

## CprE 281: Digital Logic

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http://www.ece.iastate.edu/~alexs/classes/

## Synchronous Sequential Circuits Basic Design Steps

CprE 281: Digital Logic

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## Administrative Stuff

- Homework 10 is due next Monday.
- Start thinking about your final project ideas.

First Design Pattern: Moore Machines

# Moore Machine: <br> A Type of Finite State Machine (FSM) 


[ Figure 6.3 from the textbook]


- Finite number of states (nodes).
- Discrete state transitions (edges).
- Only "in" one state at a time.
- One reset state
- Every state has an outgoing state transition for each possible input.
[ Figure 6.3 from the textbook]



## Key: The next state depends on both the current state and the current input.

[ Figure 6.3 from the textbook]


Key: The output depends only on the current state.
[ Figure 6.3 from the textbook]


| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |

## Let's do a simulation



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$

| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$

| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$

| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| :---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | :---: |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | :---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | $\square 1$ | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | :---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



In general, we need to start tracing from the beginning to know which state the FSM is in. It may not be clear from a short sequence of outputs.

Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$

| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |

## Inferring the States



Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$ $w: \begin{array}{llllllllllll}0 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1\end{array}$
$z: 0$

## Inferring the States



Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$ $w: \quad 0 \quad 1 \quad 0 \quad 1 \quad 0$
$z: 0$


Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$

| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |

## What is a State?

It is not really a memory of every past input (We might run out of space to remember it all!)

Rather, it is a characterization or snapshot of the pattern of inputs that have come before.

## Moore Machine Implementation

The state diagram is just an illustration to help us describe and reason about how the FSM will behave in each of its states.

So, how do we turn it into a circuit?

## Moore Machine Implementation



Note: The $W$ and $Z$ lines need not be wires. They can be buses.

## State Storage



Any usable "memory" of the preceding input sequence is encoded in the flip-flop array.

## FSM States

The Flip-Flop array stores an encoding of the current state.


## State Encoding

Each of the states in our design is identified by a distinct code.

If we use 3 flip-flops, then the FSM can have up to $2^{3}=8$ distinct states.

So, when the flip-flop array contains the code 011, we say that the machine is in state 011 .

## Synchronous Design



Every active clock edge causes a state transition.

## Synchronous Design



We expect the input signals to be stable before the active clock edge occurs.

## Synchronous Design



There is a whole other class of sequential circuits that are asynchronous, but we will not study them in this course.

## Sequential Circuits: Key Ideas

The current output depends on something about the preceding sequence of inputs (and maybe the current output).

Using memory elements (i.e., flip-flops), we design the circuit to remember some relevant information about the prior inputs.

## Example



We need to find both the next state logic and the output logic implied by this machine.
[ Figure 6.3 from the textbook]


| Present <br> state | Next state |  |
| :---: | :---: | :---: |
|  |  |  |
|  | $w=0 \quad w=1$ |  |
| A |  |  |
| B |  |  |
| C |  |  |



| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
| A | A | B | 0 |
| B | A | C | 0 |
| C | A | C | 1 |

[Figure 6.4 from the textbook ]

## How to represent the States?

One way is to encode each state with a 2-bit binary number

$$
\begin{aligned}
& A \sim 00 \\
& B \sim 01 \\
& C \sim 10
\end{aligned}
$$

## How to represent the states?

One way is to encode each state with a 2-bit binary number

$$
\begin{aligned}
& A \sim 00 \\
& B \sim 01 \\
& C \sim 10
\end{aligned}
$$

How many flip-flops do we need?

## Let's use two flip-flops to hold the machine's state




Let's pick D Flip-Flops.

Clock


We will call $y_{1}$ and $y_{2}$ the present state variables.
We will call $Y_{1}$ and $Y_{2}$ the next state variables.
[ Figure 6.5 from the textbook]


Two zeros on the output JOINTLY represent state A.


This flip-flop output pattern represents state B.


This flip-flop output pattern represents state C.


What does this flip-flop output pattern represent?

Clock

This would be state D, but we don't have one in this example. So this is an impossible state.


We will call $y_{1}$ and $y_{2}$ the present state variables.
We will call $Y_{1}$ and $Y_{2}$ the next state variables.
[ Figure 6.5 from the textbook ]


We will call $y_{1}$ and $y_{2}$ the present state variables.
We will call $Y_{1}$ and $Y_{2}$ the next state variables.
[ Figure 6.5 from the textbook ]


We need to find logic expressions for $Y_{1}\left(w, y_{1}, y_{2}\right), Y_{2}\left(w, y_{1}, y_{2}\right)$, and $z\left(y_{1}, y_{2}\right)$.
[ Figure 6.5 from the textbook ]


We need to find logic expressions for $Y_{1}\left(w, y_{1}, y_{2}\right), Y_{2}\left(w, y_{1}, y_{2}\right)$, and $z\left(y_{1}, y_{2}\right)$.
[ Figure 6.5 from the textbook ]

| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
| A | A | B | 0 |
| B | A | C | 0 |
| C | A | C | 1 |

Suppose we encoded our states in the same order in which they were labeled:

$$
\begin{aligned}
& A \sim 00 \\
& B \sim 01 \\
& C \sim 10
\end{aligned}
$$

| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
| A | A | B | 0 |
| B | A | C | 0 |
| C | A | C | 1 |



The finite state machine will never reach a state encoded as 11 .
[ Figure 6.6 from the textbook ]

| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
| A | A | B | 0 |
| B | A | C | 0 |
| C | A | C | 1 |


[ Figure 6.6 from the textbook ]

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ |  |
|  | $\|c\|$ <br> 00 | 00 |  |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 |  |  |
| 0 | 0 | 1 |  |  |
| 0 | 1 | 0 |  |  |
| 0 | 1 | 1 |  |  |
| 1 | 0 | 0 |  |  |
| 1 | 0 | 1 |  |  |
| 1 | 1 | 0 |  |  |
| 1 | 1 | 1 |  |  |


| $y_{2}$ | $y_{1}$ | $z$ |
| :--- | :--- | :--- |
| 0 | 0 |  |
| 0 | 1 |  |
| 1 | 0 |  |
| 1 | 1 |  |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ |  |
|  | $w=1$ |  |  |
| 00 | 00 | 01 | 0 |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 |  |  |
| 0 | 0 | 1 |  |  |
| 0 | 1 | 0 |  |  |
| 0 | 1 | 1 |  |  |
| 1 | 0 | 0 |  |  |
| 1 | 0 | 1 |  |  |
| 1 | 1 | 0 |  |  |
| 1 | 1 | 1 |  |  |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ |  |
|  | $\|c\|$ <br> 00 | 00 |  |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 |  |  |
| 0 | 0 | 1 |  |  |
| 0 | 1 | 0 |  |  |
| 0 | 1 | 1 |  |  |
| 1 | 0 | 0 |  |  |
| 1 | 0 | 1 |  |  |
| 1 | 1 | 0 |  |  |
| 1 | 1 | 1 |  |  |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  |  |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ | $z$ |
| 00 | 00 | 01 | 0 |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |  |
| 0 | 0 | 1 | 0 |  |
| 0 | 1 | 0 | 0 |  |
| 0 | 1 | 1 | d |  |
| 1 | 0 | 0 |  |  |
| 1 | 0 | 1 |  |  |
| 1 | 1 | 0 |  |  |
| 1 | 1 | 1 |  |  |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  | Output |
| :---: | :---: | ---: | :---: |
|  | $w=0$ | $w=1$ |  |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ |  |
| 00 | 00 | 01 | 0 |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |  |
| 0 | 0 | 1 | 0 |  |
| 0 | 1 | 0 | 0 |  |
| 0 | 1 | 1 | d |  |
| 1 | 0 | 0 | 0 |  |
| 1 | 0 | 1 | 1 |  |
| 1 | 1 | 0 | 1 |  |
| 1 | 1 | 1 | d |  |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  |  |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ |  |
| 00 | 00 | 01 | 0 |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook ]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |  |
| 0 | 0 | 1 | 0 |  |
| 0 | 1 | 0 | 0 |  |
| 0 | 1 | 1 | d |  |
| 1 | 0 | 0 | 0 |  |
| 1 | 0 | 1 | 1 |  |
| 1 | 1 | 0 | 1 |  |
| 1 | 1 | 1 | d |  |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  |  |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ | $z$ |
| 00 | 00 | 01 | 0 |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | d | d |
| 1 | 0 | 0 | 0 |  |
| 1 | 0 | 1 | 1 |  |
| 1 | 1 | 0 | 1 |  |
| 1 | 1 | 1 | d |  |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  | Output |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
| $y_{2}{ }_{1}$ | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ |  |
| 00 |  | 01 | 0 |
| 01 |  | 10 | 0 |
| 10 |  | 10 | 1 |
| 11 |  | dd | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | d | d |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | d | d |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

$$
Q(t)=y_{2} y_{l} \text { and } Q(t+l)=Y_{2} Y_{1}
$$

| Present <br> state | Next state |  |  |
| :---: | :---: | :---: | :---: |
|  | $w=0$ | $w=1$ |  |
|  | $Y_{2} Y_{1}$ | $Y_{2} Y_{1}$ |  |
| 00 | 00 | 01 | 0 |
| 01 | 00 | 10 | 0 |
| 10 | 00 | 10 | 1 |
| 11 | $d d$ | $d d$ | $d$ |

[ Figure 6.6 from the textbook]

| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | d | d |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | d | d |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

Note that the textbook draws these K-Maps differently from all previous K-maps
(the most significant bit indexes the rows).

$Q(t)=y_{2} y_{1}$ and $Q(t+1)=Y_{2} Y_{1}$


| $w$ | $y_{2}$ | $y_{1}$ | $Y_{2}$ | $Y_{1}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | d | d |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | d | d |


| $y_{2}$ | $y_{1}$ | $z$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | d |

## Don't care conditions simplify the combinatorial logic



Ignoring don't cares

$$
Y_{1}=w \bar{y}_{1} \bar{y}_{2}
$$

$$
Y_{2}=w y_{1} \bar{y}_{2}+\bar{w}_{1} y_{2}
$$

$$
z=\bar{y}_{1} y_{2}
$$

Using don't cares
$Y_{1}=w \bar{y}_{1} \bar{y}_{2}$
$\begin{aligned} Y_{2} & =w y_{1}+w y_{2} \\ & =w\left(y_{1}+y_{2}\right)\end{aligned}$
$z=y_{2}$
[ Figure 6.7 from the textbook ]

[ Figure 6.8 from the textbook]

[ Figure 6.8 from the textbook]

State $A=00$

Finally, we add a reset signal. When it is equal to zero it puts the machine back to its start state, which is state 00 in this case.

[ Figure 6.8 from the textbook ]

State $A=00$

Finally, we add a reset signal. When it is equal to zero it puts the machine back to its start state, which is state 00 in this case.

[ Figure 6.8 from the textbook]

State $A=00$


State $A=00$


State $A=00$


State $A=00$


State $A=00$


State $B=01$


State $B=01$


State $B=01$


State $A=00$


State $A=00$


State $A=00$


State $B=01$


State $B=01$


State $B=01$


State $C=10$


State $C=10$


State $C=10$


State $A=00$


| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |


[ Figure 6.9 from the textbook]

Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$ $w: \begin{array}{llllllllllll}0 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1\end{array}$ $z: 0$


Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$ $w: ~ 0 \begin{array}{lllllllllll} & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1\end{array}$ $z: 0$


| Clockcycle: | $\mathrm{t}_{0}$ | $\mathrm{t}_{1}$ | $\mathrm{t}_{2}$ | $\mathrm{t}_{3}$ | $\mathrm{t}_{4}$ | $\mathrm{t}_{5}$ | $\mathrm{t}_{6}$ | $\mathrm{t}_{7}$ | $\mathrm{t}_{8}$ | $\mathrm{t}_{9}$ | $\mathrm{t}_{10}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



Clockcycle: $\begin{array}{llllllllllll}\mathrm{t}_{0} & \mathrm{t}_{1} & \mathrm{t}_{2} & \mathrm{t}_{3} & \mathrm{t}_{4} & \mathrm{t}_{5} & \mathrm{t}_{6} & \mathrm{t}_{7} & \mathrm{t}_{8} & \mathrm{t}_{9} & \mathrm{t}_{10}\end{array}$ | $w:$ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $z:$ | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |



## Summary: Designing a Moore Machine

- Obtain the circuit specification.
- Derive a state diagram.
- Derive the state table.
- Decide on a state encoding.
- Encode the state table.
- Derive the output logic and next-state logic.
- Add a reset signal.


## Questions?

## THE END

