## A four-bit synchronous up-counter



(b) Timing diagram
[ Figure 5.21 from the textbook ]

## Derivation of the synchronous up-counter

| Clock cycle | $\mathrm{Q}_{2} \mathrm{Q}_{1} \mathrm{Q}_{0}$ |  |  |
| :---: | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 |
| 2 | 0 | 1 | 0 |
| 3 |  |  |  |
| 4 | 0 | 1 | 1 |
| 5 | 1 | 0 | 0 |
| 6 | 1 | 0 | 1 |
| 7 | 1 | 1 | 0 |
| 8 | 1 | 1 | 1 |

## Derivation of the synchronous up-counter

| Clock cycle | $\mathrm{Q}_{2} \mathrm{Q}_{1} \mathrm{Q}_{0}$ | $\mathrm{Q}_{1}$ changes |
| :---: | :---: | :---: |
| 0 | 0 0 0 |  |
| 1 | $\begin{array}{llll}0 & 0 & 1\end{array}$ | $\square \mathrm{Q}_{2}$ changes |
| 2 | $0 \quad 1 \quad 0 \leftharpoonup$ - |  |
| 3 | $\begin{array}{lll}0 & 1 & 1\end{array}$ |  |
| 4 | 100 |  |
| 5 | $1 \begin{array}{lll}1 & 0 & 1\end{array}$ |  |
| 6 | $110 \longrightarrow$ |  |
| 7 | $\begin{array}{lll}1 & 1 & 1\end{array}$ |  |
| 8 | $0 \quad 0 \quad 0 \leftharpoonup$ |  |
| $\mathrm{T}_{0}=1$ |  |  |
| $\mathrm{T}_{1}=\mathrm{Q}_{0}$ |  |  |
| $\mathrm{T}_{2}=\mathrm{Q}_{0} \mathrm{Q}_{1}$ |  |  |

[ Table 5.1 from the textbook ]

## A four-bit synchronous up-counter



$$
\begin{aligned}
& \mathrm{T}_{0}=1 \\
& \mathrm{~T}_{1}=\mathrm{Q}_{0} \\
& \mathrm{~T}_{2}=\mathrm{Q}_{0} \mathrm{Q}_{1}
\end{aligned}
$$

[ Figure 5.21 from the textbook ]

## In general we have

$$
\begin{aligned}
& \mathrm{T}_{0}=1 \\
& \mathrm{~T}_{1}=\mathrm{Q}_{0} \\
& \mathrm{~T}_{2}=\mathrm{Q}_{0} \mathrm{Q}_{1} \\
& \mathrm{~T}_{3}=\mathrm{Q}_{0} \mathrm{Q}_{1} \mathrm{Q}_{2} \\
& \ldots \\
& \mathrm{~T}_{\mathrm{n}}=\mathrm{Q}_{0} \mathrm{Q}_{1} \mathrm{Q}_{2} \ldots \mathrm{Q}_{\mathrm{n}-1}
\end{aligned}
$$

## Synchronous v.s. Asynchronous Clear

## 2-Bit Synchronous Up-Counter (without clear capability)



## 2-Bit Synchronous Up-Counter (with asynchronous clear)



## 2-Bit Synchronous Up-Counter (with asynchronous clear)



This is the same circuit but uses D Flip-Flops.

## 2-Bit Synchronous Up-Counter (with synchronous clear)



This counter can be cleared only on the positive clock edge.

## Adding Enable Capability

## A four-bit synchronous up-counter


[ Figure 5.21 from the textbook ]

## Inclusion of Enable and Clear Capability


[ Figure 5.22 from the textbook ]

## Inclusion of Enable and Clear Capability


[ Figure 5.22 from the textbook ]

## Providing an enable input for a D flip-flop


(a) Using a multiplexer

(b) Clock gating

## Synchronous Counter (with D Flip-Flops)

## A four-bit counter with D flip-flops


[ Figure 5.23 from the textbook ]

## Counters with Parallel Load

## A 4-bit up-counter with D flip-flops


[ Figure 5.23 from the textbook ]

## A 4-bit up-counter with D flip-flops


[ Figure 5.23 from the textbook ]

## Equivalent to this circuit with T flip-flops



## Equivalent to this circuit with T flip-flops



But has one extra output called Z , which can be used to connect two 4-bit counters to make an 8 -bit counter.

When $\mathrm{Z}=1$ the counter will go 0000 on the next clock edge, i.e., the outputs of all flip-flops are currently 1 (maximum count value).

## Counters with Parallel Load

## A counter with parallel-load capability


[ Figure 5.24 from the textbook]

## How to load the initial count value

Set the initial count on the parallel load lines
(in this case 5).


## How to zero a counter

Set "Load" to 1, to open the


## How to zero a counter



## Reset Synchronization

## Motivation

- An n-bit counter counts from $\mathbf{0 , 1}, \ldots, \mathbf{2 n}^{\mathbf{n}} \mathbf{- 1}$
- For example a 3-bit counter counts up as follow
- $0,1,2,3,4,5,6,7,0,1,2, \ldots$
- What if we want it to count like this
- $0,1,2,3,4,5,0,1,2,3,4,5,0,1, \ldots$
- In other words, what is the cycle is not a power of 2?


## What does this circuit do?


[ Figure 5.25a from the textbook]

## A modulo-6 counter with synchronous reset


(a) Circuit

(b) Timing diagram

## A modulo-6 counter with asynchronous reset


[ Figure 5.26 from the textbook ]

## A modulo-6 counter with asynchronous reset


(b) Timing diagram
[ Figure 5.26 from the textbook ]

## Questions?

## THE END

