

# **CprE 281: Digital Logic**

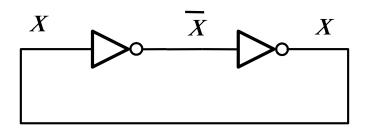
Instructor: Alexander Stoytchev

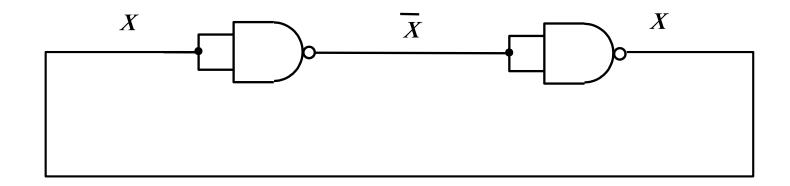
http://www.ece.iastate.edu/~alexs/classes/

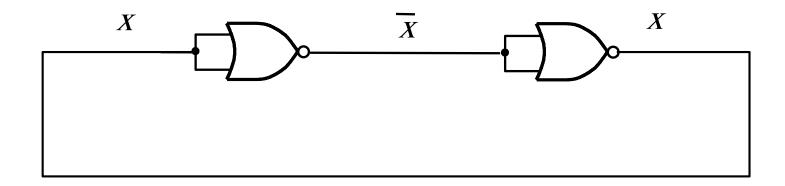
# D Flip-Flops

CprE 281: Digital Logic Iowa State University, Ames, IA Copyright © Alexander Stoytchev

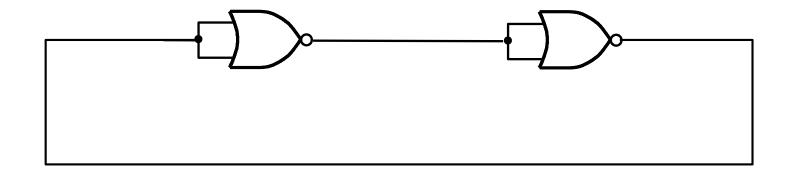
## **Quick Review**

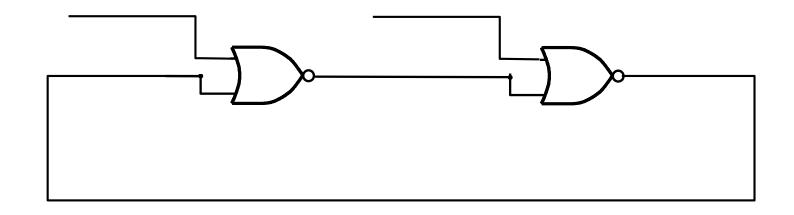


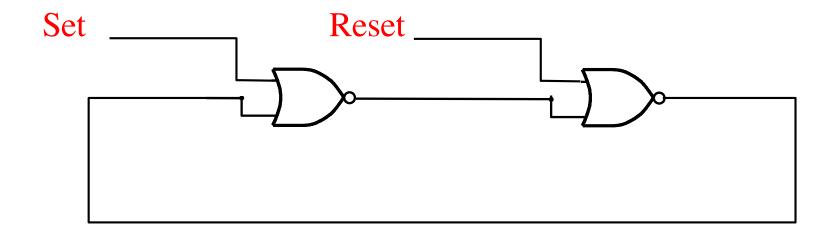




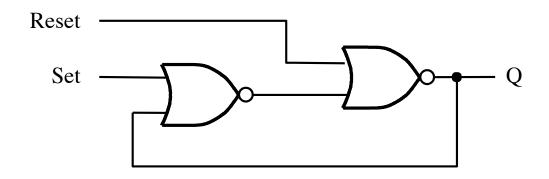
## **Basic Latch**



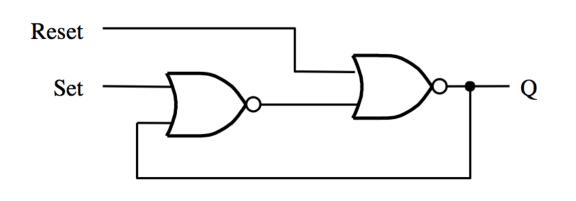


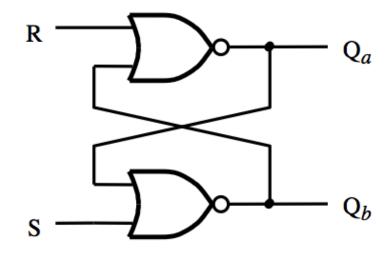


## A memory element with NOR gates



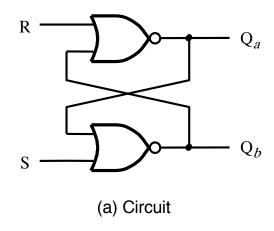
## Two Different Ways to Draw the Same Circuit

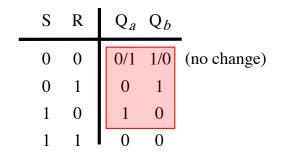




#### Circuit and Characteristic Table for the Basic Latch

Note that Q<sub>a</sub> and Q<sub>b</sub> are inverses of each other!



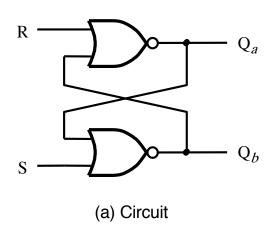


#### Oscillations and Undesirable States

- When S=1 and R=1 both outputs of the latch are equal to 0, i.e., Q<sub>a</sub>=0 and Q<sub>b</sub>=0.
- Thus, the two outputs are no longer complements of each other.
- This is undesirable as many of the circuits that we will build later with these latches rely on the assumption that the two outputs are always complements of each other.
- (This is obviously not the case for the basic latch, but we will patch it later to eliminate this problem).

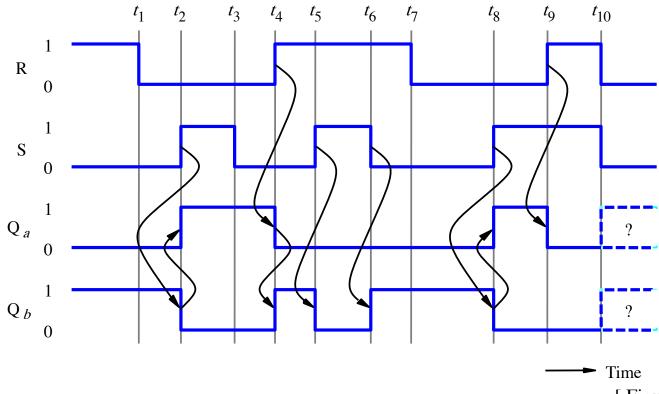
#### Oscillations and Undesirable States

- An even bigger problem occurs when we transition from S=R=1 to S=R=0.
- When S=R=1 we have  $Q_a=Q_b=0$ . After the transition to S=R=0, however, we get  $Q_a=Q_b=1$ , which would immediately cause  $Q_a=Q_b=0$ , and so on.
- If the gate delays and the wire lengths are identical, then this oscillation will continue forever.
- In practice, the oscillation dies down and the output settles into either  $Q_a=1$  and  $Q_b=0$  or  $Q_a=0$  and  $Q_b=1$ .
- The problem is that we can't predict which one of these two it will settle into.



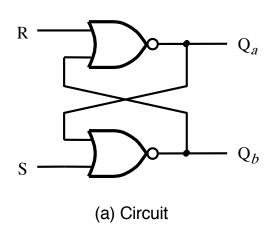
S	R	$Q_a$	$Q_b$	_
0	0	0/1	1/0	(no change)
0	1	0	1	
1	0	1	0	
1	1	0	0	

(b) Characteristic table



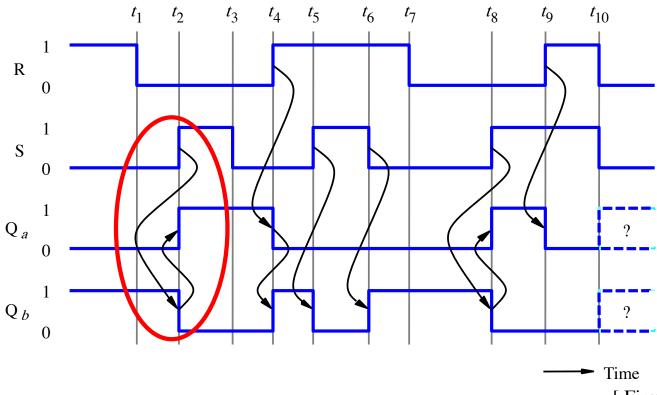
(c) Timing diagram

[ Figure 5.4 from the textbook ]



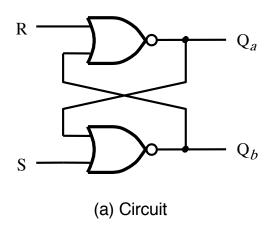
	$Q_b$	$Q_a$	R	S	
(no change)	1/0	0/1	0	0	
	1	0	1	0	
	0	1	0	1	
	0	0	1	1	

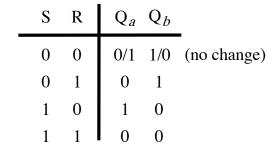
(b) Characteristic table



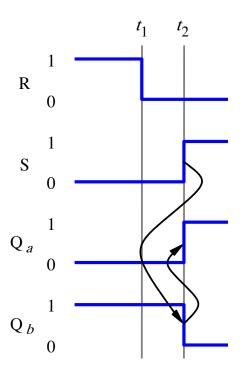
(c) Timing diagram

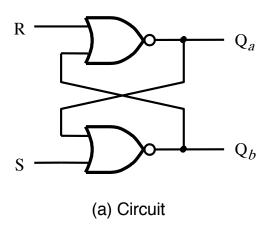
[ Figure 5.4 from the textbook ]

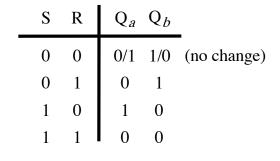




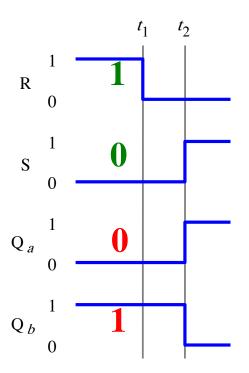
(b) Characteristic table

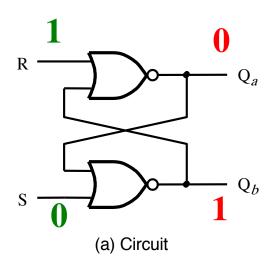


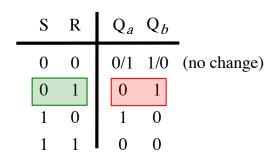




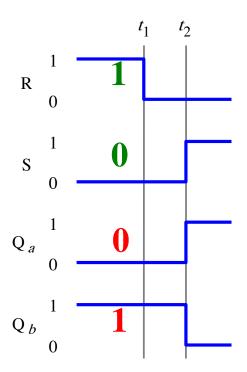
(b) Characteristic table

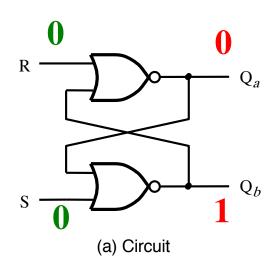


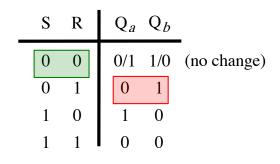




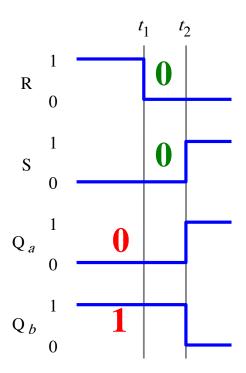
(b) Characteristic table

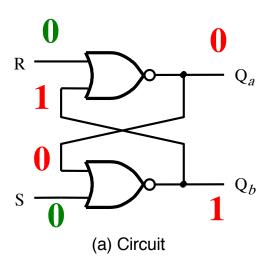


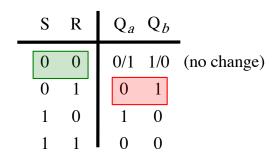




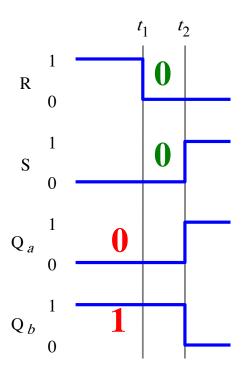
(b) Characteristic table

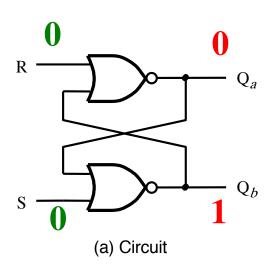


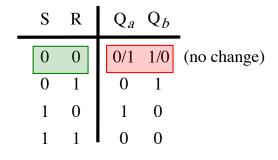


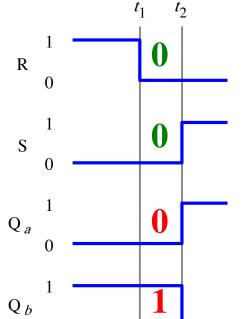


(b) Characteristic table

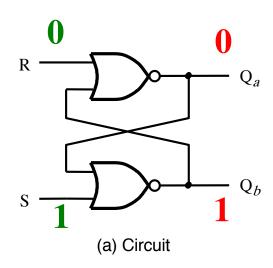


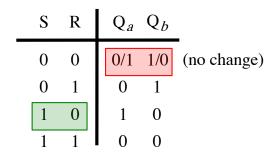




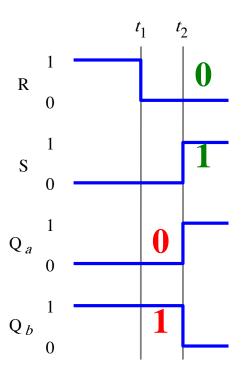


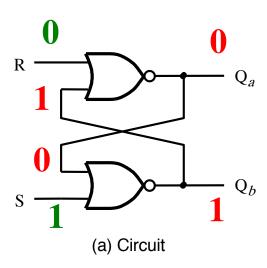
(b) Characteristic table

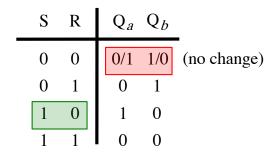




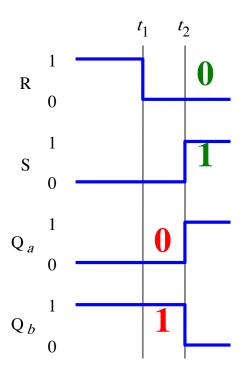
(b) Characteristic table

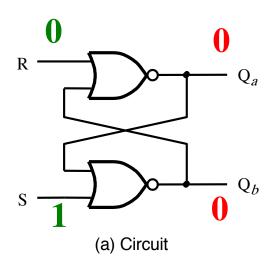






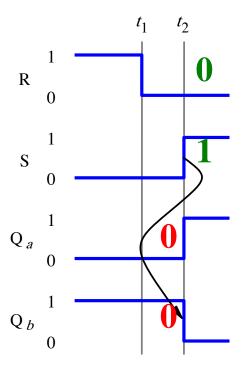
(b) Characteristic table



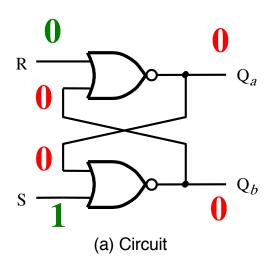


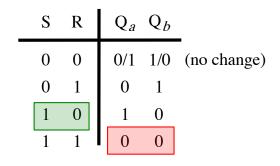
	S	R	$Q_a$	$Q_b$	_
•	0	0	0/1	1/0	(no change)
	0	1	0	1	
	1	0	_ 1	0	
	1	1	0	0	

(b) Characteristic table

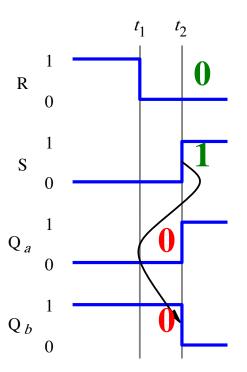


For a brief moment the latch goes through the undesirable state  $Q_a=0$  and  $Q_b=0$ .

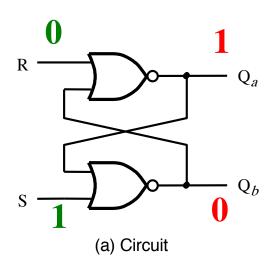


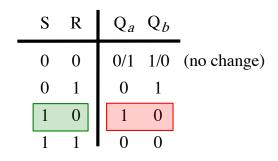


(b) Characteristic table

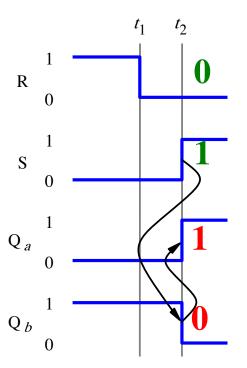


But these zeros loop around ...

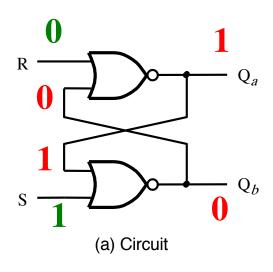


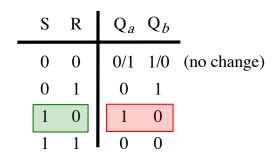


(b) Characteristic table

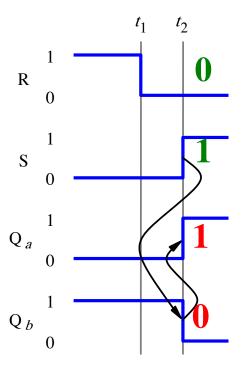


... and set it to  $Q_a=1$  and  $Q_b=0$ .

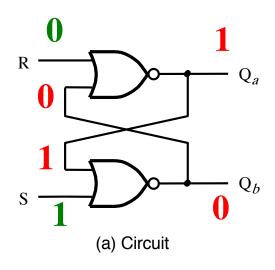


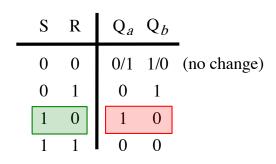


(b) Characteristic table

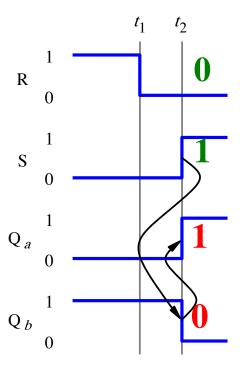


The new values also loop around ...

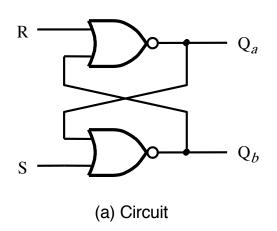




(b) Characteristic table

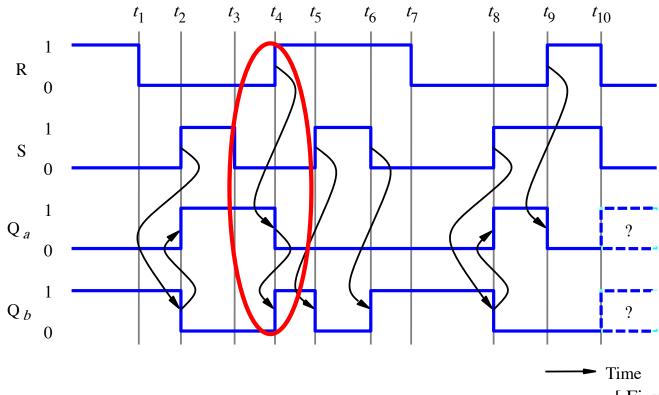


... but they leave the outputs the same.



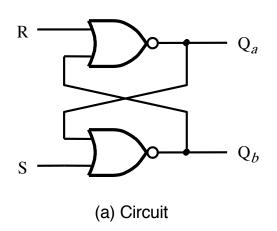
	$Q_b$	$Q_a$	R	S	
(no change)	1/0	0/1	0	0	
	1	0	1	0	
	0	1	0	1	
	0	0	1	1	

(b) Characteristic table

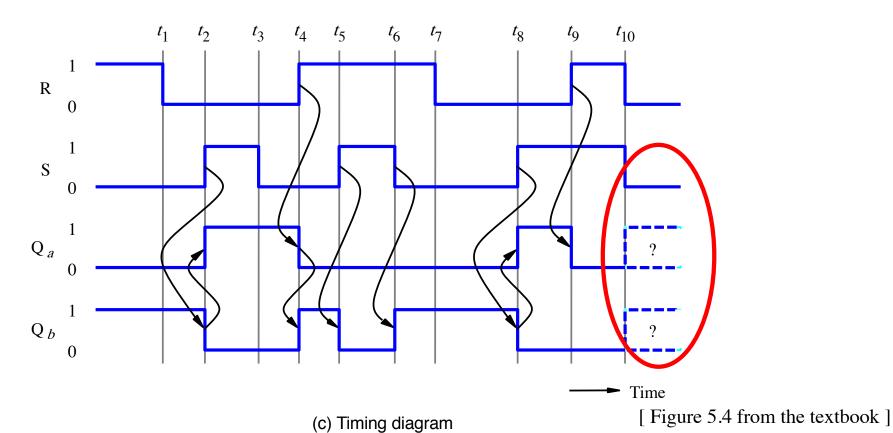


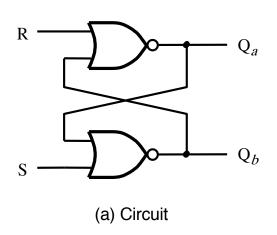
(c) Timing diagram

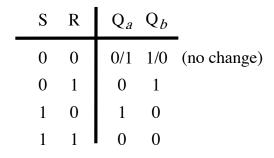
[ Figure 5.4 from the textbook ]

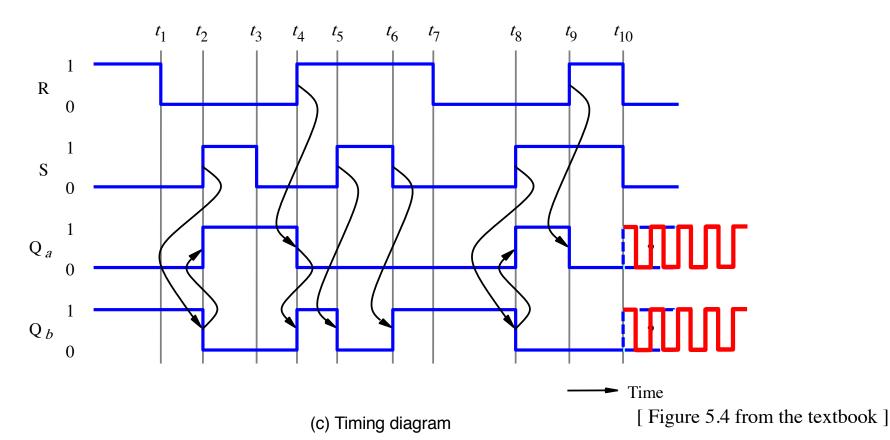


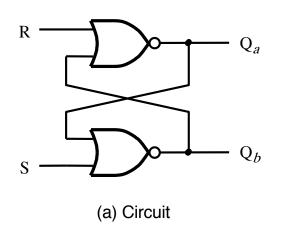
	S	R	$Q_a$	$Q_b$	_
•	0	0	0/1	1/0	(no change)
	0	1	0	1	
	1	0	1	0	
	1	1	0	0	

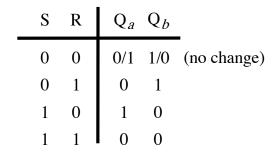


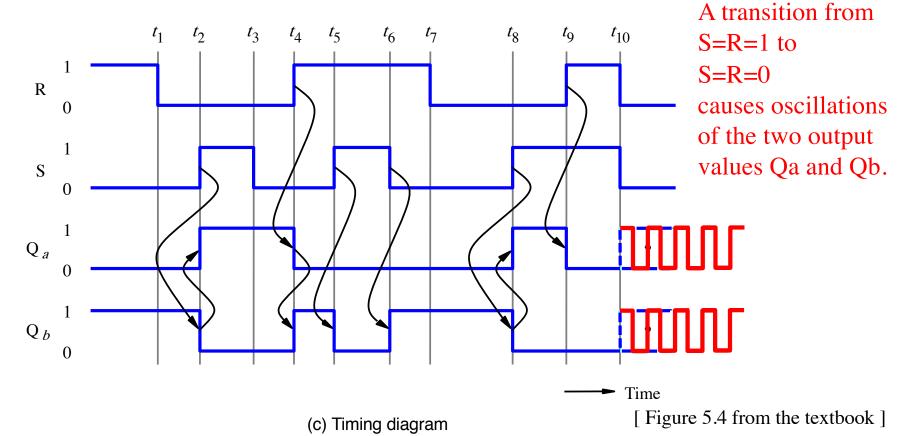






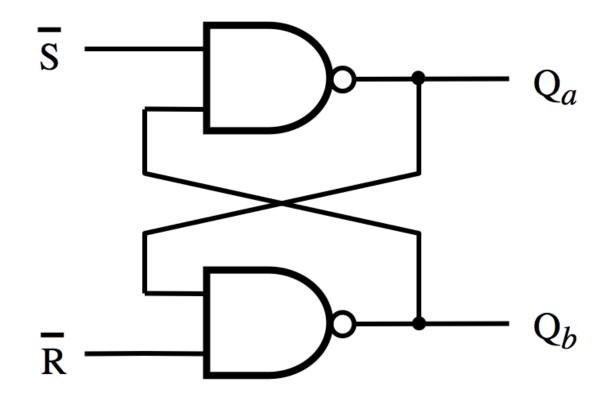




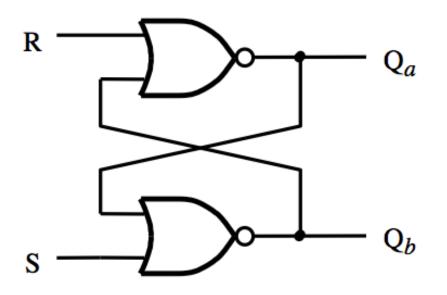


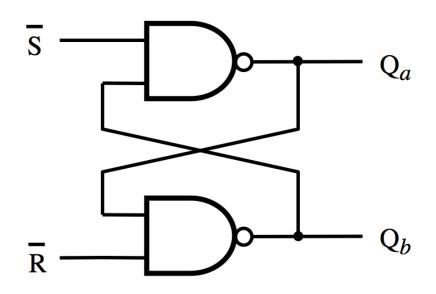
## **Basic Latch with NAND Gates**

#### **Circuit for the Basic Latch with NAND Gates**







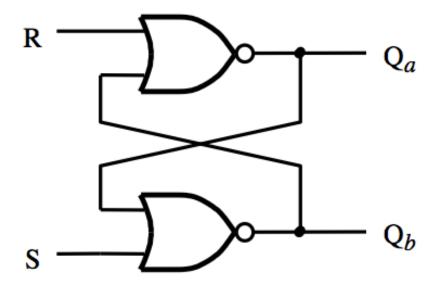


Notice that in the NAND case the two inputs are swapped and negated.

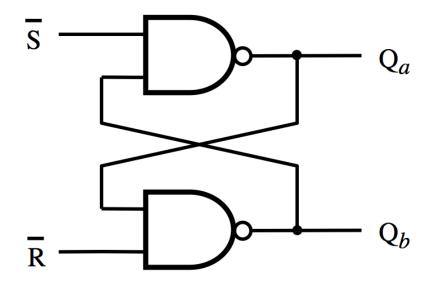
The labels of the outputs are the same in both cases.

# **Basic Latch**

# (with NOR Gates)



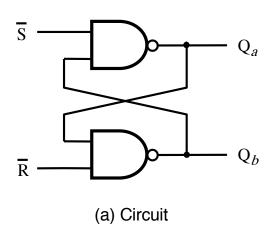
### **Basic Latch** (with NAND Gates)



SR Latch

SR Latch

#### **Circuit and Characteristic Table**



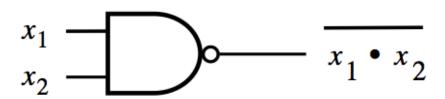
S	R	$Q_a$	$Q_b$	_
0	0	1	1	_
0	1	1	0	
1	0	0	1	
1	1	0/1	1/0	(no change)

(b) Characteristic table
(version 1)

S	R	$Q_a$	$Q_b$	_
0	0	0/1	1/0	(no change)
0	1	0	1	
1	0	1	0	
1	1	1	1	

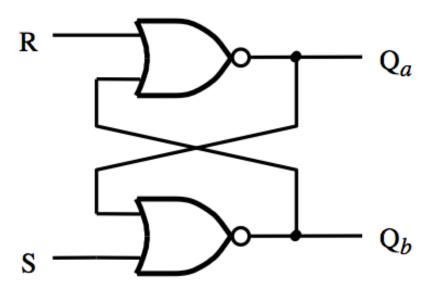
(c) Characteristic table (version 2)

#### NAND Gate



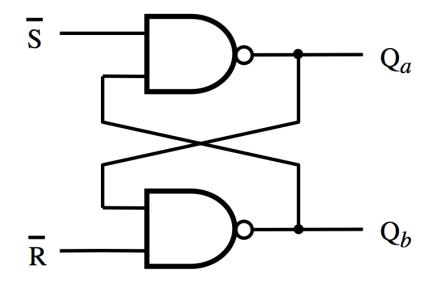
#### NAND Gate Truth table

$x_1$	$x_2$	f
0	0	1
0	1	1
1	0	1
1	1	0

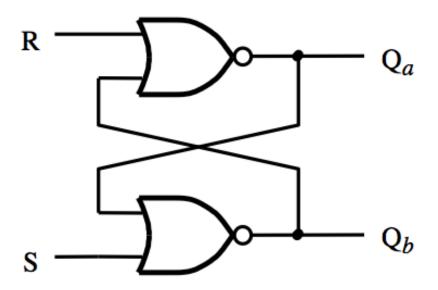


	S	R	$Q_a$	$Q_b$	_
•	0	0	0/1	1/0	(no change)
	0	1	0	1	
	1	0	1	0	
	1	1	0	0	

# Basic Latch (with NAND Gates)

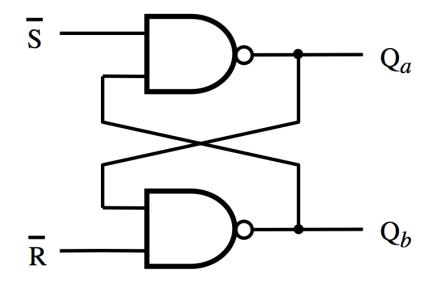


S	R	$Q_a$	$Q_b$	_
0	0	0/1	1/0	(no change)
0	1	0	1	
1	0	1	0	
1	1	1	1	

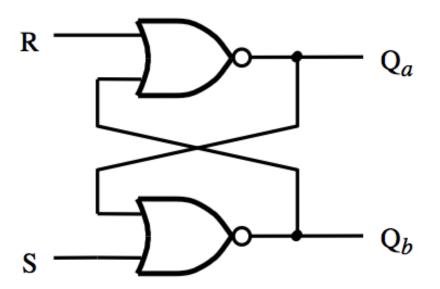


S	R	$Q_a$	$Q_b$	_	
0	0	0/1	1/0	(no change)	Latch
0	1	0 1 0	1		Reset
1	0	1	0		Set
1	1	0	0		Undesirable

# Basic Latch (with NAND Gates)

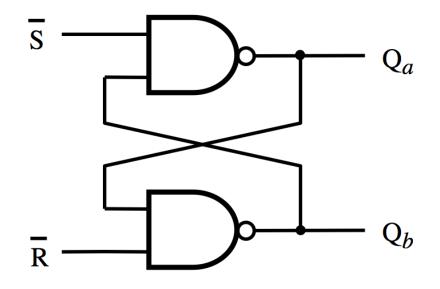


	S	R	$Q_a$	$Q_b$	_	
•	0	0	0/1	1/0	(no change)	Latch
	0	1	0	1		Reset
	1	0	1	0		Set
	1	1	1	1		Undesirable



S	R	$Q_a$	$Q_b$	_	
0	0	0/1	1/0	(no change)	Latch
0	1	0	1		Reset
1	0	$\begin{bmatrix} 0/1 \\ 0 \\ 1 \\ 0 \end{bmatrix}$	0		Set
1	1	0	0		Undesirable

# Basic Latch (with NAND Gates)



	S	R	$Q_a$	$Q_b$	_	
•	0	0	0/1	1/0	(no change)	Latch
	0	1	0	1		Reset
	1	0	1	0		Set
	1	1	1	1		Undesirable

The two characteristic tables are the same (except for the last row, which is the undesirable configuration).

#### Oscillations and Undesirable States

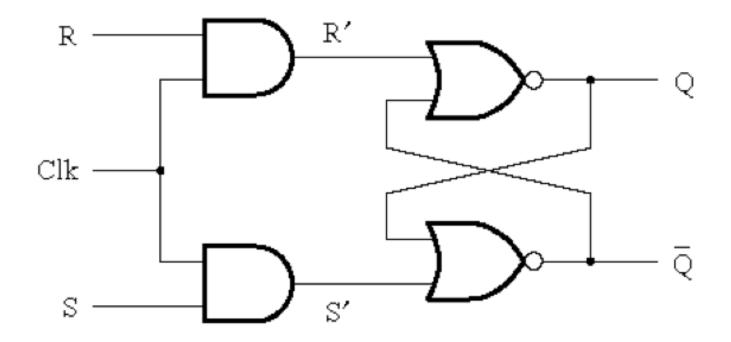
 The basic latch with NAND gates also suffers form oscillation problems, similar to the basic latch implemented with NOR gates.

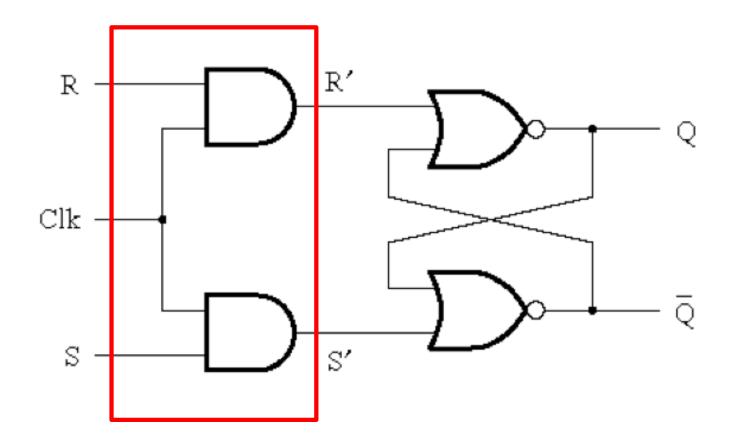
Try to do this analysis on your own.

# **Gated SR Latch**

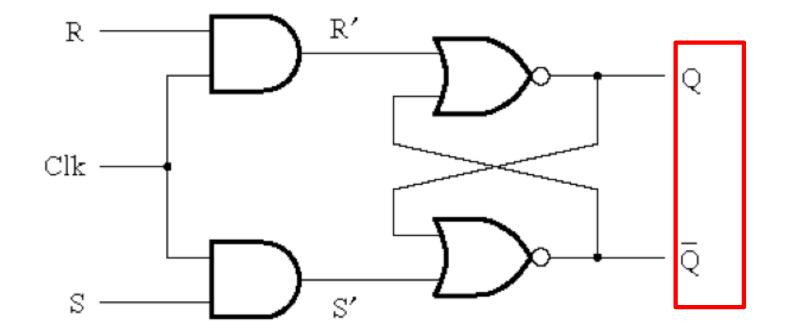
#### **Motivation**

- The basic latch changes its state when the input signals change
- It is hard to control when these input signals will change and thus it is hard to know when the latch may change its state.
- We want to have something like an Enable input.
- In this case it is called the "Clock" input because it is desirable for the state changes to be synchronized.



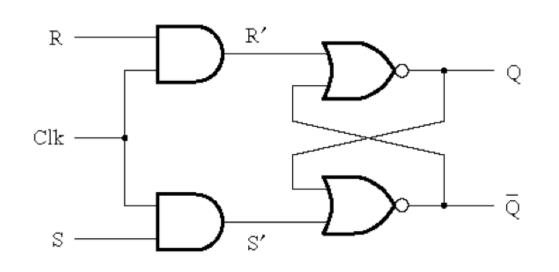


This is the "gate" of the gated latch



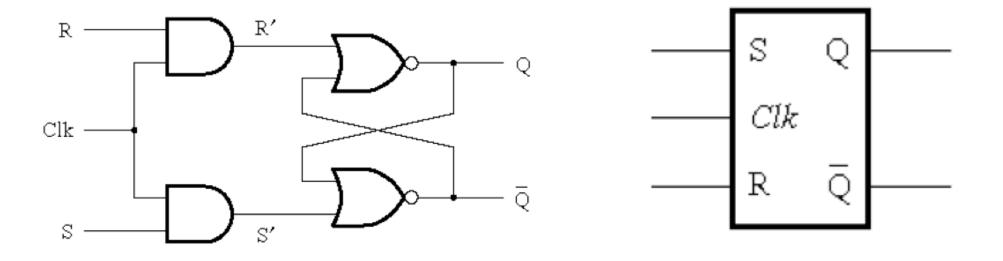
Notice that these are complements of each other

# Circuit Diagram and Characteristic Table for the Gated SR Latch

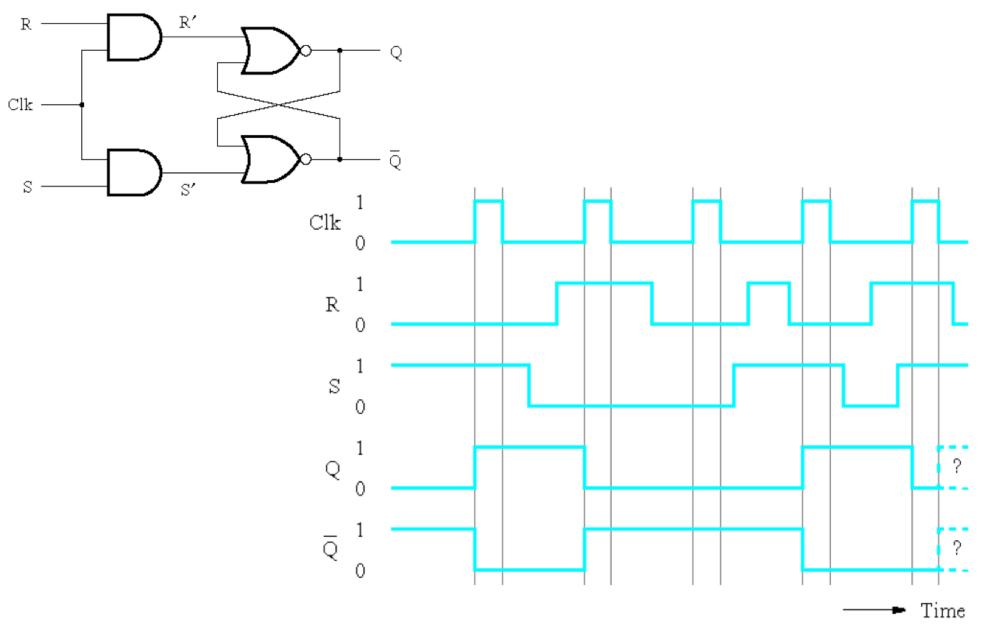


Clk	S	R	Q(t+1)
0	х	х	Q(t) (no change)
1	0	0	Q(t) (no change)
1	0	1	0
1	1	0	1
1	1	1	x

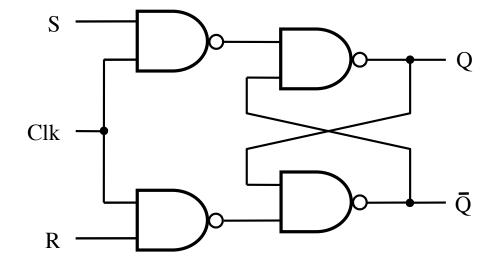
# Circuit Diagram and Graphical Symbol for the Gated SR Latch

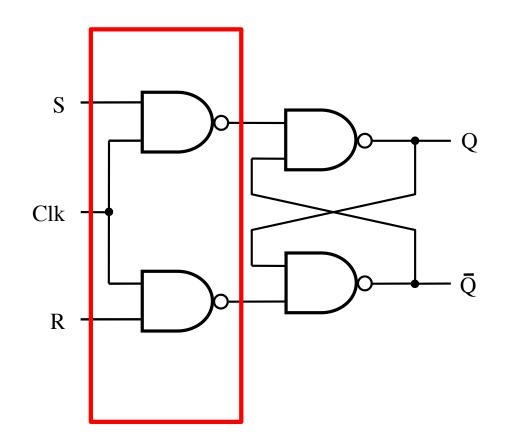


### Timing Diagram for the Gated SR Latch

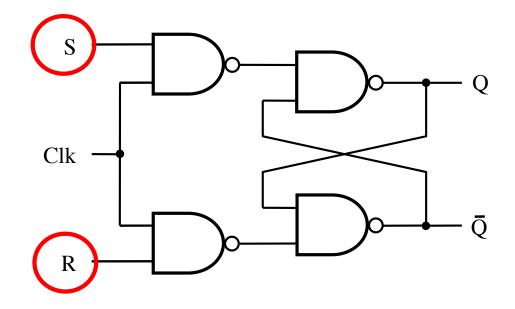


[ Figure 5.5c from the textbook ]

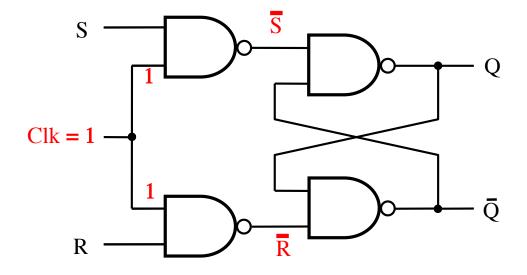




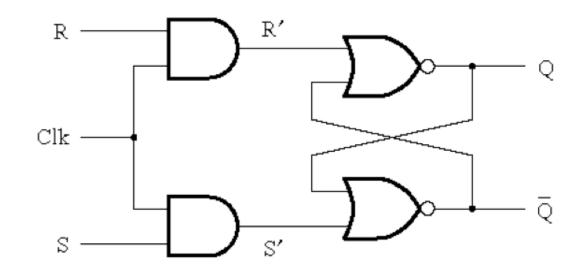
In this case the "gate" is constructed using NAND gates! Not AND gates.



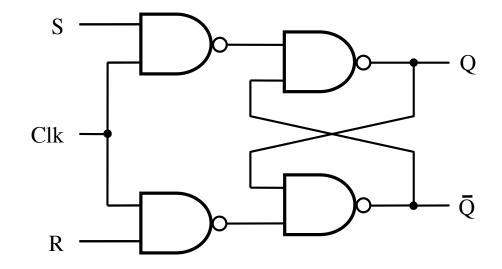
Also, notice that the positions of S and R are now swapped.

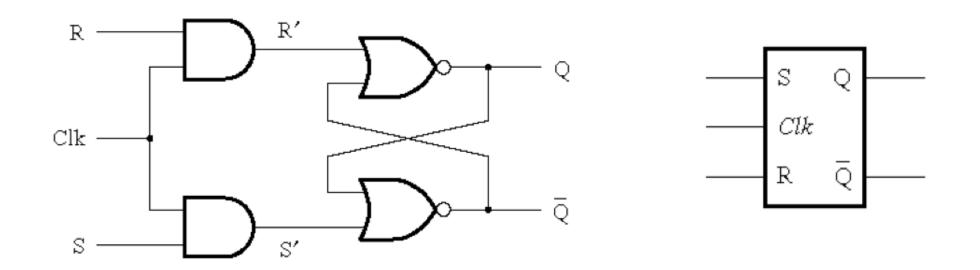


Finally, notice that when Clk=1 this turns into the basic latch with NAND gates, i.e., the  $\overline{SR}$  Latch.

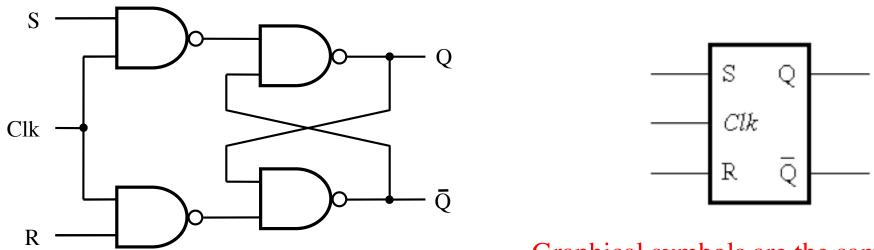


#### Gated SR latch with NAND gates

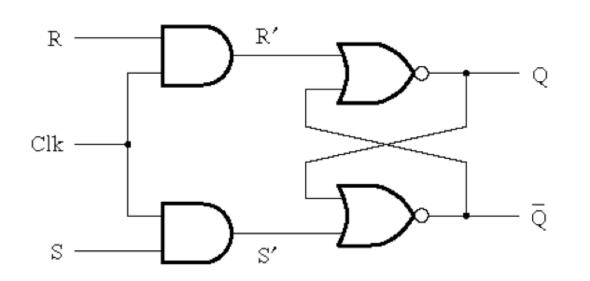




#### Gated SR latch with NAND gates

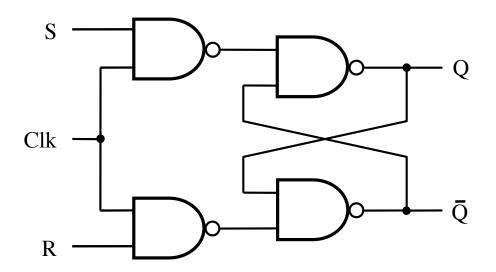


Graphical symbols are the same



Clk	S	R	Q(t+1)
0	x	x	Q(t) (no change)
1	0	0	Q(t) (no change)
1	0	1	0
1	1	0	1
1	1	1	x (undesirable)

#### Gated SR latch with NAND gates



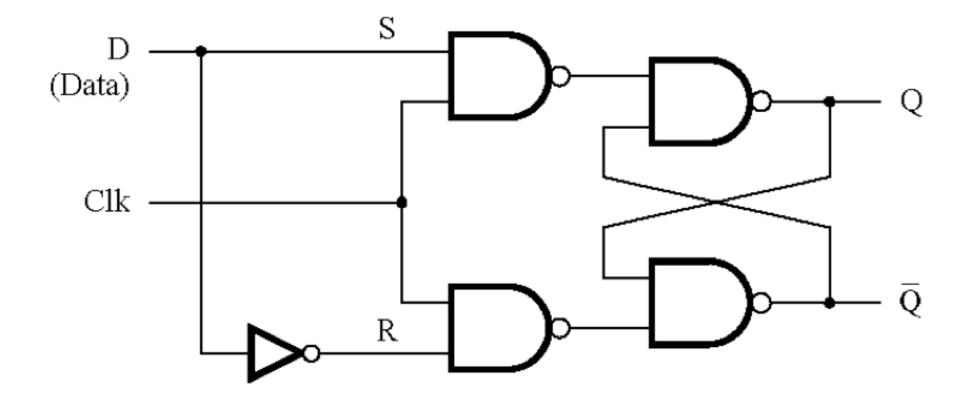
Clk	S	R	Q(t+1)
0	x	x	Q(t) (no change)
1	0	0	Q(t) (no change)
1	0	1	0
1	1	0	1
1	1	1	x (undesirable)

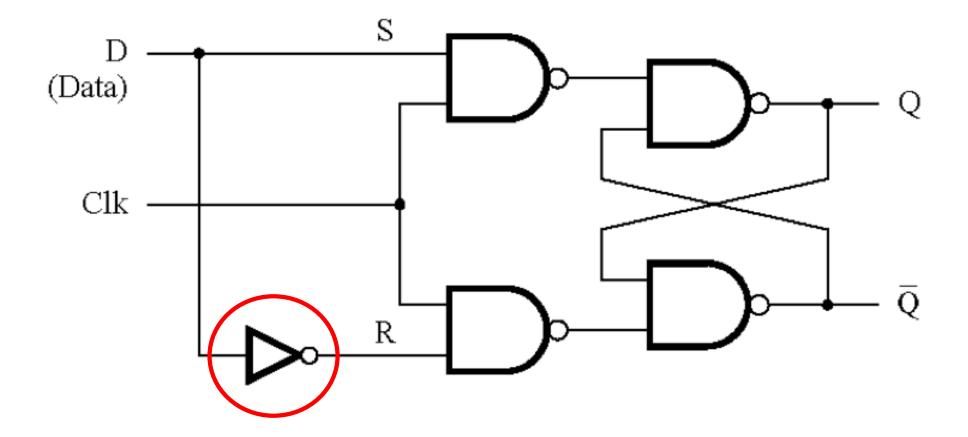
Characteristic tables are the same

## **Gated D Latch**

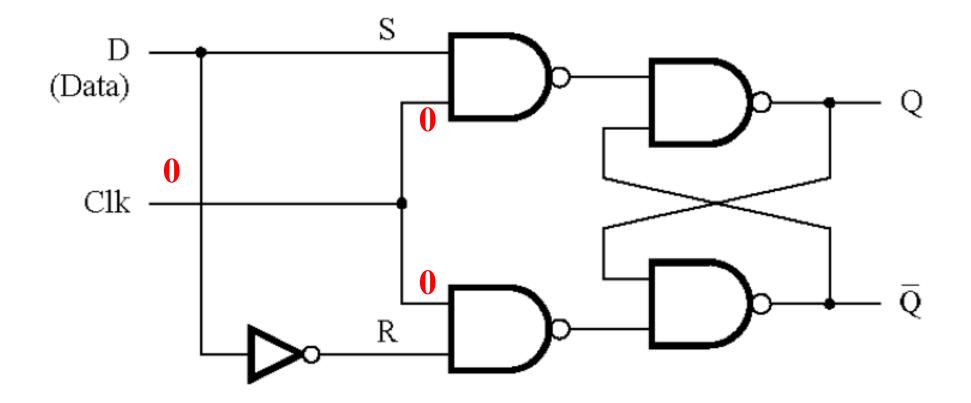
#### **Motivation**

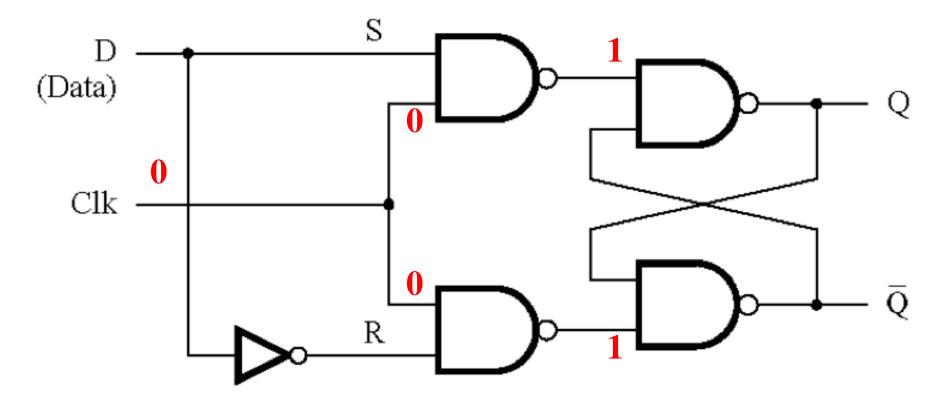
- Dealing with two inputs (S and R) could be messy.
   For example, we may have to reset the latch before some operations in order to store a specific value but the reset may not be necessary depending on the current state of the latch.
- Why not just have one input and call it D.
- The D latch can be constructed using a simple modification of the SR latch.



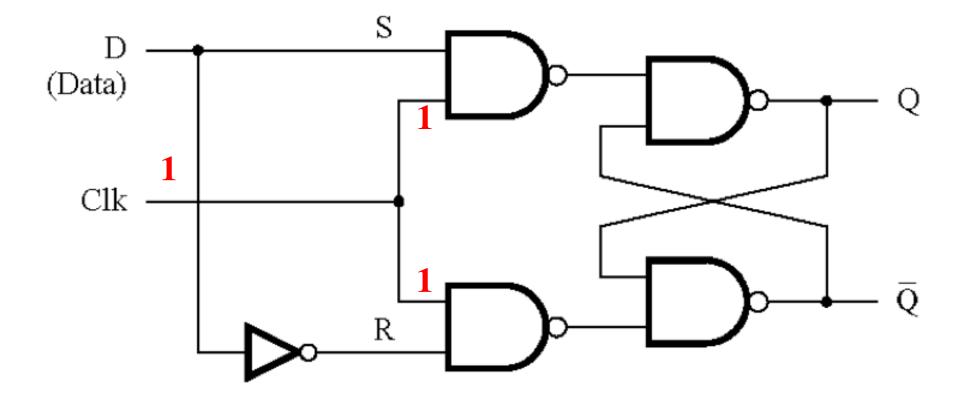


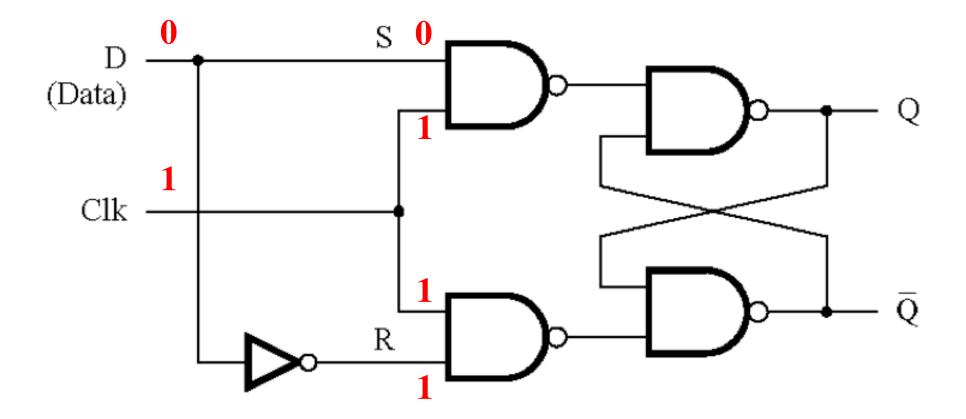
This is the only new thing here.

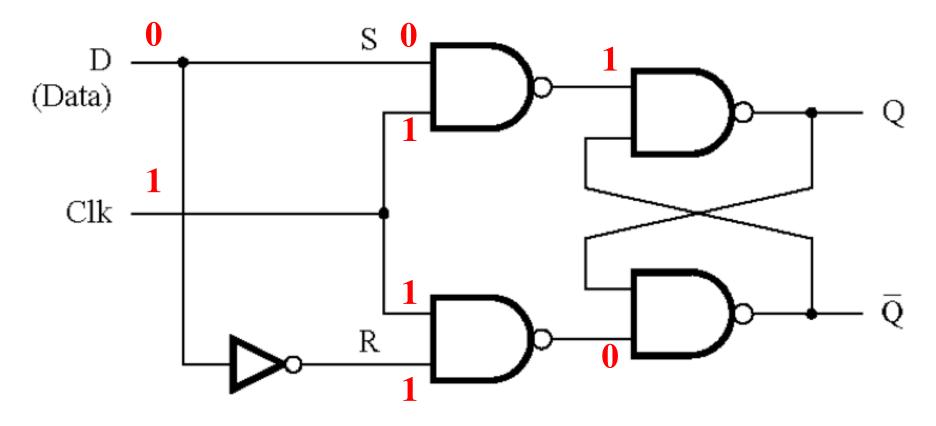




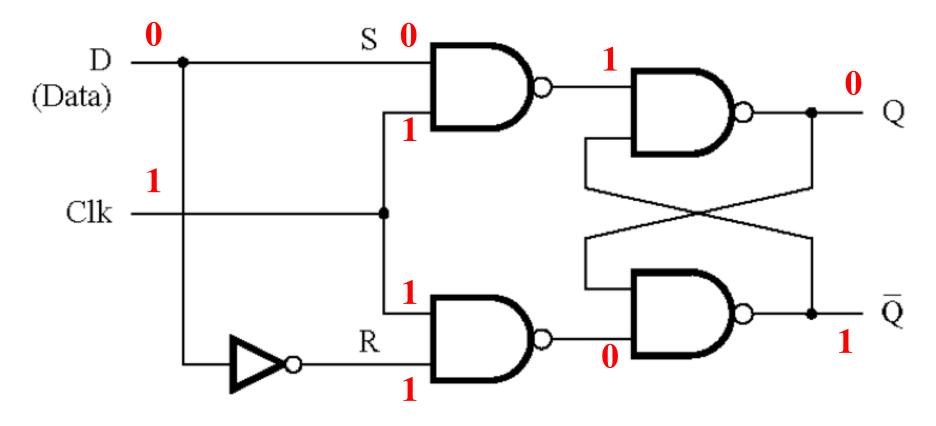
$\overline{S}$ $\overline{R}$	$Q_a Q_b$	
0 0	1 1	•
0 1	1 0	
1 0	0 1	
1 1	0/1 1/0	(no change)



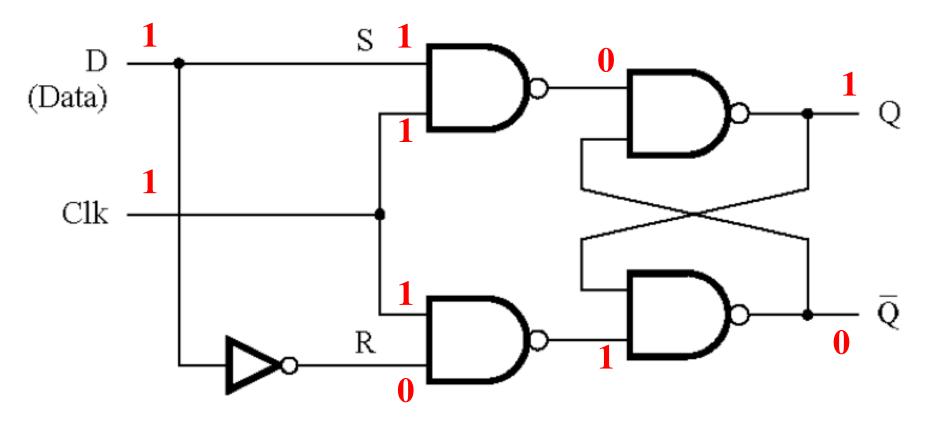




S R	$Q_a Q_b$	
0 0	1 1	•
0 1	1 0	
1 0	0 1	
1 1	0/1 1/0	(no change)

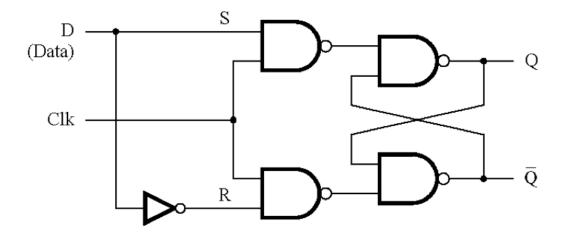


S R	$Q_a Q_b$	_
0 0	1 1	
0 1	1 0	
1 0	0 1	
1 1	0/1 1/0	(no change)



_	$Q_b$	$Q_a$	R	s	
	1	1	0	0	
	0	1	1	0	
	1	0	0	1	
(no change)	1/0	0/1	1	1	

# Circuit Diagram and Characteristic Table for the Gated D Latch



Clk
 D
 
$$Q(t+1)$$

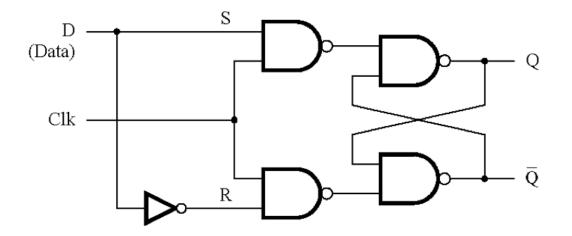
 0
 x
  $Q(t)$ 

 1
 0
 0

 1
 1
 1

Note that it is now impossible to have S=R=1.

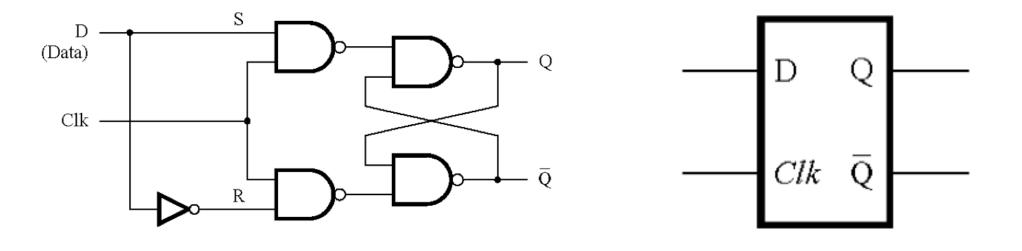
# Circuit Diagram and Characteristic Table for the Gated D Latch



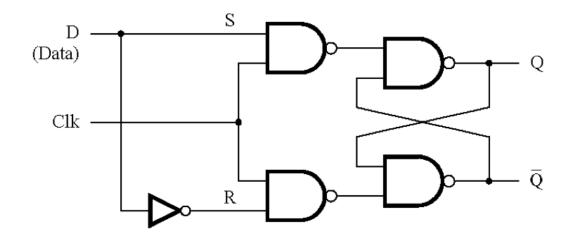
	- 1)
$ \begin{array}{c cccc} 0 & x & Q(t) \\ 1 & 0 & 0 \\ 1 & 1 & 1 \end{array} $	

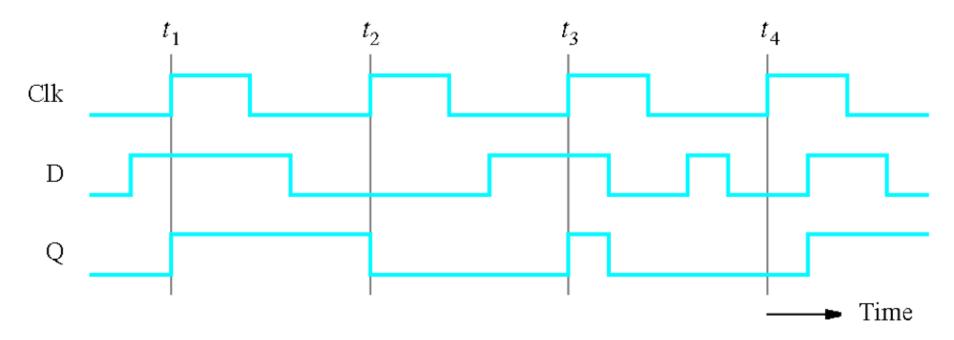
When Clk=1 the output follows the D input. When Clk=0 the output cannot be changed.

## Circuit Diagram and Graphical Symbol for the Gated D Latch



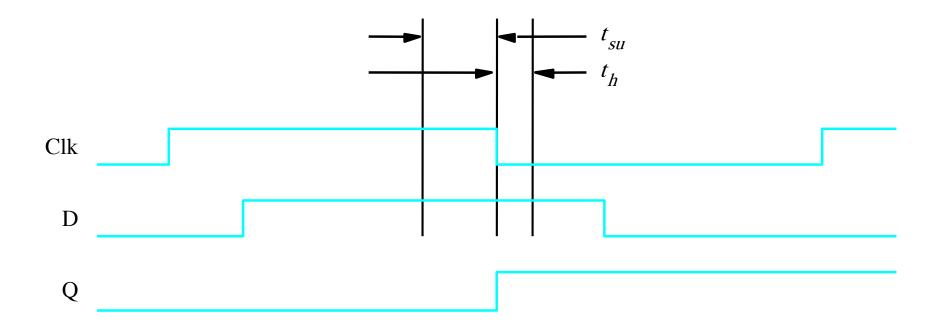
### Timing Diagram for the Gated D Latch





[ Figure 5.7d from the textbook ]

### Setup and hold times

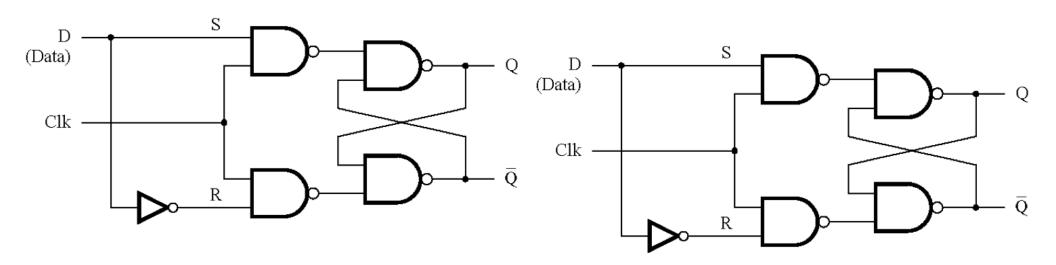


Setup time  $(t_{su})$  – the minimum time that the D signal must be stable prior to the negative edge of the Clock signal.

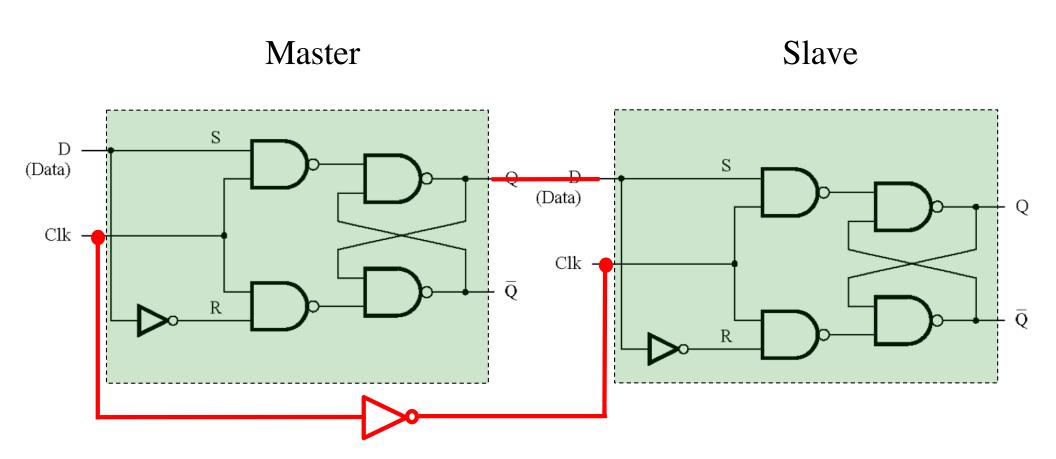
Hold time  $(t_h)$  – the minimum time that the D signal must remain stable after the negative edge of the Clock signal.

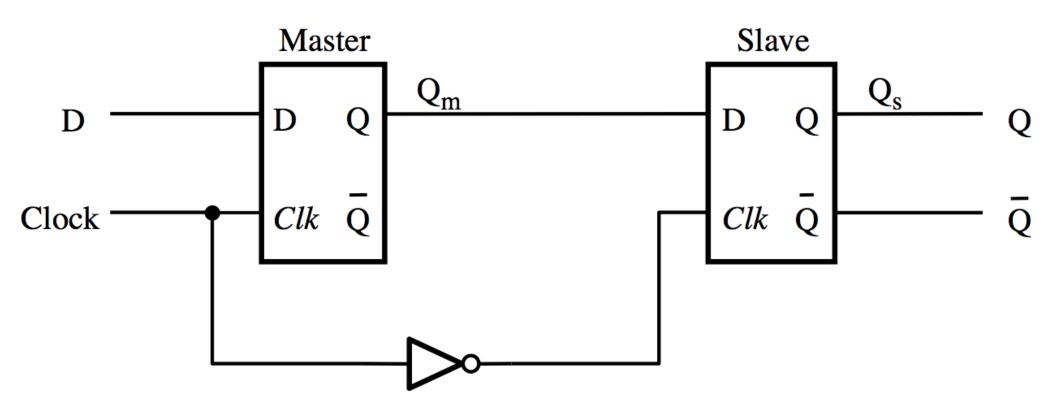
### **Master-Slave D Flip-Flop**

Master Slave



Master Slave (Data) (Data) Clk





## Constructing a Master-Slave D Flip-Flop From one D Latch and one Gated SR Latch

(This version uses one less NOT gate)

Slave

Master

## Constructing a Master-Slave D Flip-Flop From one D Latch and one Gated SR Latch

(This version uses one less NOT gate)

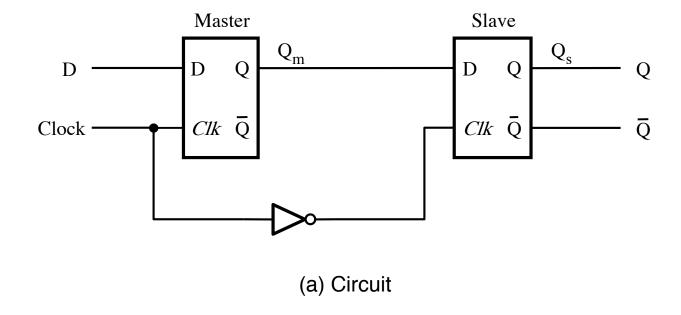
Master Slave (Data) Clk Clk

### **Edge-Triggered D Flip-Flops**

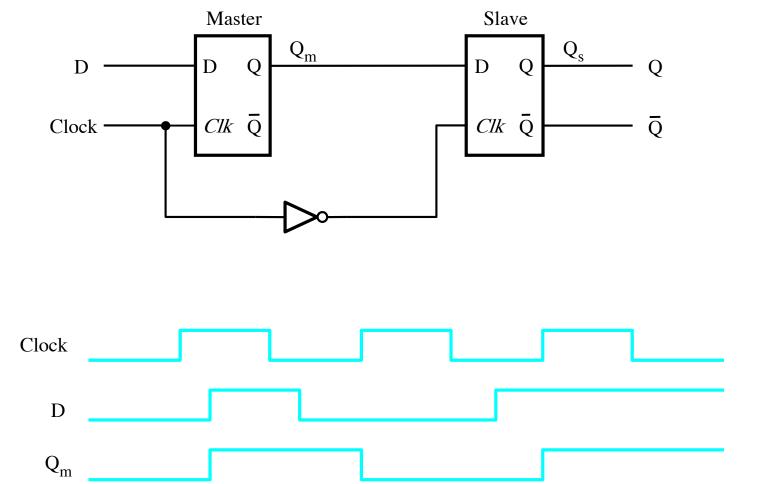
#### **Motivation**

In some cases we need to use a memory storage device that can change its state no more than once during each clock cycle.

### **Master-Slave D Flip-Flop**

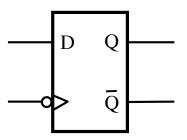


#### Timing Diagram for the Master-Slave D Flip-Flop

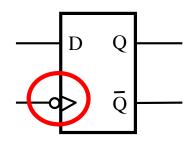


 $Q = Q_s$ 

### **Graphical Symbol for the Master-Slave D Flip-Flop**

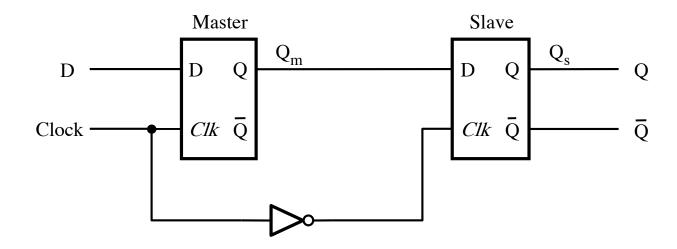


#### **Graphical Symbol for the Master-Slave D Flip-Flop**

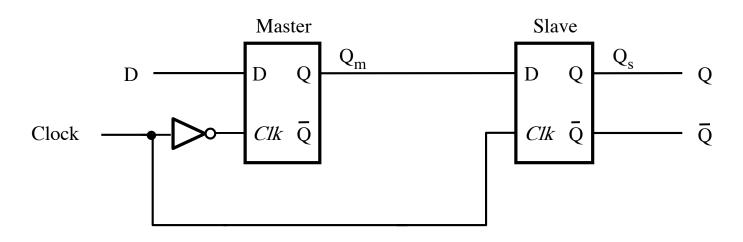


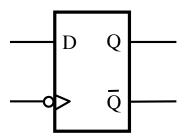
The > means that this is edge-triggered

The small circle means that is is the negative edge

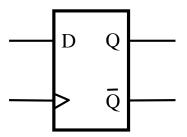


#### Positive-Edge-Triggered Master-Slave D Flip-Flop

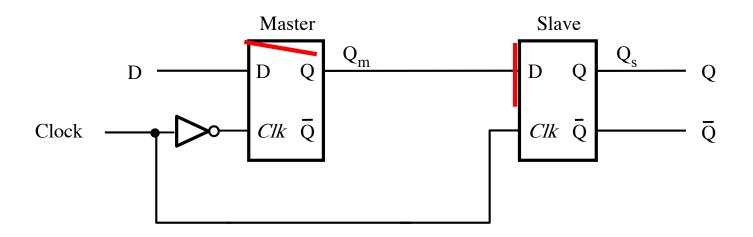




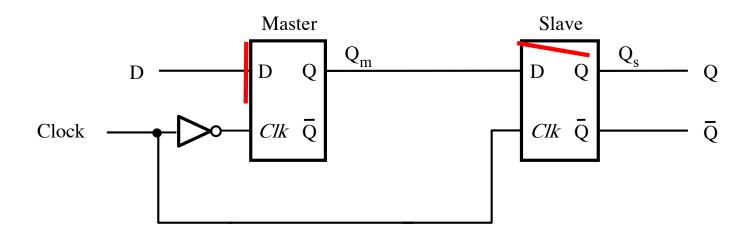
#### Positive-Edge-Triggered Master-Slave D Flip-Flop



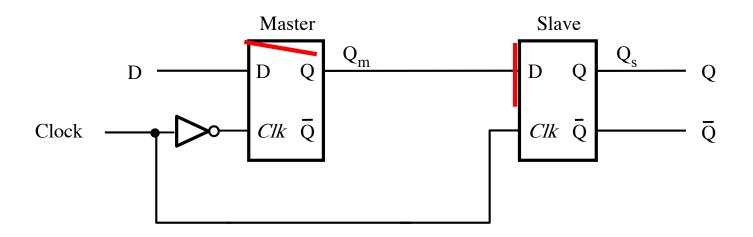
## D Flip-Flop: A Double Door Analogy

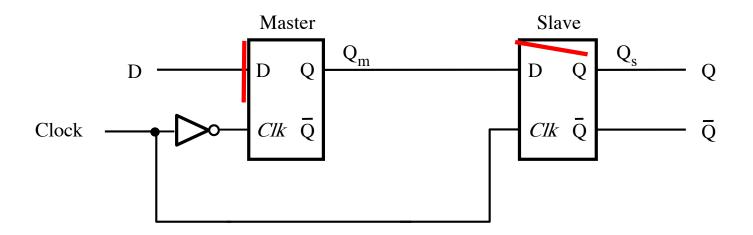


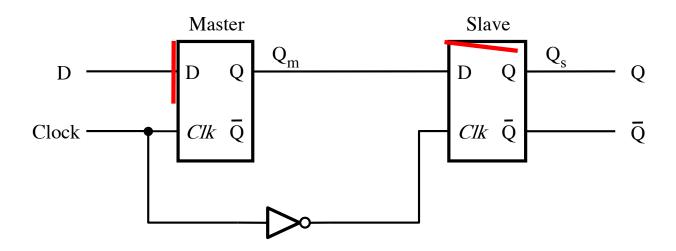




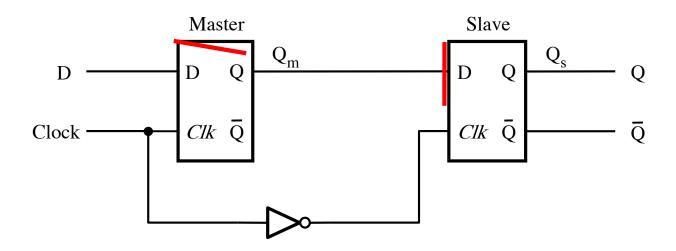


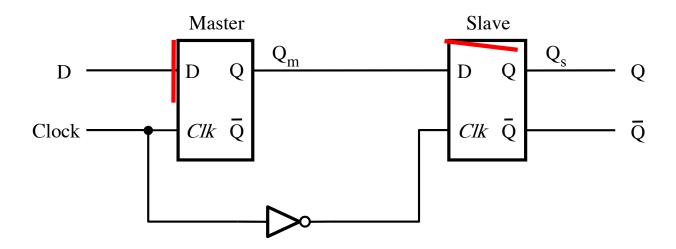


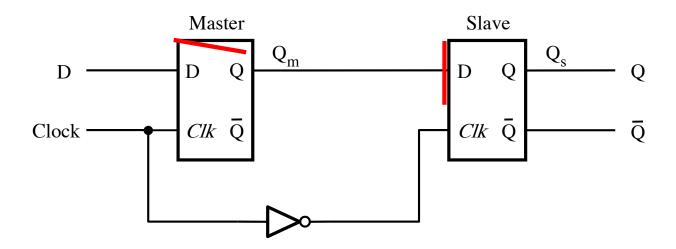




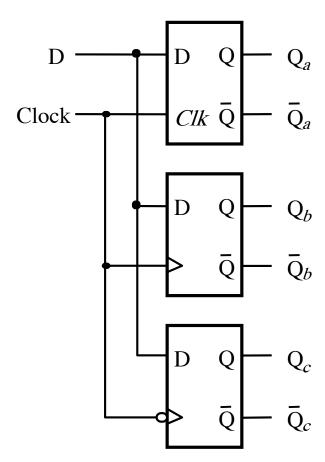




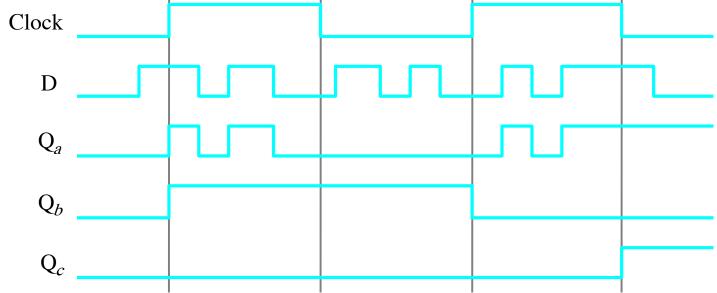


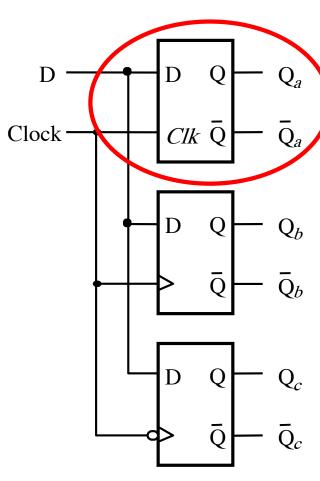


# Other Types of Edge-Triggered D Flip-Flops



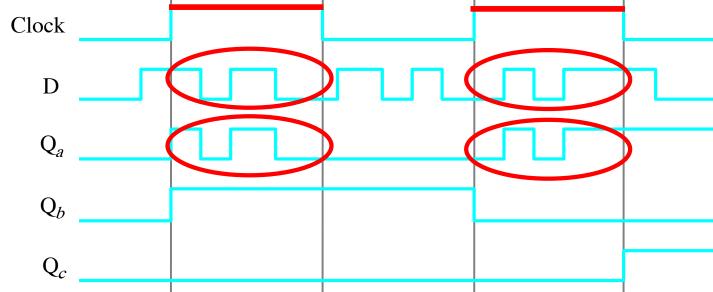
## Comparison of level-sensitive and edge-triggered D storage elements

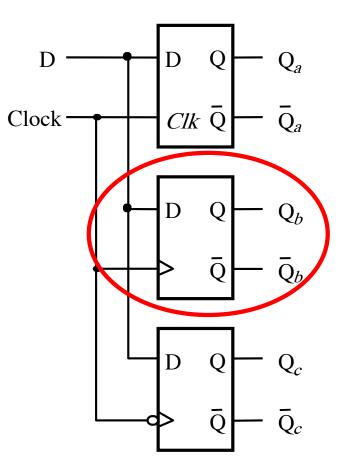




## Comparison of level-sensitive and edge-triggered D storage elements

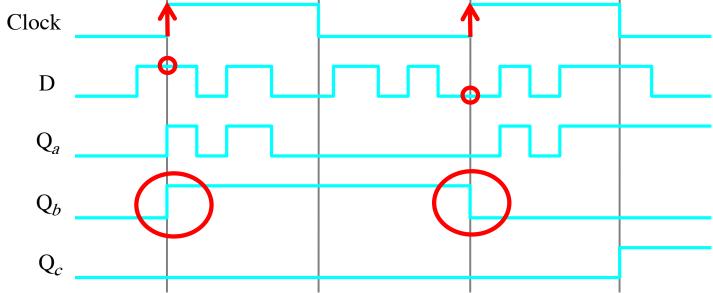
Level-sensitive (the output mirrors the D input when Clk=1)

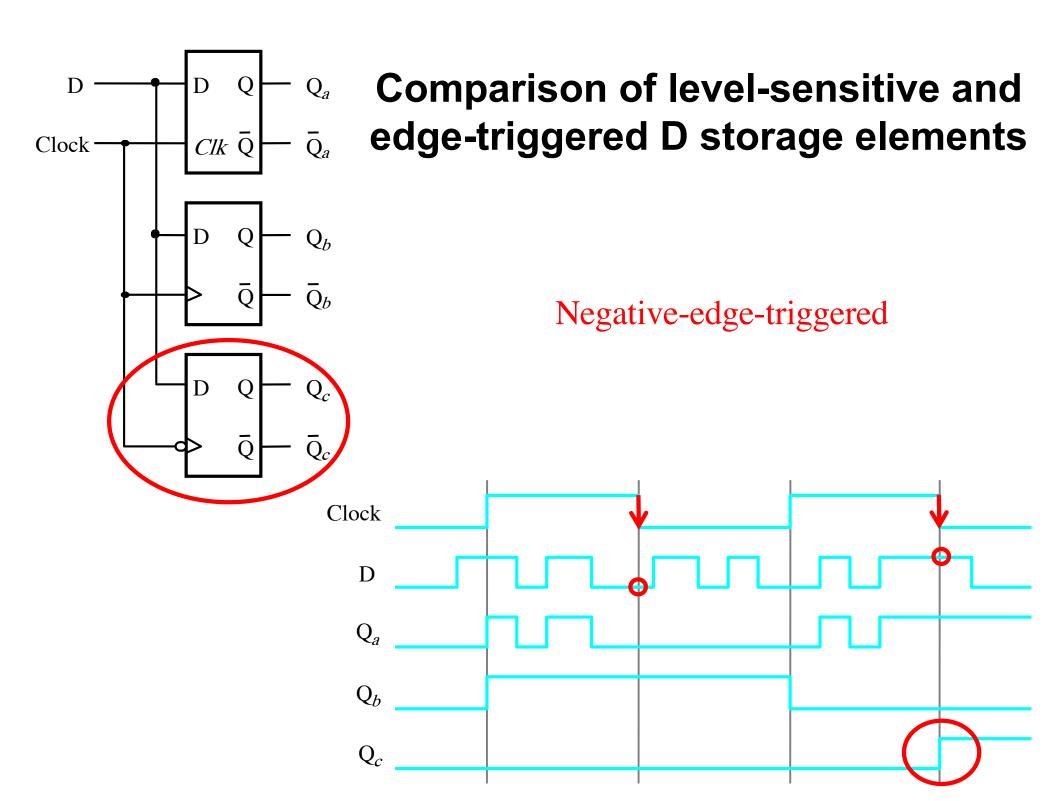




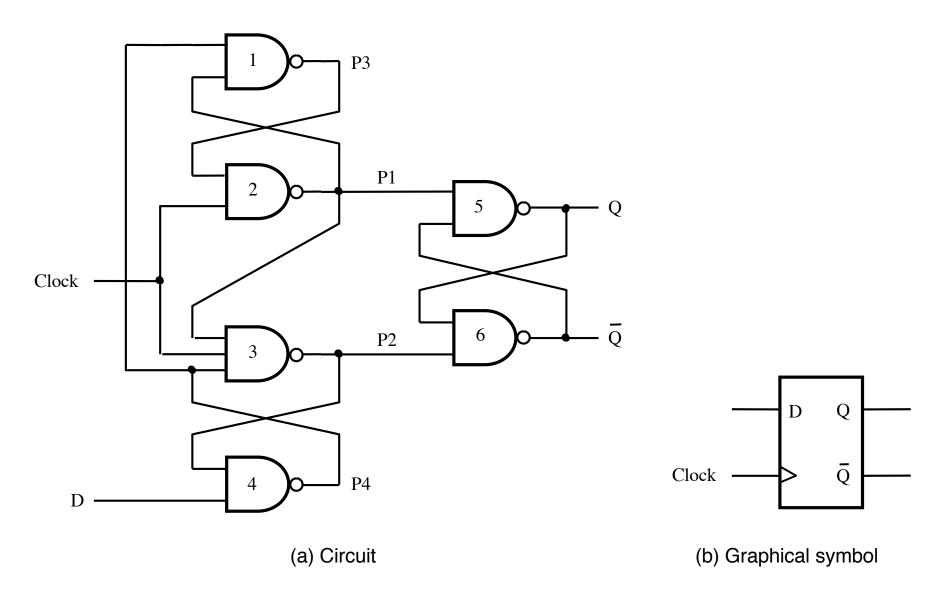
## Comparison of level-sensitive and edge-triggered D storage elements

Positive-edge-triggered



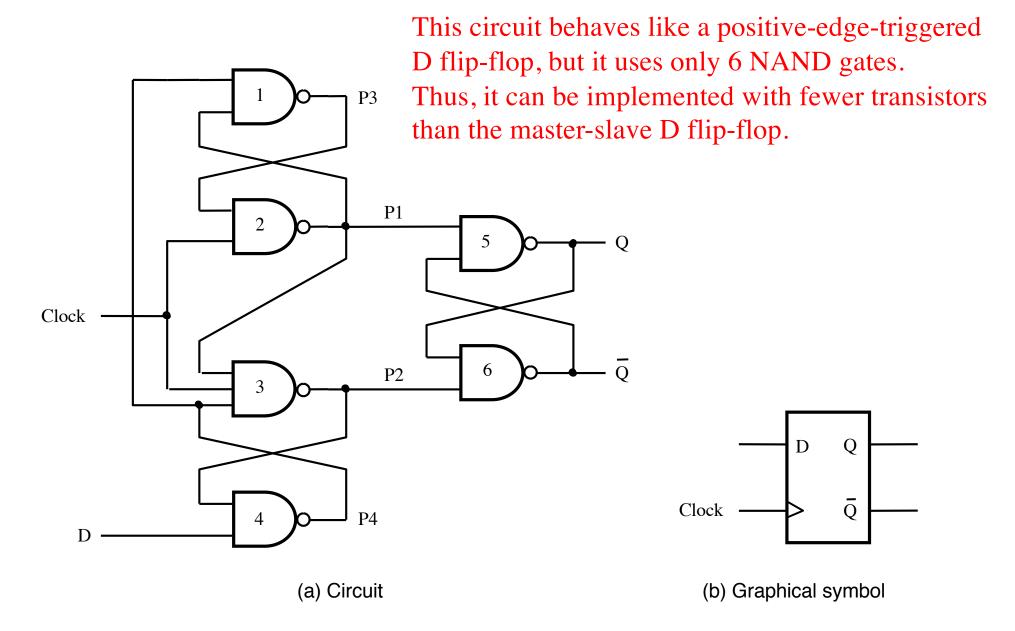


#### A positive-edge-triggered D flip-flop



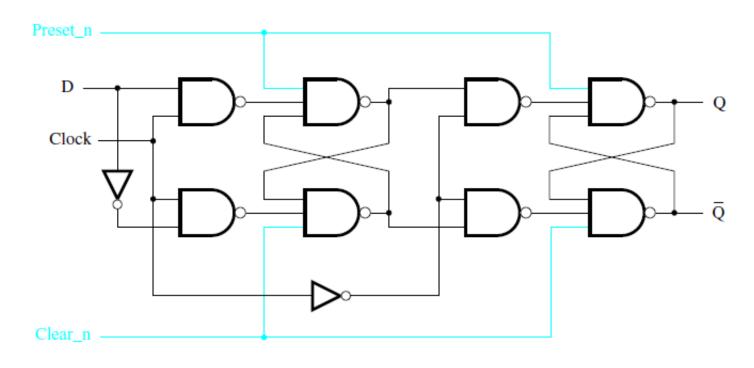
[ Figure 5.11 from the textbook ]

#### A positive-edge-triggered D flip-flop

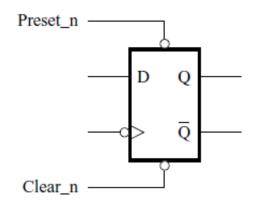


[ Figure 5.11 from the textbook ]

#### Master-slave D flip-flop with Clear and Preset



(a) Circuit

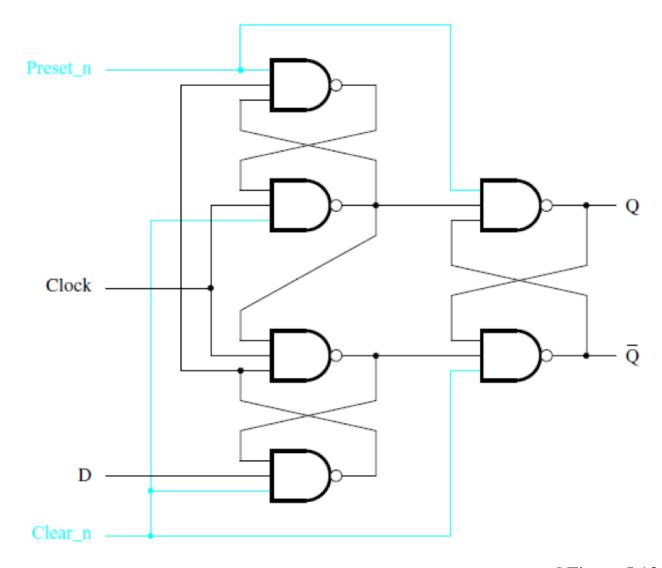


(b) Graphical symbol

[ Figure 5.12 from the textbook ]

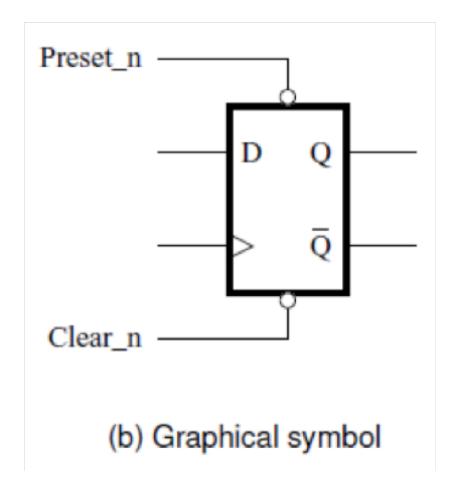
# Positive-edge-triggered D flip-flop with Clear and Preset

## Positive-edge-triggered D flip-flop with Clear and Preset

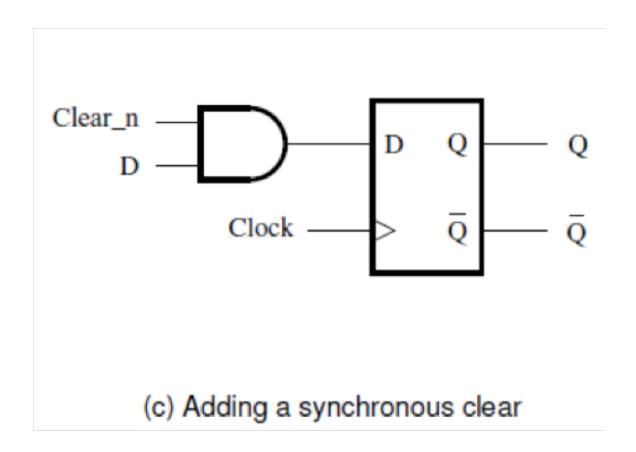


[ Figure 5.13a from the textbook ]

## Positive-edge-triggered D flip-flop with Clear and Preset

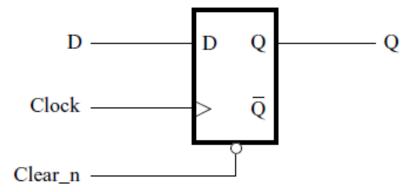


## Positive-edge-triggered D flip-flop with Synchronous Clear

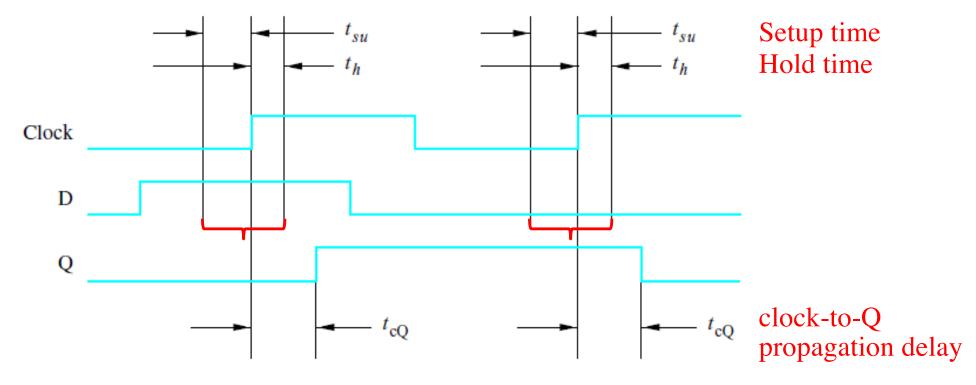


The output Q can be cleared only on the positive clock edge.

### **Flip-Flop Timing Parameters**



(a) D flip-flop with asynchronous clear



[ Figure 5.14 from the textbook ]

(b) Timing diagram

- Basic Latch is a feedback connection of two NOR gates or two NAND gates, which can store one bit of information. It can be set using the S input and reset to 0 using the R input.
- Gated Latch is a basic latch that includes input gating and a control input signal. The latch retains its existing state when the control input is equal to 0. Its state may be changed when the control signal is equal to 1.

- Two types of gated latches (the control input is the clock):
- Gated SR Latch uses the S and R inputs to set the latch to 1 or reset it to 0.
- Gated D Latch uses the D input to force the latch into a state that has the same logic value as the D input.

 Flip-Flop – is a storage element that can have its output state changed only on the edge of the controlling clock signal.

- Positive-edge triggered if the state changes when the clock signal goes from 0 to 1.
- Negative-edge triggered if the state changes when the clock signal goes from 1 to 0.

The word *latch* is mainly used for storage elements, while clocked devices are described as *flip-flops*.

A **latch** is level-sensitive, whereas a **flip-flop** is edge-sensitive. That is, when a latch is enabled it becomes transparent, while a flip flop's output only changes on a single type (positive going or negative going) of clock edge.

### **Questions?**

### THE END