

# CprE 281: Digital Logic

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**<http://www.ece.iastate.edu/~alexs/classes/>**

# Design Examples

*CprE 281: Digital Logic  
Iowa State University, Ames, IA  
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# **Administrative Stuff**

- HW3 is due on Monday Sep 16 @ 10pm
- Please write clearly on the first page the following three things:
  - Your First and Last Name
  - Your Student ID Number
  - Your Lab Section Letter
- Submit on Canvas as \*one\* PDF file.
- Please orient your pages such that the text can be read without the need to rotate the page.

# **Quick Review**

# Axioms of Boolean Algebra

1a.  $0 \cdot 0 = 0$

1b.  $1 + 1 = 1$

2a.  $1 \cdot 1 = 1$

2b.  $0 + 0 = 0$

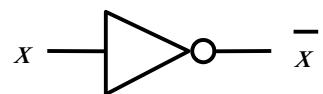
3a.  $0 \cdot 1 = 1 \cdot 0 = 0$

3b.  $1 + 0 = 0 + 1 = 1$

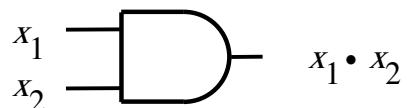
4a. If  $x=0$ , then  $\bar{x} = 1$

4b. If  $x=1$ , then  $\bar{x} = 0$

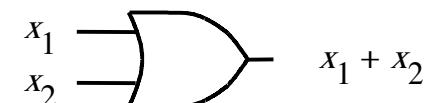
# The Three Basic Logic Gates



NOT gate



AND gate



OR gate

[ Figure 2.8 from the textbook ]

# Single-Variable Theorems

$$5a. \quad x \cdot 0 = 0$$

$$5b. \quad x + 1 = 1$$

$$6a. \quad x \cdot 1 = x$$

$$6b. \quad x + 0 = x$$

$$7a. \quad x \cdot x = x$$

$$7b. \quad x + x = x$$

$$8a. \quad x \cdot \bar{x} = 0$$

$$8b. \quad x + \bar{x} = 1$$

$$9. \quad \bar{\bar{x}} = x$$

# Two- and Three-Variable Properties

10a.  $x \cdot y = y \cdot x$  Commutative

10b.  $x + y = y + x$

11a.  $x \cdot (y \cdot z) = (x \cdot y) \cdot z$  Associative

11b.  $x + (y + z) = (x + y) + z$

12a.  $x \cdot (y + z) = x \cdot y + x \cdot z$  Distributive

12b.  $x + y \cdot z = (x + y) \cdot (x + z)$

13a.  $x + x \cdot y = x$  Absorption

13b.  $x \cdot (x + y) = x$

# Two- and Three-Variable Properties

$$14a. \quad x \cdot y + x \cdot \bar{y} = x \quad \text{Combining}$$

$$14b. \quad (x + y) \cdot (x + \bar{y}) = x$$

$$15a. \quad \overline{x \cdot y} = \bar{x} + \bar{y} \quad \text{DeMorgan's}$$

$$15b. \quad \overline{x + y} = \bar{x} \cdot \bar{y} \quad \text{theorem}$$

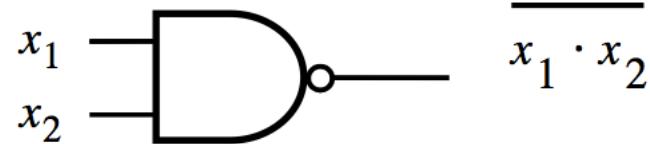
$$16a. \quad x + \bar{x} \cdot y = x + y$$

$$16b. \quad x \cdot (\bar{x} + y) = x \cdot y$$

$$17a. \quad x \cdot y + y \cdot z + \bar{x} \cdot z = x \cdot y + \bar{x} \cdot z \quad \text{Consensus}$$

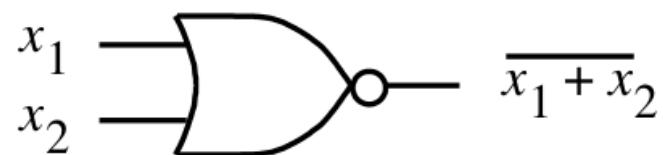
$$17b. \quad (x+y) \cdot (y+z) \cdot (\bar{x}+z) = (x+y) \cdot (\bar{x}+z)$$

# NAND Gate



$x_1$	$x_2$	f
0	0	1
0	1	1
1	0	1
1	1	0

# NOR Gate



$x_1$	$x_2$	f
0	0	1
0	1	0
1	0	0
1	1	0

# **Why do we need two more gates?**

**They can be implemented with fewer transistors.**

**Each of the new gates can be used to implement  
the three basic logic gates: NOT, AND, OR.**

# **Implications**

**Any Boolean function can be implemented  
with only NAND gates!**

# **Implications**

**Any Boolean function can be implemented  
with only NAND gates!**

**The same is also true for NOR gates!**

**minterms  
(for two variables)**

# The Four minterms

x	y	$m_0$
0	0	1
0	1	0
1	0	0
1	1	0

$m_0(x, y)$

x	y	$m_1$
0	0	0
0	1	1
1	0	0
1	1	0

$m_1(x, y)$

x	y	$m_2$
0	0	0
0	1	0
1	0	1
1	1	0

$m_2(x, y)$

x	y	$m_3$
0	0	0
0	1	0
1	0	0
1	1	1

$m_3(x, y)$

# The Four minterms

x	y	$m_0$
0	0	1
0	1	0
1	0	0
1	1	0

$m_0(x, y)$

x	y	$m_1$
0	0	0
0	1	1
1	0	0
1	1	0

$m_1(x, y)$

x	y	$m_2$
0	0	0
0	1	0
1	0	1
1	1	0

$m_2(x, y)$

x	y	$m_3$
0	0	0
0	1	0
1	0	0
1	1	1

$m_3(x, y)$

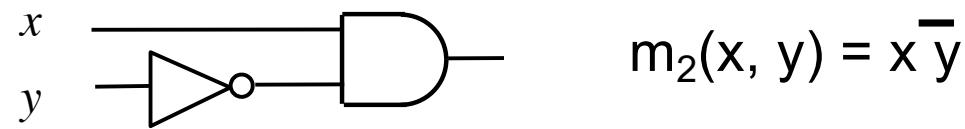
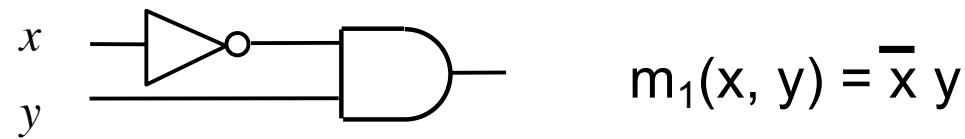
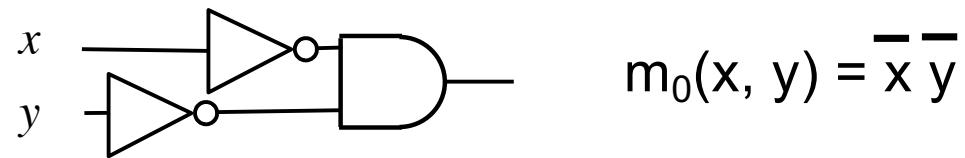
# The Four minterms

x	y		$m_0(x, y)$	$m_1(x, y)$	$m_2(x, y)$	$m_3(x, y)$
0	0		1	0	0	0
0	1		0	1	0	0
1	0		0	0	1	0
1	1		0	0	0	1

# The Four minterms

x	y	$\bar{x}\bar{y}$	$\bar{x}y$	$x\bar{y}$	$xy$
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

# Circuits for the four minterms



# **Maxterms (for two variables)**

# The Four Maxterms

<b>x</b>	<b>y</b>	<b>M<sub>0</sub></b>
0	0	0
0	1	1
1	0	1
1	1	1

$M_0(x, y)$

<b>x</b>	<b>y</b>	<b>M<sub>1</sub></b>
0	0	1
0	1	0
1	0	1
1	1	1

$M_1(x, y)$

<b>x</b>	<b>y</b>	<b>M<sub>2</sub></b>
0	0	1
0	1	1
1	0	0
1	1	1

$M_2(x, y)$

<b>x</b>	<b>y</b>	<b>M<sub>3</sub></b>
0	0	1
0	1	1
1	0	1
1	1	0

$M_3(x, y)$

# The Four Maxterms

x	y	M <sub>0</sub>
0	0	0
0	1	1
1	0	1
1	1	1

M<sub>0</sub>(x, y)

x	y	M <sub>1</sub>
0	0	1
0	1	0
1	0	1
1	1	1

M<sub>1</sub>(x, y)

x	y	M <sub>2</sub>
0	0	1
0	1	1
1	0	0
1	1	1

M<sub>2</sub>(x, y)

x	y	M <sub>3</sub>
0	0	1
0	1	1
1	0	1
1	1	0

M<sub>3</sub>(x, y)

# The Four Maxterms

x	y		$M_0(x, y)$	$M_1(x, y)$	$M_2(x, y)$	$M_3(x, y)$
0	0		0	1	1	1
0	1		1	0	1	1
1	0		1	1	0	1
1	1		1	1	1	0

# The Four Maxterms

x	y	$x + y$	$x + \bar{y}$	$\bar{x} + y$	$\bar{x} + \bar{y}$
0	0	0	1	1	1
0	1	1	0	1	1
1	0	1	1	0	1
1	1	1	1	1	0

# **minterms and Maxterms (for two variables)**

# minterms and Maxterms

x	y	$m_0$
0	0	1
0	1	0
1	0	0
1	1	0

x	y	$m_1$
0	0	0
0	1	1
1	0	0
1	1	0

x	y	$m_2$
0	0	0
0	1	0
1	0	1
1	1	0

x	y	$m_3$
0	0	0
0	1	0
1	0	0
1	1	1

x	y	$M_0$
0	0	0
0	1	1
1	0	1
1	1	1

x	y	$M_1$
0	0	1
0	1	0
1	0	1
1	1	1

x	y	$M_2$
0	0	1
0	1	1
1	0	0
1	1	1

x	y	$M_3$
0	0	1
0	1	1
1	0	1
1	1	0

# minterms and Maxterms

$$m_0(x, y) = \overline{x} \overline{y}$$

$$M_0(x, y) = x + y$$

$$m_1(x, y) = \overline{x} y$$

$$M_1(x, y) = x + \overline{y}$$

$$m_2(x, y) = x \overline{y}$$

$$M_2(x, y) = \overline{x} + y$$

$$m_3(x, y) = x y$$

$$M_3(x, y) = \overline{x} + \overline{y}$$

**minterms  
(for three variables)**

# The Eight minterms

# The Eight minterms

# Expressions for the minterms

$$m_0 = \overline{x} \ \overline{y} \ \overline{z}$$

$$m_1 = \overline{x} \ \overline{y} \ z$$

$$m_2 = \overline{x} \ y \ \overline{z}$$

$$m_3 = \overline{x} \ y \ z$$

$$m_4 = x \ \overline{y} \ \overline{z}$$

$$m_5 = x \ \overline{y} \ z$$

$$m_6 = x \ y \ \overline{z}$$

$$m_7 = x \ y \ z$$

# Expressions for the minterms

$$0 \ 0 \ 0 \quad m_0 = \overline{x} \ \overline{y} \ \overline{z}$$

$$0 \ 0 \ 1 \quad m_1 = \overline{x} \ \overline{y} \ z$$

$$0 \ 1 \ 0 \quad m_2 = \overline{x} \ y \ \overline{z}$$

$$0 \ 1 \ 1 \quad m_3 = \overline{x} \ y \ z$$

$$1 \ 0 \ 0 \quad m_4 = x \ \overline{y} \ \overline{z}$$

$$1 \ 0 \ 1 \quad m_5 = x \ \overline{y} \ z$$

$$1 \ 1 \ 0 \quad m_6 = x \ y \ \overline{z}$$

$$1 \ 1 \ 1 \quad m_7 = x \ y \ z$$

The bars coincide  
with the 0's  
in the binary expansion  
of the minterm sub-index

# **Maxterms (for three variables)**

# The Eight Maxterms

# The Eight Maxterms

# Expressions for the Maxterms

$$M_0 = x + y + z$$

$$M_1 = x + y + \bar{z}$$

$$M_2 = x + \bar{y} + z$$

$$M_3 = x + \bar{y} + \bar{z}$$

$$M_4 = \bar{x} + y + z$$

$$M_5 = \bar{x} + y + \bar{z}$$

$$M_6 = \bar{x} + \bar{y} + z$$

$$M_7 = \bar{x} + \bar{y} + \bar{z}$$

# Expressions for the Maxterms

$$0 \ 0 \ 0 \quad M_0 = x + y + z$$

$$0 \ 0 \ 1 \quad M_1 = x + y + \bar{z}$$

$$0 \ 1 \ 0 \quad M_2 = x + \bar{y} + z$$

$$0 \ 1 \ 1 \quad M_3 = x + \bar{y} + \bar{z}$$

$$1 \ 0 \ 0 \quad M_4 = \bar{x} + y + z$$

$$1 \ 0 \ 1 \quad M_5 = \bar{x} + y + \bar{z}$$

$$1 \ 1 \ 0 \quad M_6 = \bar{x} + \bar{y} + z$$

$$1 \ 1 \ 1 \quad M_7 = \bar{x} + \bar{y} + \bar{z}$$

The bars coincide  
with the 1's  
in the binary expansion  
of the maxterm sub-index

# **minterms and Maxterms (for three variables)**

# minterms and Maxterms

$$m_0 = \overline{x} \ \overline{y} \ \overline{z}$$

$$m_1 = \overline{x} \ \overline{y} \ z$$

$$m_2 = \overline{x} \ y \ \overline{z}$$

$$m_3 = \overline{x} \ y \ z$$

$$m_4 = x \ \overline{y} \ \overline{z}$$

$$m_5 = x \ \overline{y} \ z$$

$$m_6 = x \ y \ \overline{z}$$

$$m_7 = x \ y \ z$$

$$M_0 = x + y + z$$

$$M_1 = x + y + \overline{z}$$

$$M_2 = x + \overline{y} + z$$

$$M_3 = x + \overline{y} + \overline{z}$$

$$M_4 = \overline{x} + y + z$$

$$M_5 = \overline{x} + y + \overline{z}$$

$$M_6 = \overline{x} + \overline{y} + z$$

$$M_7 = \overline{x} + \overline{y} + \overline{z}$$

# **Synthesis Example**

# Truth table for a three-way light control

$x_1$	$x_2$	$x_3$	$f$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

[ Figure 2.31 from the textbook ]

# Let's Derive the SOP form

$x_1$	$x_2$	$x_3$	$f$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

# minterms and Maxterms (with three variables)

Row number	$x_1$	$x_2$	$x_3$	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1 \bar{x}_2 \bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1 \bar{x}_2 x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1 x_2 \bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1 x_2 x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1 \bar{x}_2 \bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1 \bar{x}_2 x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1 x_2 \bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1 x_2 x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

[ Figure 2.22 from the textbook ]

# Let's Derive the SOP form

$x_1$	$x_2$	$x_3$	$f$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

# Let's Derive the SOP form

$x_1$	$x_2$	$x_3$	$f$	
0	0	0	0	
0	0	1	1	$x_1 \ x_2 \ x_3$
0	1	0	1	$x_1 \ x_2 \ x_3$
0	1	1	0	
1	0	0	1	$x_1 \ x_2 \ x_3$
1	0	1	0	
1	1	0	0	
1	1	1	1	$x_1 \ x_2 \ x_3$

# Let's Derive the SOP form

$x_1$	$x_2$	$x_3$	$f$	
0	0	0	0	
0	0	1	1	$\overline{x}_1 \overline{x}_2 x_3$
0	1	0	1	$\overline{x}_1 x_2 \overline{x}_3$
0	1	1	0	
1	0	0	1	$x_1 \overline{x}_2 \overline{x}_3$
1	0	1	0	
1	1	0	0	
1	1	1	1	$x_1 x_2 x_3$

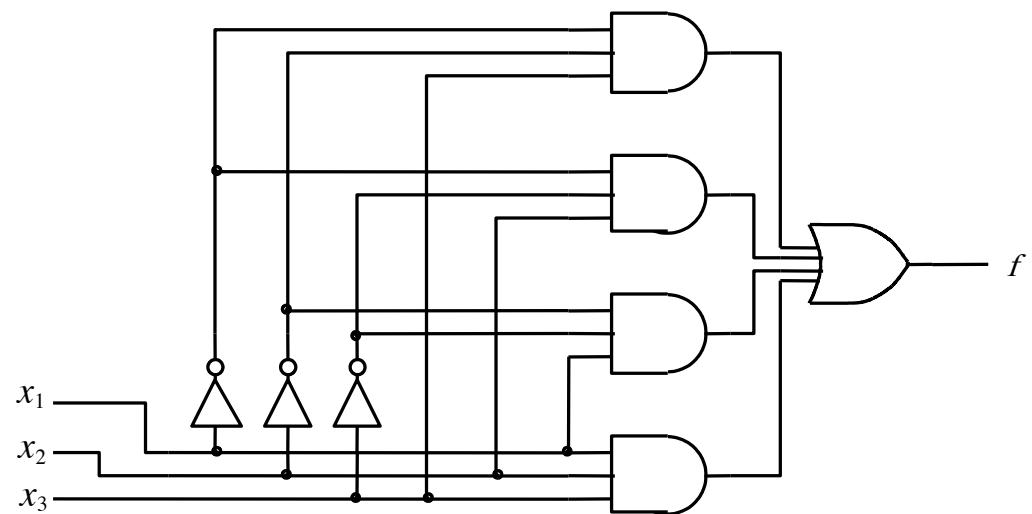
# Let's Derive the SOP form

$x_1$	$x_2$	$x_3$	$f$	
0	0	0	0	
0	0	1	1	$\bar{x}_1 \bar{x}_2 x_3$
0	1	0	1	$\bar{x}_1 x_2 \bar{x}_3$
0	1	1	0	
1	0	0	1	$x_1 \bar{x}_2 \bar{x}_3$
1	0	1	0	
1	1	0	0	
1	1	1	1	$x_1 x_2 x_3$

$$f = m_1 + m_2 + m_4 + m_7$$

$$= \bar{x}_1 \bar{x}_2 x_3 + \bar{x}_1 x_2 \bar{x}_3 + x_1 \bar{x}_2 \bar{x}_3 + x_1 x_2 x_3$$

# Sum-of-products realization



[ Figure 2.32a from the textbook ]

# Let's Derive the POS form

$x_1$	$x_2$	$x_3$	$f$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

[ Figure 2.31 from the textbook ]

# Let's Derive the POS form

$x_1$	$x_2$	$x_3$	$f$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

# minterms and Maxterms (with three variables)

Row number	$x_1$	$x_2$	$x_3$	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1 \bar{x}_2 \bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1 \bar{x}_2 x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1 x_2 \bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1 x_2 x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1 \bar{x}_2 \bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1 \bar{x}_2 x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1 x_2 \bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1 x_2 x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

[ Figure 2.22 from the textbook ]

# Let's Derive the POS form

$x_1$	$x_2$	$x_3$	$f$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

# Let's Derive the POS form

$x_1$	$x_2$	$x_3$	$f$	
0	0	0	0	$(x_1 + x_2 + x_3)$
0	0	1	1	
0	1	0	1	
0	1	1	0	$(x_1 + x_2 + x_3)$
1	0	0	1	
1	0	1	0	$(x_1 + x_2 + x_3)$
1	1	0	0	$(x_1 + x_2 + x_3)$
1	1	1	1	

# Let's Derive the POS form

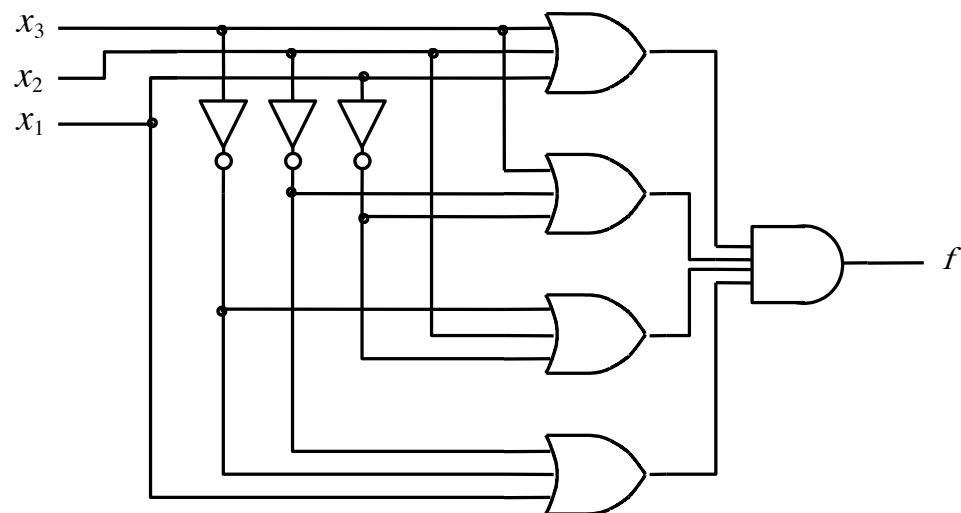
$x_1$	$x_2$	$x_3$	$f$	
0	0	0	0	$(x_1 + x_2 + x_3)$
0	0	1	1	
0	1	0	1	
0	1	1	0	$(x_1 + \bar{x}_2 + \bar{x}_3)$
1	0	0	1	
1	0	1	0	$(\bar{x}_1 + x_2 + \bar{x}_3)$
1	1	0	0	$(\bar{x}_1 + \bar{x}_2 + x_3)$
1	1	1	1	

# Let's Derive the POS form

$x_1$	$x_2$	$x_3$	$f$	
0	0	0	0	$(x_1 + x_2 + x_3)$
0	0	1	1	
0	1	0	1	
0	1	1	0	$(x_1 + \bar{x}_2 + \bar{x}_3)$
1	0	0	1	
1	0	1	0	$(\bar{x}_1 + x_2 + \bar{x}_3)$
1	1	0	0	$(\bar{x}_1 + \bar{x}_2 + x_3)$
1	1	1	1	

$$\begin{aligned}f &= M_0 \cdot M_3 \cdot M_5 \cdot M_6 \\&= (x_1 + x_2 + x_3)(x_1 + \bar{x}_2 + \bar{x}_3)(\bar{x}_1 + x_2 + \bar{x}_3)(\bar{x}_1 + \bar{x}_2 + x_3)\end{aligned}$$

# Product-of-sums realization



[ Figure 2.32b from the textbook ]

# **Function Synthesis**

## **Example 2.10**

Implement the function  $f(x_1, x_2, x_3) = \sum m(2, 3, 4, 6, 7)$

# Minterms and Maxterms (with three variables)

Row number	$x_1$	$x_2$	$x_3$	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1 \bar{x}_2 \bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1 \bar{x}_2 x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1 x_2 \bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1 x_2 x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1 \bar{x}_2 \bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1 \bar{x}_2 x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1 x_2 \bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1 x_2 x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

[ Figure 2.22 from the textbook ]

# minterms and Maxterms (with three variables)

Row number	$x_1$	$x_2$	$x_3$	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1 \bar{x}_2 \bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1 \bar{x}_2 x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1 x_2 \bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1 x_2 x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1 \bar{x}_2 \bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1 \bar{x}_2 x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1 x_2 \bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1 x_2 x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

$$f(x_1, x_2, x_3) = \sum m(2, 3, 4, 6, 7)$$

- The SOP expression is:

$$\begin{aligned} f &= m_2 + m_3 + m_4 + m_6 + m_7 \\ &= \bar{x}_1 x_2 \bar{x}_3 + \bar{x}_1 x_2 x_3 + x_1 \bar{x}_2 \bar{x}_3 + x_1 x_2 \bar{x}_3 + x_1 x_2 x_3 \end{aligned}$$

- This could be simplified as follows:

$$\begin{aligned} f &= \bar{x}_1 x_2 (\bar{x}_3 + x_3) + x_1 (\bar{x}_2 + x_2) \bar{x}_3 + x_1 x_2 (\bar{x}_3 + x_3) \\ &= \bar{x}_1 x_2 + x_1 \bar{x}_3 + x_1 x_2 \\ &= (\bar{x}_1 + x_1) x_2 + x_1 \bar{x}_3 \\ &= x_2 + x_1 \bar{x}_3 \end{aligned}$$

# Recall Property 14a

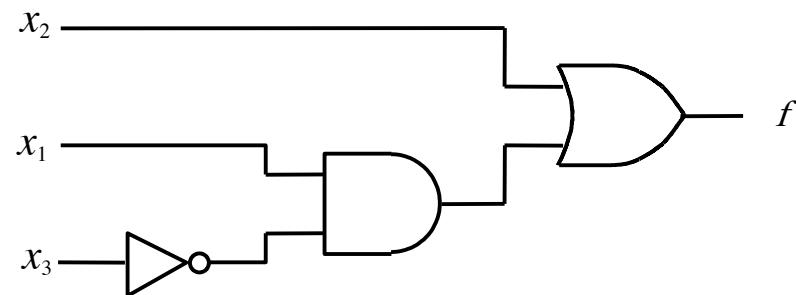
$$14a. \quad x \cdot y + x \cdot \bar{y} = x$$

Combining

$$14b. \quad (x + y) \cdot (x + \bar{y}) = x$$

# SOP realization of the function

The SOP expression is:  $f = x_2 + x_1\bar{x}_3$



[ Figure 2.30a from the textbook ]

## **Example 2.12**

Implement the function  $f(x_1, x_2, x_3) = \prod M(0, 1, 5)$ ,

which is equivalent to  $f(x_1, x_2, x_3) = \sum m(2, 3, 4, 6, 7)$

# minterms and Maxterms (with three variables)

Row number	$x_1$	$x_2$	$x_3$	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1 \bar{x}_2 \bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1 \bar{x}_2 x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1 x_2 \bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1 x_2 x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1 \bar{x}_2 \bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1 \bar{x}_2 x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1 x_2 \bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1 x_2 x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

$$f(x_1, x_2, x_3) = \Pi M(0, 1, 5)$$

- **The POS expression is:**

$$\begin{aligned} f &= M_0 \cdot M_1 \cdot M_5 \\ &= (x_1 + x_2 + x_3)(x_1 + x_2 + \bar{x}_3)(\bar{x}_1 + x_2 + \bar{x}_3) \end{aligned}$$

- **This could be simplified as follows:**

$$\begin{aligned} f &= (x_1 + x_2 + x_3)(x_1 + x_2 + \bar{x}_3)(x_1 + x_2 + \bar{x}_3)(\bar{x}_1 + x_2 + \bar{x}_3) \\ &= ((x_1 + x_2) + x_3)((x_1 + x_2) + \bar{x}_3)(x_1 + (x_2 + \bar{x}_3))(\bar{x}_1 + (x_2 + \bar{x}_3)) \\ &= ((x_1 + x_2) + x_3\bar{x}_3)(x_1\bar{x}_1 + (x_2 + \bar{x}_3)) \\ &= (x_1 + x_2)(x_2 + \bar{x}_3) \end{aligned}$$

# Recall Property 14b

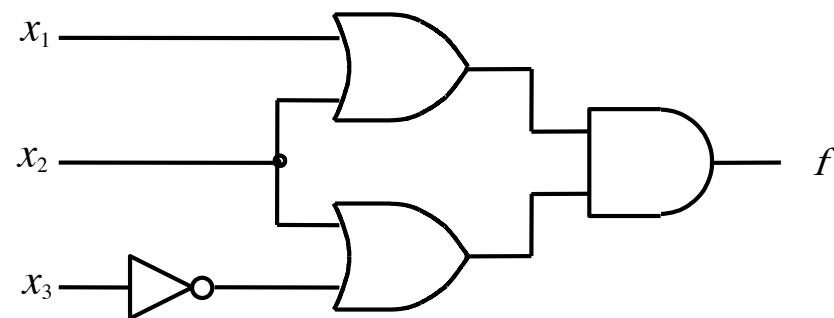
$$14a. \quad x \cdot y + x \cdot \bar{y} = x$$

Combining

$$14b. \quad (x + y) \cdot (x + \bar{y}) = x$$

# POS realization of the function

The POS expression is:  $f = (x_1 + x_2)(x_2 + \bar{x}_3)$



[ Figure 2.29a from the textbook ]

# **More Examples**

## **Example 2.14**

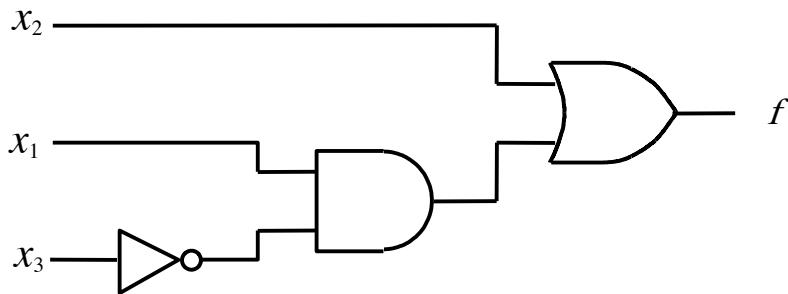
Implement the function  $f(x_1, x_2, x_3) = \sum m(2, 3, 4, 6, 7)$  using only NAND gates.

## **Example 2.14**

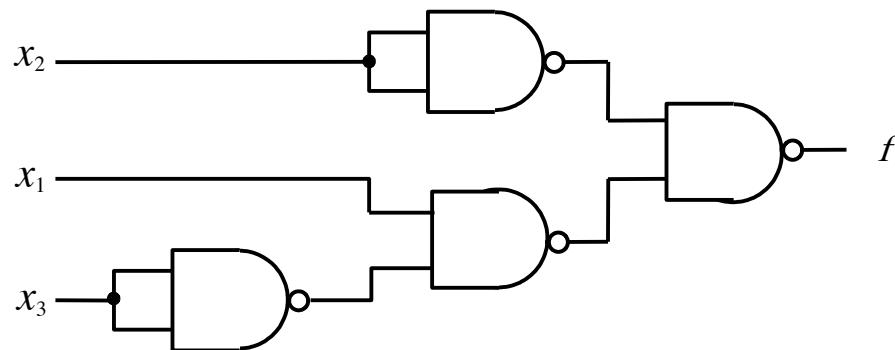
Implement the function  $f(x_1, x_2, x_3) = \sum m(2, 3, 4, 6, 7)$  using only NAND gates.

The SOP expression is:  $f = x_2 + x_1\bar{x}_3$

# NAND-gate realization of the function



(a) SOP implementation



(b) NAND implementation

[ Figure 2.30 from the textbook ]

## **Example 2.13**

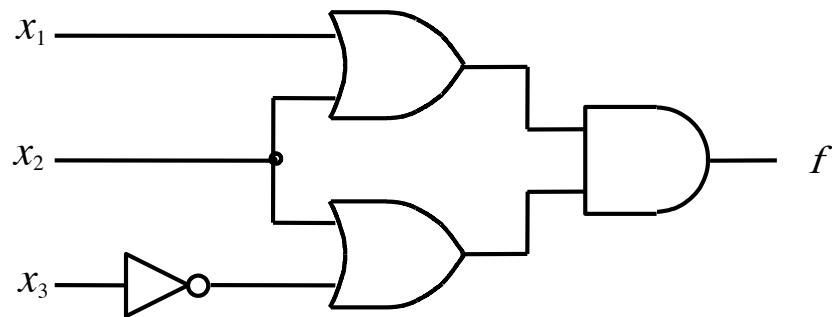
Implement the function  $f(x_1, x_2, x_3) = \sum m(2, 3, 4, 6, 7)$  using only NOR gates.

## Example 2.13

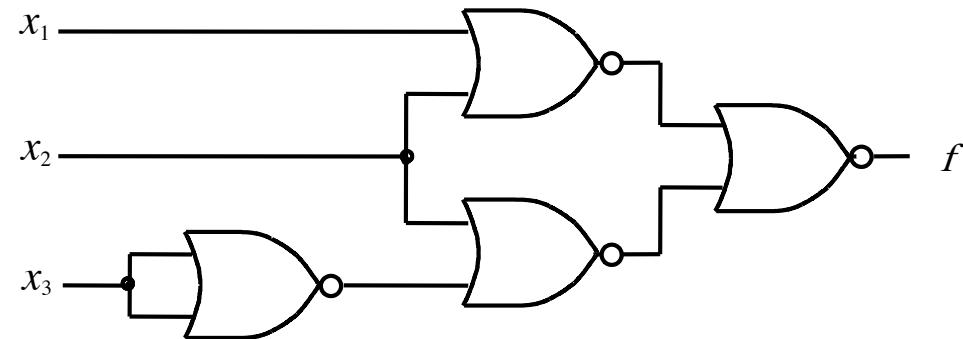
Implement the function  $f(x_1, x_2, x_3) = \sum m(2, 3, 4, 6, 7)$  using only NOR gates.

The POS expression is:  $f = (x_1 + x_2)(x_2 + \bar{x}_3)$

# NOR-gate realization of the function



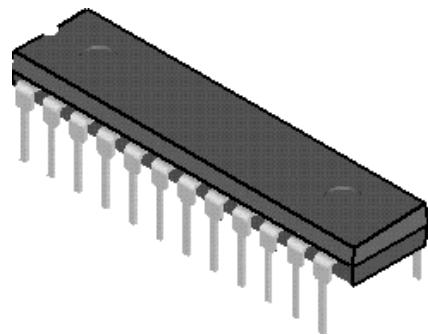
(a) POS implementation



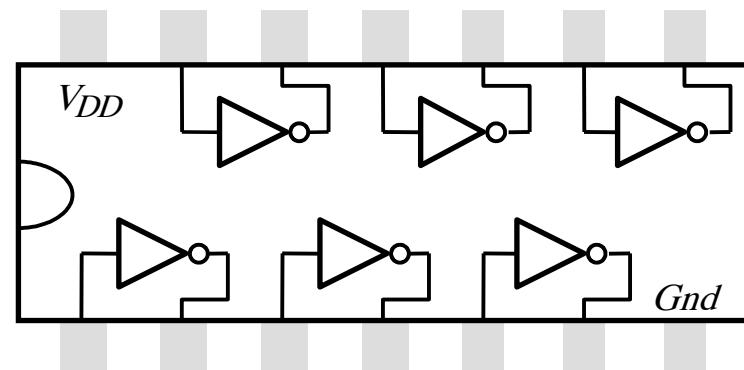
(b) NOR implementation

[ Figure 2.29 from the textbook ]

# **Implementation with Chips**



(a) Dual-inline package



(b) Structure of 7404 chip

Figure B.21. A 7400-series chip.

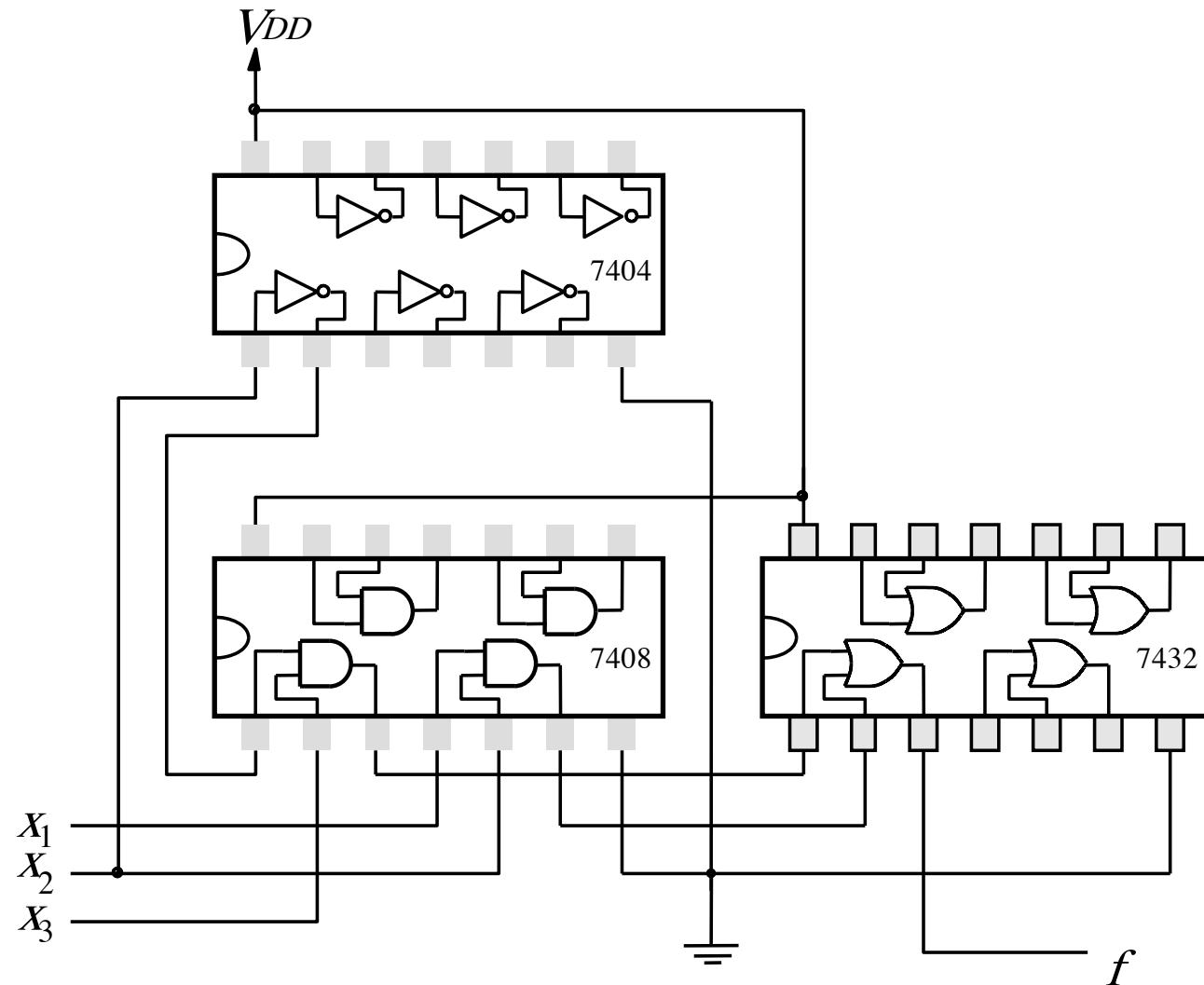


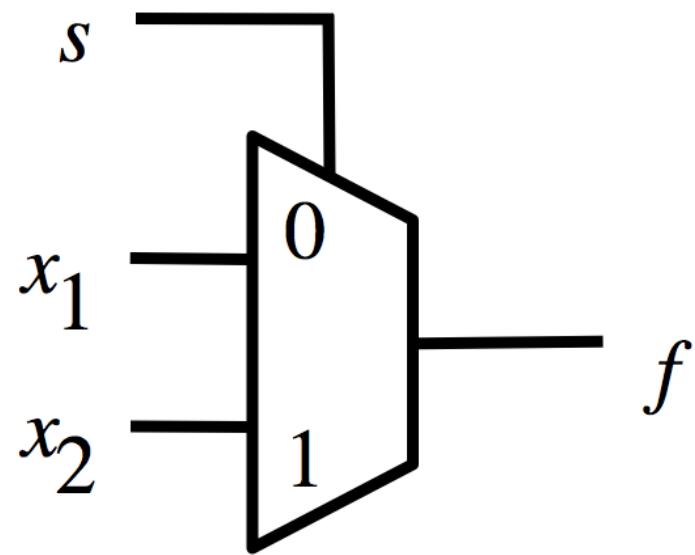
Figure B.22. An implementation of  $f = x_1x_2 + \bar{x}_2x_3$ .

# **Multiplexers**

## 2-to-1 Multiplexer (Definition)

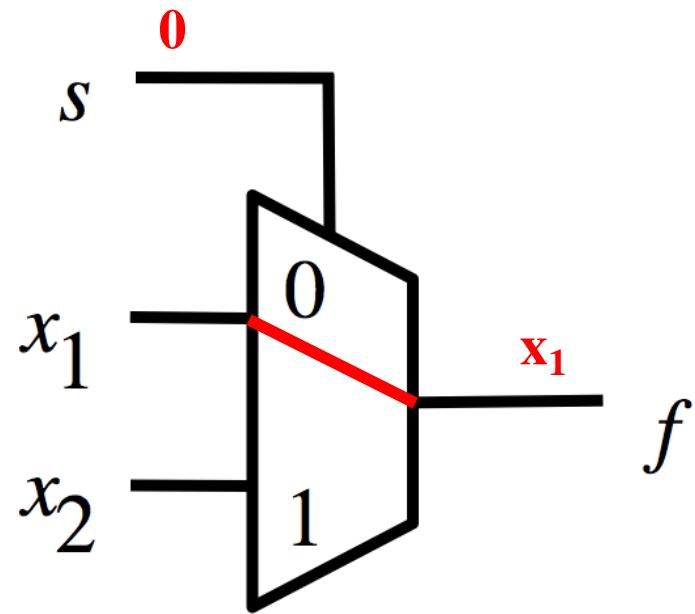
- Has two inputs:  $x_1$  and  $x_2$
- Also has another input line  $s$
- If  $s=0$ , then the output is equal to  $x_1$
- If  $s=1$ , then the output is equal to  $x_2$

# Graphical Symbol for a 2-to-1 Multiplexer

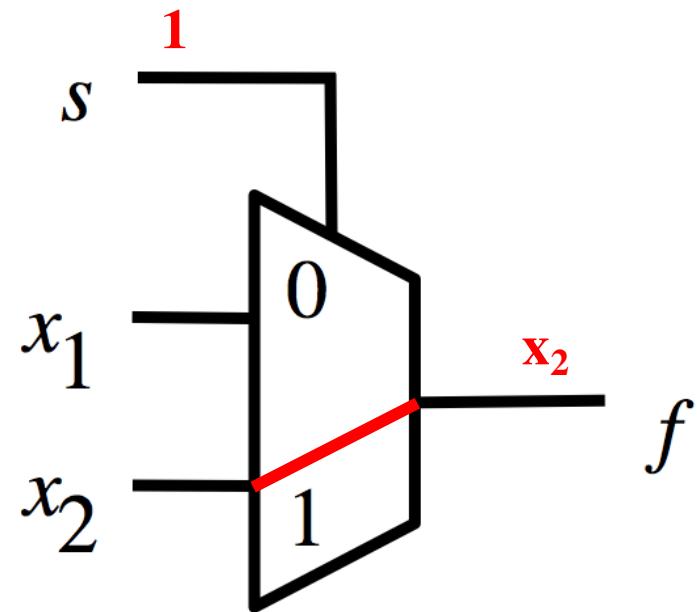


[ Figure 2.33c from the textbook ]

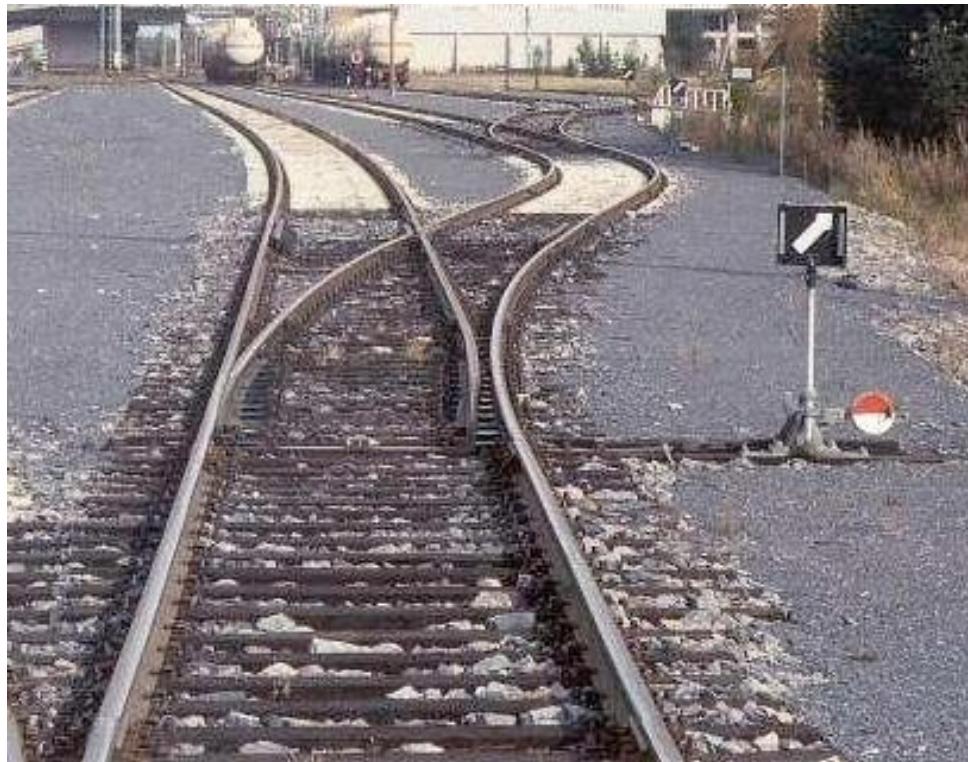
# Analysis of the 2-to-1 Multiplexer (when the input $s=0$ )



# Analysis of the 2-to-1 Multiplexer (when the input $s=1$ )

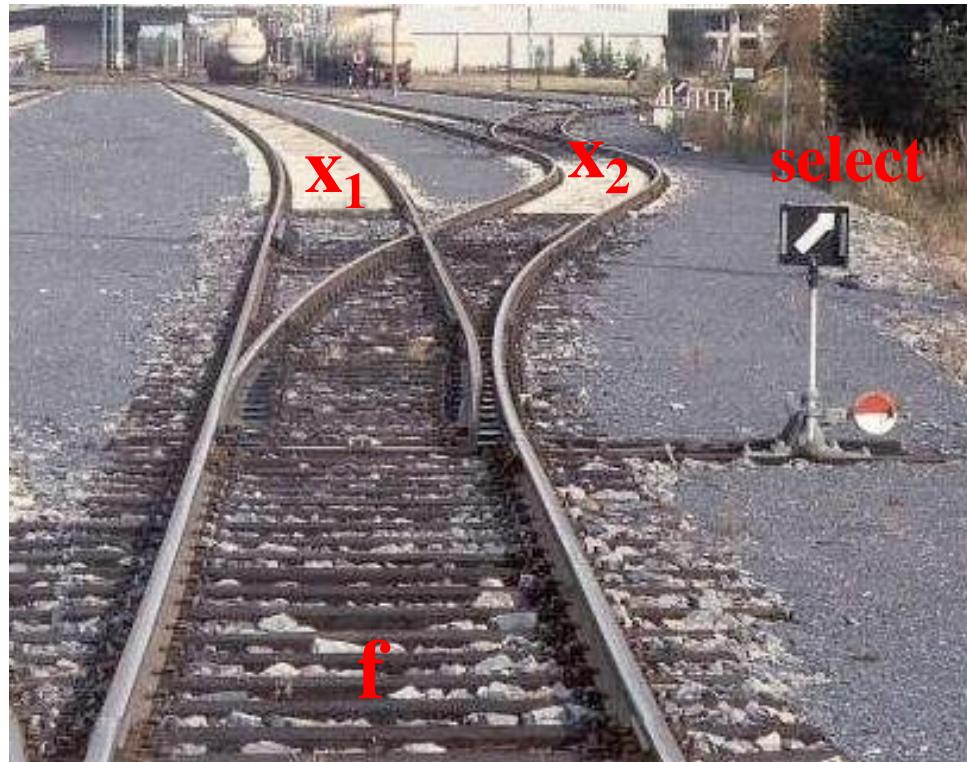


# Analogy: Railroad Switch



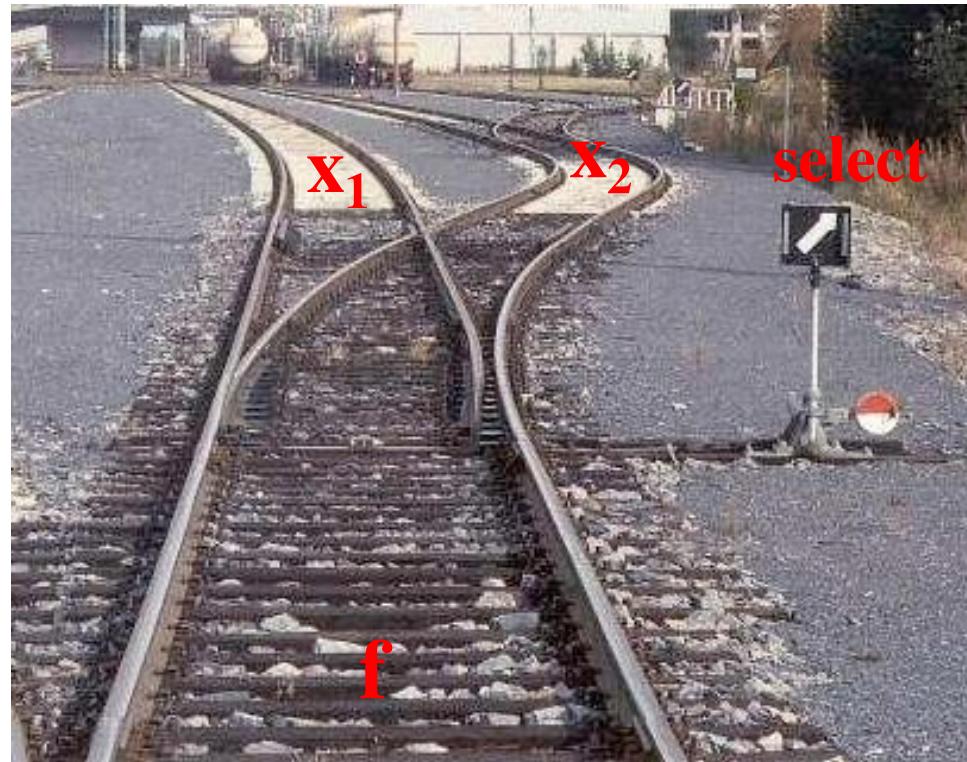
[http://en.wikipedia.org/wiki/Railroad\\_switch\]](http://en.wikipedia.org/wiki/Railroad_switch)

# Analogy: Railroad Switch



[http://en.wikipedia.org/wiki/Railroad\\_switch\]](http://en.wikipedia.org/wiki/Railroad_switch)

# Analogy: Railroad Switch



This is not a perfect analogy because the trains can go in either direction, while the multiplexer would only allow them to go from top to bottom.

[http://en.wikipedia.org/wiki/Railroad\\_switch](http://en.wikipedia.org/wiki/Railroad_switch)

# Truth Table for a 2-to-1 Multiplexer

$s\ x_1\ x_2$	$f(s, x_1, x_2)$
0 0 0	0
0 0 1	0
0 1 0	1
0 1 1	1
1 0 0	0
1 0 1	1
1 1 0	0
1 1 1	1

[ Figure 2.33a from the textbook ]

# Let's Derive the SOP form

$s \ x_1 \ x_2$	$f(s, x_1, x_2)$
0 0 0	0
0 0 1	0
0 1 0	1
0 1 1	1
1 0 0	0
1 0 1	1
1 1 0	0
1 1 1	1

# Let's Derive the SOP form

$s \ x_1 \ x_2$	$f(s, x_1, x_2)$
0 0 0	0
0 0 1	0
0 1 0	1
0 1 1	1
1 0 0	0
1 0 1	1
1 1 0	0
1 1 1	1

# Let's Derive the SOP form

$s \ x_1 \ x_2$	$f(s, x_1, x_2)$
0 0 0	0
0 0 1	0
0 1 0	1
0 1 1	1
1 0 0	0
1 0 1	1
1 1 0	0
1 1 1	1

Where should we  
put the negation signs?

$s \ x_1 \ x_2$

# Let's Derive the SOP form

$s \ x_1 \ x_2$	$f(s, x_1, x_2)$	
0 0 0	0	
0 0 1	0	
0 1 0	1	$\bar{s} \ x_1 \ \bar{x}_2$
0 1 1	1	$\bar{s} \ x_1 \ x_2$
1 0 0	0	
1 0 1	1	$s \ \bar{x}_1 \ x_2$
1 1 0	0	
1 1 1	1	$s \ x_1 \ x_2$

# Let's Derive the SOP form

$s \ x_1 \ x_2$	$f(s, x_1, x_2)$
0 0 0	0
0 0 1	0
0 1 0	1
0 1 1	1
1 0 0	0
1 0 1	1
1 1 0	0
1 1 1	1

$$f(s, x_1, x_2) = \overline{s}x_1\overline{x}_2 + \overline{s}x_1x_2 + s\overline{x}_1x_2 + sx_1x_2$$

**Let's simplify this expression**

$$f(s, x_1, x_2) = \bar{s}x_1\bar{x}_2 + \bar{s}x_1x_2 + s\bar{x}_1x_2 + sx_1x_2$$

# Let's simplify this expression

$$f(s, x_1, x_2) = \bar{s}x_1\bar{x}_2 + \bar{s}x_1x_2 + s\bar{x}_1x_2 + sx_1x_2$$

$$f(s, x_1, x_2) = \bar{s}x_1(\bar{x}_2 + x_2) + s(\bar{x}_1 + x_1)x_2$$

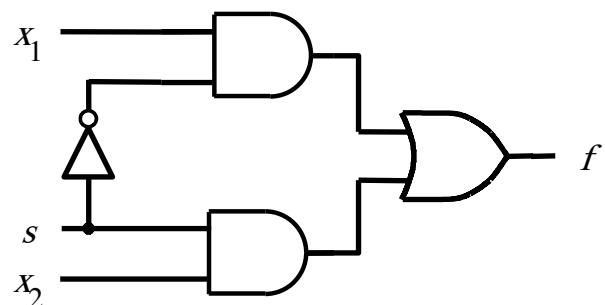
# Let's simplify this expression

$$f(s, x_1, x_2) = \bar{s}x_1\bar{x}_2 + \bar{s}x_1x_2 + s\bar{x}_1x_2 + sx_1x_2$$

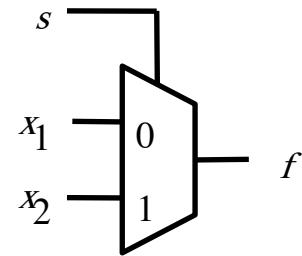
$$f(s, x_1, x_2) = \bar{s}x_1(\bar{x}_2 + x_2) + s(\bar{x}_1 + x_1)x_2$$

$$f(s, x_1, x_2) = \bar{s}x_1 + s x_2$$

# Circuit for 2-to-1 Multiplexer



(b) Circuit

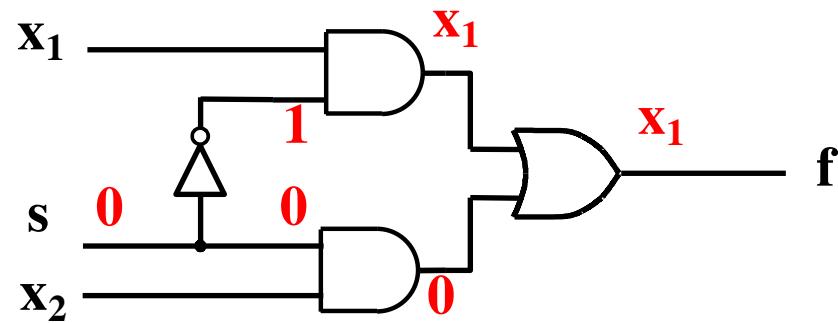


(c) Graphical symbol

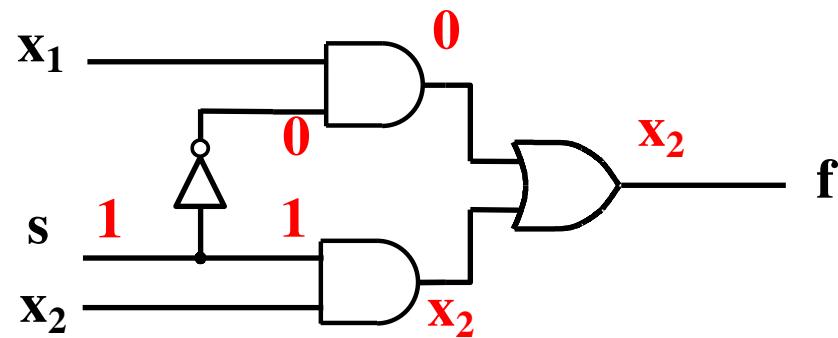
$$f(s, x_1, x_2) = \bar{s}x_1 + s x_2$$

[ Figure 2.33b-c from the textbook ]

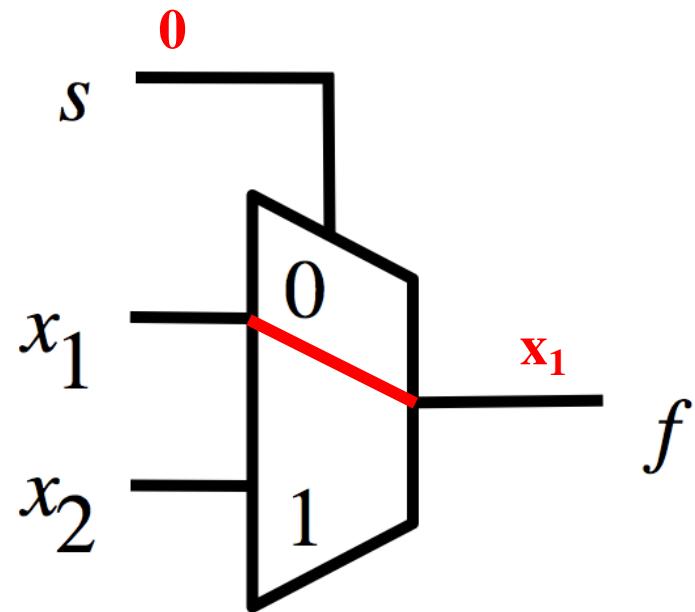
# Analysis of the 2-to-1 Multiplexer (when the input s=0)



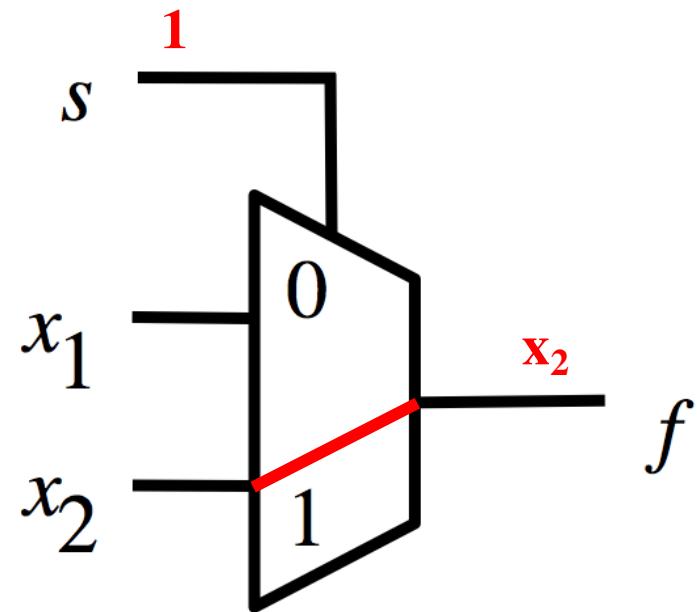
# Analysis of the 2-to-1 Multiplexer (when the input s=1)



# Analysis of the 2-to-1 Multiplexer (when the input $s=0$ )



# Analysis of the 2-to-1 Multiplexer (when the input $s=1$ )



# More Compact Truth-Table Representation

$s$	$x_1$	$x_2$	$f(s, x_1, x_2)$
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

(a) Truth table

$s$	$f(s, x_1, x_2)$
0	$x_1$
1	$x_2$

[ Figure 2.33 from the textbook ]

# **4-to-1 Multiplexer**

# 4-to-1 Multiplexer (Definition)

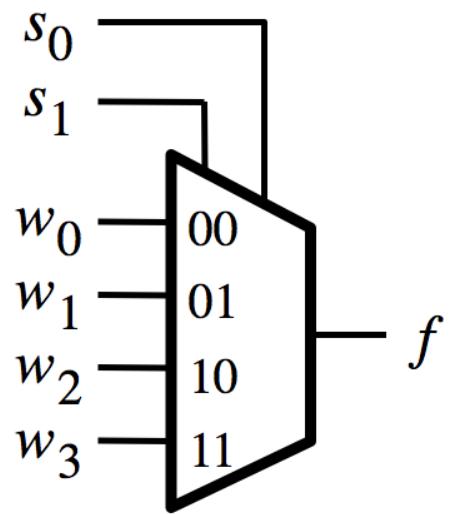
- Has four inputs:  $w_0$ ,  $w_1$ ,  $w_2$ ,  $w_3$
- Also has two select lines:  $s_1$  and  $s_0$
- If  $s_1=0$  and  $s_0=0$ , then the output f is equal to  $w_0$
- If  $s_1=0$  and  $s_0=1$ , then the output f is equal to  $w_1$
- If  $s_1=1$  and  $s_0=0$ , then the output f is equal to  $w_2$
- If  $s_1=1$  and  $s_0=1$ , then the output f is equal to  $w_3$

# 4-to-1 Multiplexer (Definition)

- Has four inputs:  $w_0$ ,  $w_1$ ,  $w_2$ ,  $w_3$
- Also has two select lines:  $s_1$  and  $s_0$
- If  $s_1=0$  and  $s_0=0$ , then the output f is equal to  $w_0$
- If  $s_1=0$  and  $s_0=1$ , then the output f is equal to  $w_1$
- If  $s_1=1$  and  $s_0=0$ , then the output f is equal to  $w_2$
- If  $s_1=1$  and  $s_0=1$ , then the output f is equal to  $w_3$

We'll talk more about this when we get to chapter 4, but here is a quick preview.

# Graphical Symbol and Truth Table



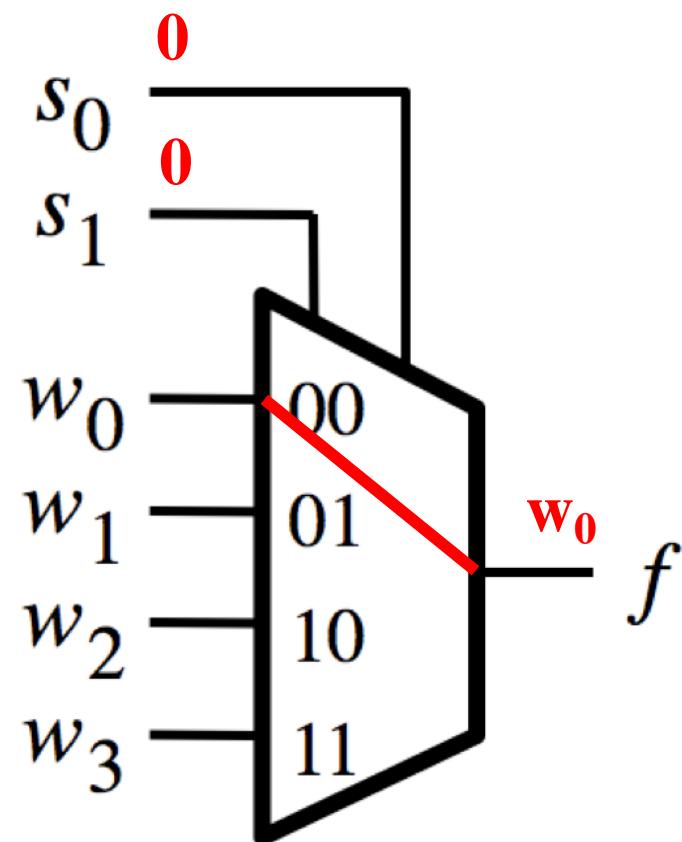
(a) Graphic symbol

$s_1$	$s_0$	$f$
0	0	$w_0$
0	1	$w_1$
1	0	$w_2$
1	1	$w_3$

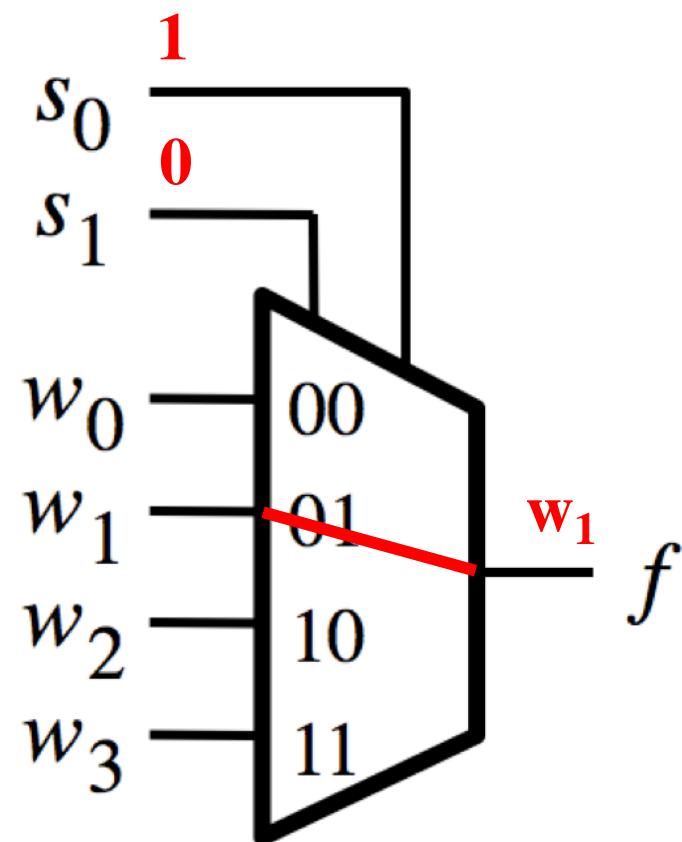
(b) Truth table

[ Figure 4.2a-b from the textbook ]

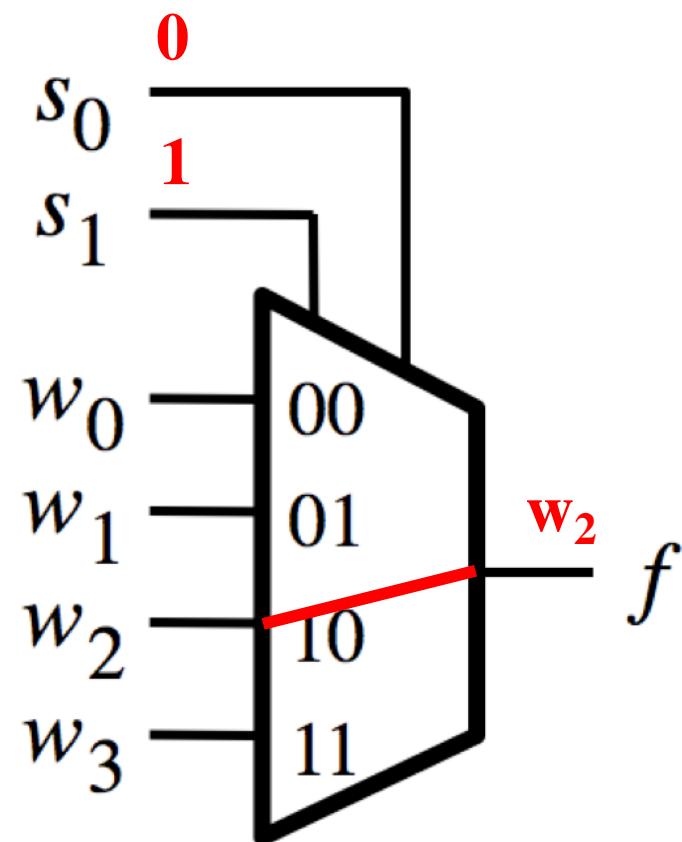
# Analysis of the 4-to-1 Multiplexer ( $s_1=0$ and $s_0=0$ )



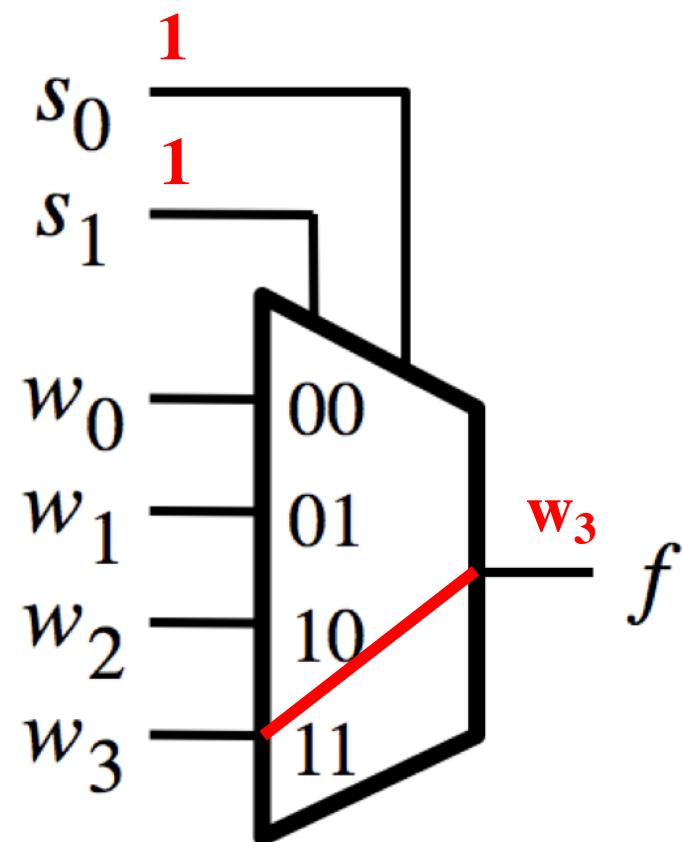
# Analysis of the 4-to-1 Multiplexer ( $s_1=0$ and $s_0=1$ )



# Analysis of the 4-to-1 Multiplexer ( $s_1=1$ and $s_0=0$ )



# Analysis of the 4-to-1 Multiplexer ( $s_1=1$ and $s_0=1$ )



# **The long-form truth table**

# The long-form truth table

$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F
0 0	0	0	0	0	0	0 1	0	0	0	0	0	1 0	0	0	0	0	0	1 1	0	0	0	0	0
	0	0	0	1	1		0	0	0	1	0		0	0	0	1	0		0	0	0	1	0
	0	0	1	0	0		0	0	1	0	1		0	0	1	0	0		0	0	1	0	0
	0	0	1	1	1		0	0	1	1	1		0	0	1	1	0		0	0	1	1	0
	0	1	0	0	0		0	1	0	0	0		0	1	0	0	1		0	1	0	0	0
	0	1	0	1	1		0	1	0	1	0		0	1	0	1	1		0	1	0	1	0
	0	1	1	0	0		0	1	1	0	1		0	1	1	0	1		0	1	1	0	0
	0	1	1	1	1		0	1	1	1	1		0	1	1	1	1		0	1	1	1	0
	1	0	0	0	0		1	0	0	0	0		1	0	0	0	0		1	0	0	0	1
	1	0	0	1	1		1	0	0	1	0		1	0	0	1	0		1	0	0	1	1
	1	0	1	0	0		1	0	1	0	1		1	0	1	0	0		1	0	1	0	1
	1	0	1	1	1		1	0	1	1	1		1	0	1	1	0		1	0	1	1	1
	1	1	0	0	0		1	1	0	0	0		1	1	0	0	1		1	1	0	0	1
	1	1	0	1	1		1	1	0	1	0		1	1	0	1	1		1	1	0	1	1
	1	1	1	0	0		1	1	1	0	1		1	1	1	0	1		1	1	1	0	1
	1	1	1	1	1		1	1	1	1	1		1	1	1	1	1		1	1	1	1	1

# The long-form truth table

$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F
0 0	0	0	0	0	0	0 1	0	0	0	0	0	1 0	0	0	0	0	0	1 1	0	0	0	0	0
	0	0	0	1	1		0	0	0	1	0		0	0	0	1	0		0	0	0	1	0
	0	0	1	0	0		0	0	1	0	1		0	0	1	0	0		0	0	1	0	0
	0	0	1	1	1		0	0	1	1	1		0	0	1	1	0		0	0	1	1	0
	0	1	0	0	0		0	1	0	0	0		0	1	0	0	1		0	1	0	0	0
	0	1	0	1	1		0	1	0	1	0		0	1	0	1	1		0	1	0	1	0
	0	1	1	0	0		0	1	1	0	1		0	1	1	0	1		0	1	1	0	0
	0	1	1	1	1		0	1	1	1	1		0	1	1	1	1		0	1	1	1	0
	1	0	0	0	0		1	0	0	0	0		1	0	0	0	0		1	0	0	0	1
	1	0	0	1	1		1	0	0	1	0		1	0	0	1	0		1	0	0	1	1
	1	0	1	0	0		1	0	1	0	1		1	0	1	0	0		1	0	1	0	1
	1	0	1	1	1		1	0	1	1	1		1	0	1	1	0		1	0	1	1	1
	1	1	0	0	0		1	1	0	0	0		1	1	0	0	1		1	1	0	0	1
	1	1	0	1	1		1	1	0	1	0		1	1	0	1	1		1	1	0	1	1
	1	1	1	0	0		1	1	1	0	1		1	1	1	0	1		1	1	1	0	1
	1	1	1	1	1		1	1	1	1	1		1	1	1	1	1		1	1	1	1	1

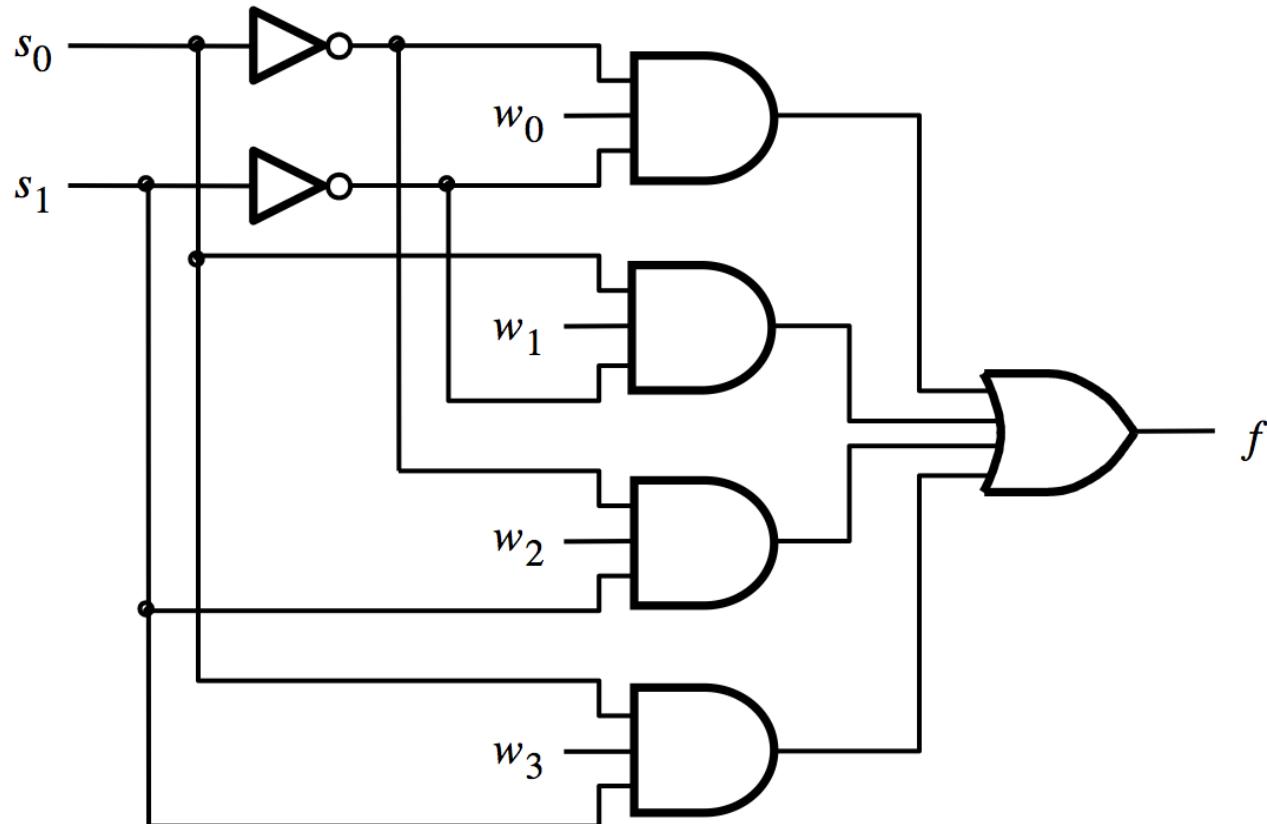
# The long-form truth table

$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F
0 0	0	0	0	0	0	0 1	0	0	0	0	0	1 0	0	0	0	0	0	1 1	0	0	0	0	0
	0	0	0	1	1		0	0	0	1	0		0	0	0	1	0		0	0	0	1	0
	0	0	1	0	0		0	0	1	0	1		0	0	1	0	0		0	0	1	0	0
	0	0	1	1	1		0	0	1	1	1		0	0	1	1	0		0	0	1	1	0
	0	1	0	0	0		0	1	0	0	0		0	1	0	0	1		0	1	0	0	0
	0	1	0	1	1		0	1	0	1	0		0	1	0	1	1		0	1	0	1	0
	0	1	1	0	0		0	1	1	0	1		0	1	1	0	1		0	1	1	0	0
	0	1	1	1	1		0	1	1	1	1		0	1	1	1	1		0	1	1	1	0
	1	0	0	0	0		1	0	0	0	0		1	0	0	0	0		1	0	0	0	1
	1	0	0	1	1		1	0	0	1	0		1	0	0	1	0		1	0	0	1	1
	1	0	1	0	0		1	0	1	0	1		1	0	1	0	0		1	0	1	0	1
	1	0	1	1	1		1	0	1	1	1		1	0	1	1	0		1	0	1	1	1
	1	1	0	0	0		1	1	0	0	0		1	1	0	0	1		1	1	0	0	1
	1	1	0	1	1		1	1	0	1	0		1	1	0	1	1		1	1	0	1	1
	1	1	1	0	0		1	1	1	0	1		1	1	1	0	1		1	1	1	0	1
	1	1	1	1	1		1	1	1	1	1		1	1	1	1	1		1	1	1	1	1

# The long-form truth table

$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F	$S_1 S_0$	I <sub>3</sub>	I <sub>2</sub>	I <sub>1</sub>	I <sub>0</sub>	F
0 0	0	0	0	0	0	0 1	0	0	0	0	0	1 0	0	0	0	0	0	1 1	0	0	0	0	0
	0	0	0	1	1		0	0	0	1	0		0	0	0	1	0		0	0	0	1	0
	0	0	1	0	0		0	0	1	0	1		0	0	1	0	0		0	0	1	0	0
	0	0	1	1	1		0	0	1	1	1		0	0	1	1	0		0	0	1	1	0
	0	1	0	0	0		0	1	0	0	0		0	1	0	0	1		0	1	0	0	0
	0	1	0	1	1		0	1	0	1	0		0	1	0	1	1		0	1	0	1	0
	0	1	1	0	0		0	1	1	0	1		0	1	1	0	1		0	1	1	0	0
	0	1	1	1	1		0	1	1	1	1		0	1	1	1	1		0	1	1	1	0
	1	0	0	0	0		1	0	0	0	0		1	0	0	0	0		1	0	0	0	1
	1	0	0	1	1		1	0	0	1	0		1	0	0	1	0		1	0	0	1	1
	1	0	1	0	0		1	0	1	0	1		1	0	1	0	0		1	0	1	0	1
	1	0	1	1	1		1	0	1	1	1		1	0	1	1	0		1	0	1	1	1
	1	1	0	0	0		1	1	0	0	0		1	1	0	0	1		1	1	0	0	1
	1	1	0	1	1		1	1	0	1	0		1	1	0	1	1		1	1	0	1	1
	1	1	1	0	0		1	1	1	0	1		1	1	1	0	1		1	1	1	0	1
	1	1	1	1	1		1	1	1	1	1		1	1	1	1	1		1	1	1	1	1

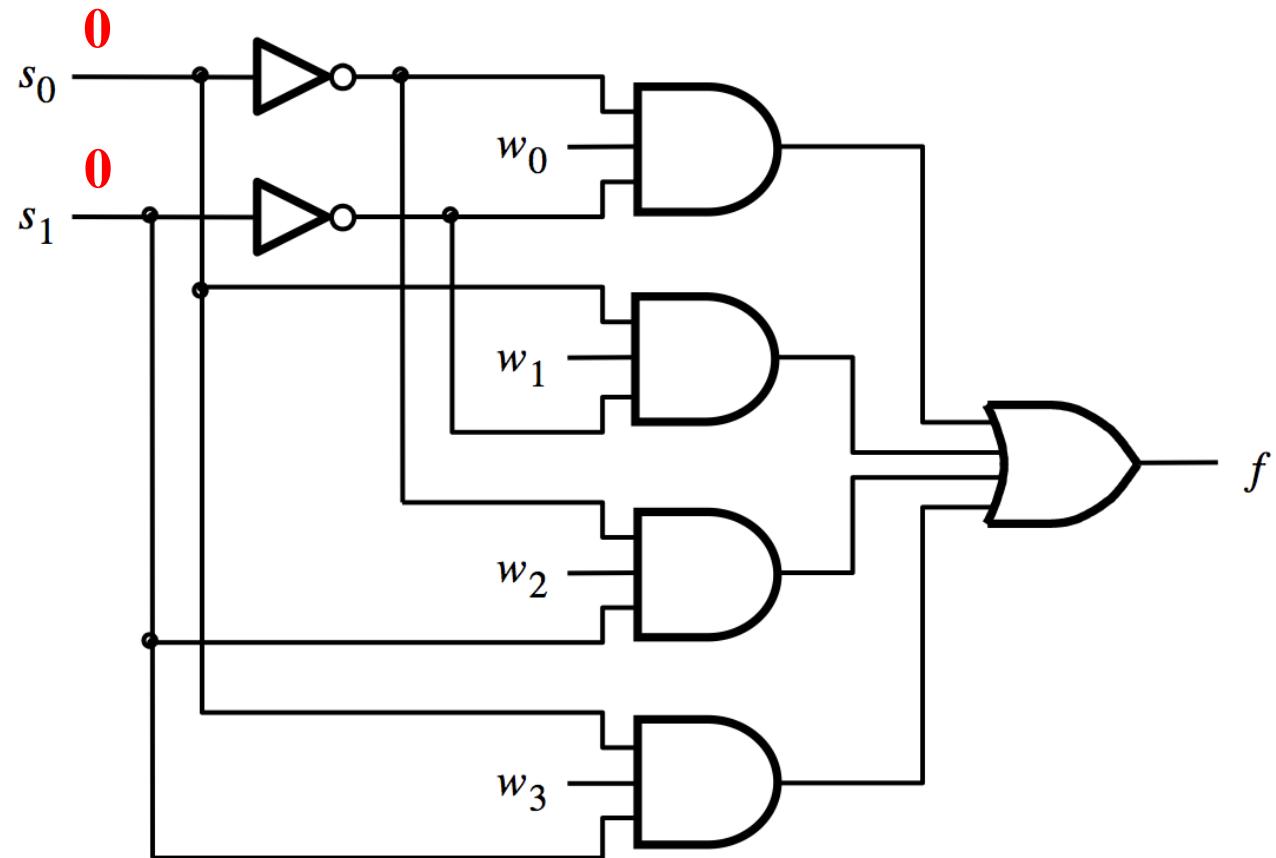
# 4-to-1 Multiplexer (SOP circuit)



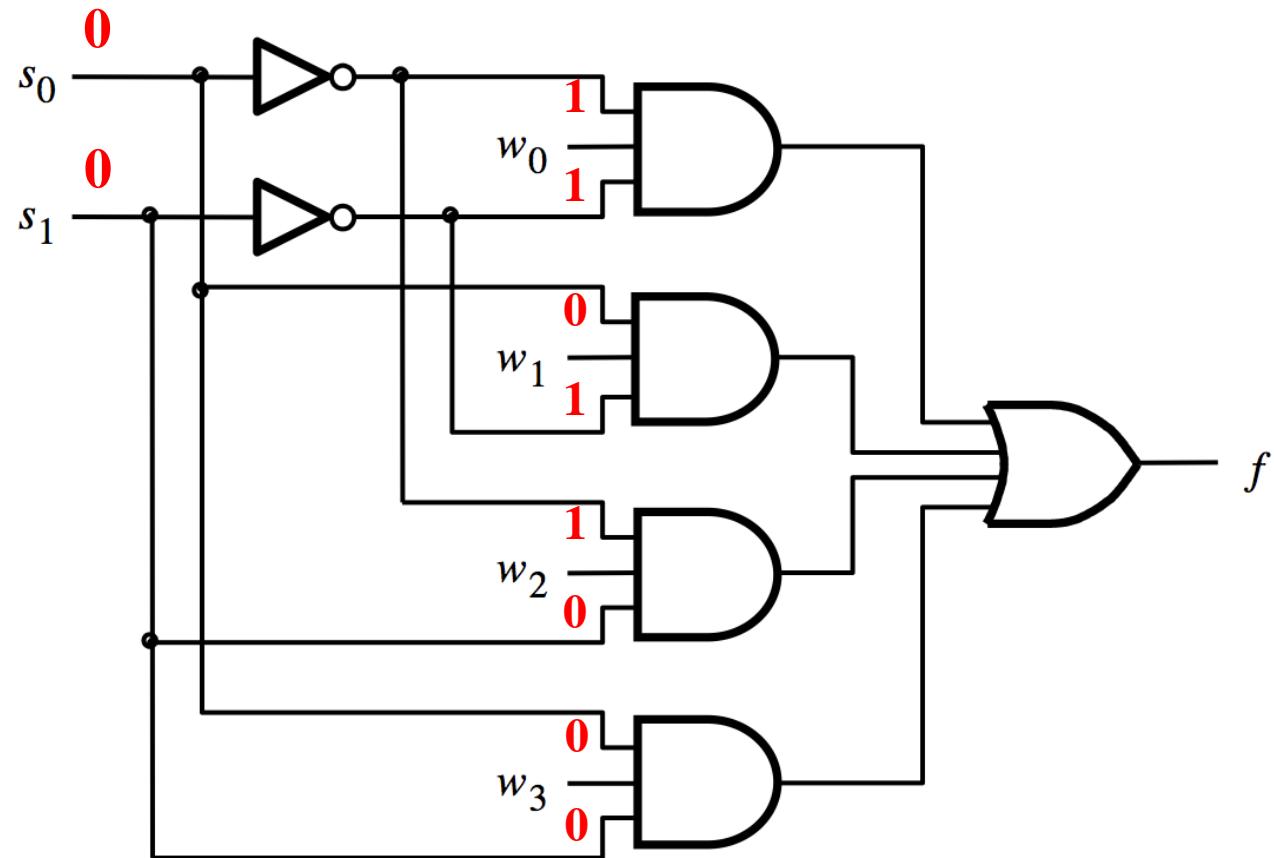
$$f = \overline{s_1} \overline{s_0} w_0 + \overline{s_1} s_0 w_1 + s_1 \overline{s_0} w_2 + s_1 s_0 w_3$$

[ Figure 4.2c from the textbook ]

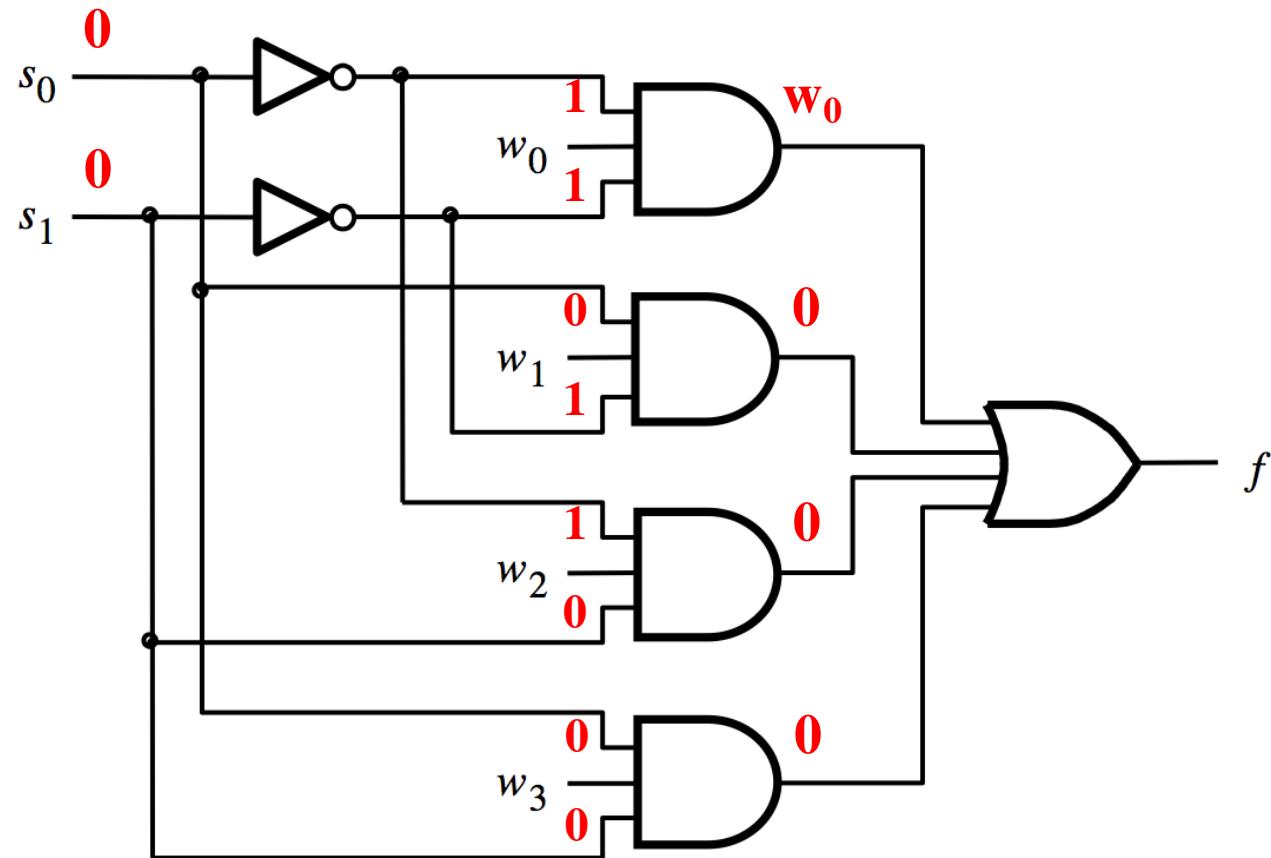
# Analysis of the 4-to-1 Multiplexer ( $s_1=0$ and $s_0=0$ )



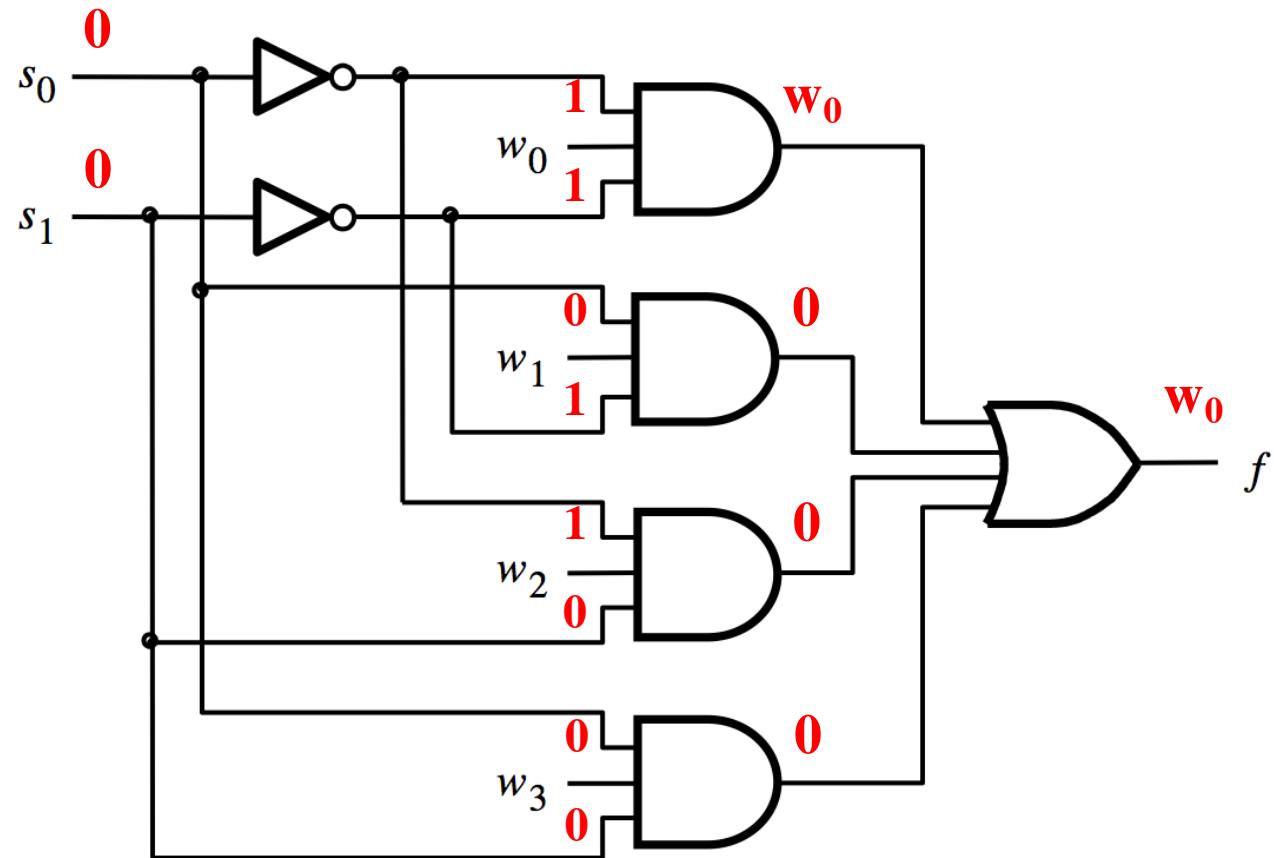
# Analysis of the 4-to-1 Multiplexer ( $s_1=0$ and $s_0=0$ )



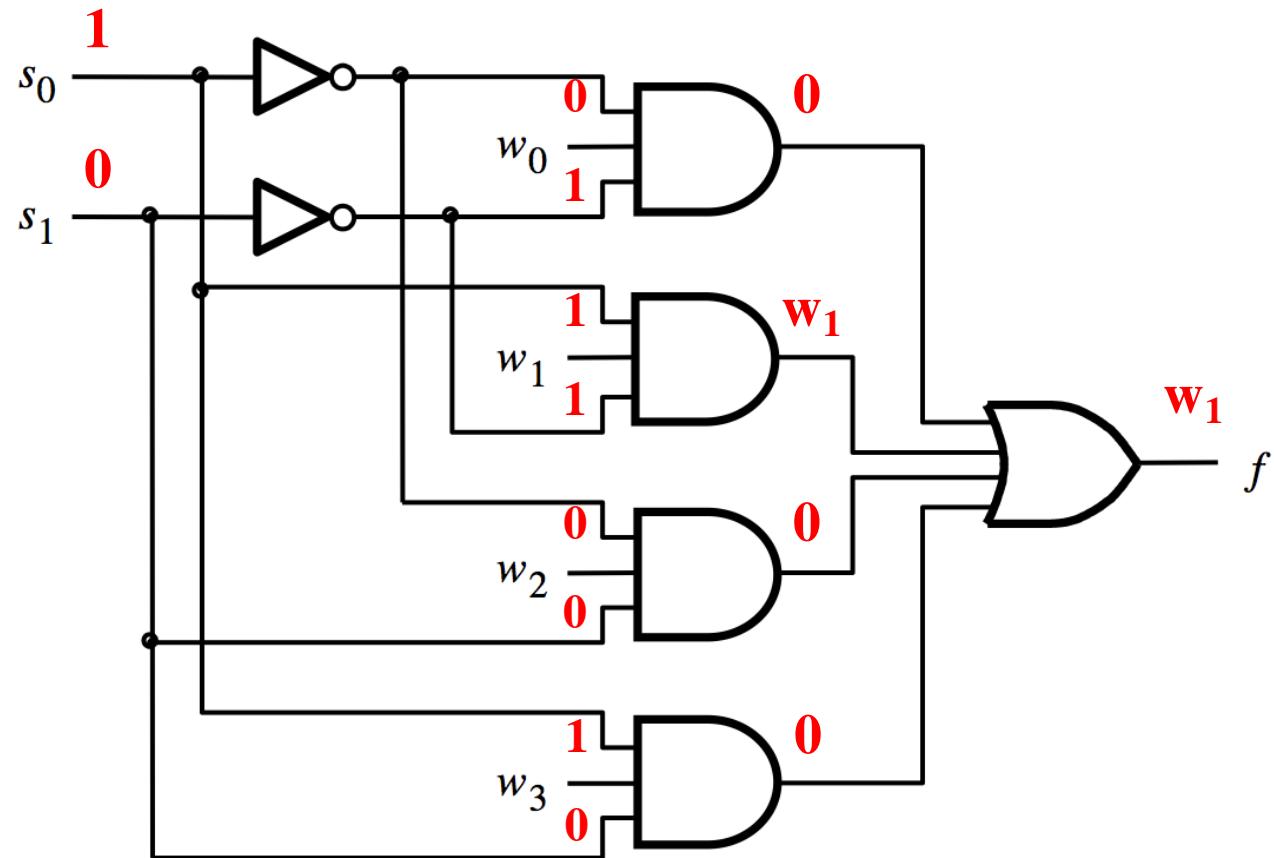
# Analysis of the 4-to-1 Multiplexer ( $s_1=0$ and $s_0=0$ )



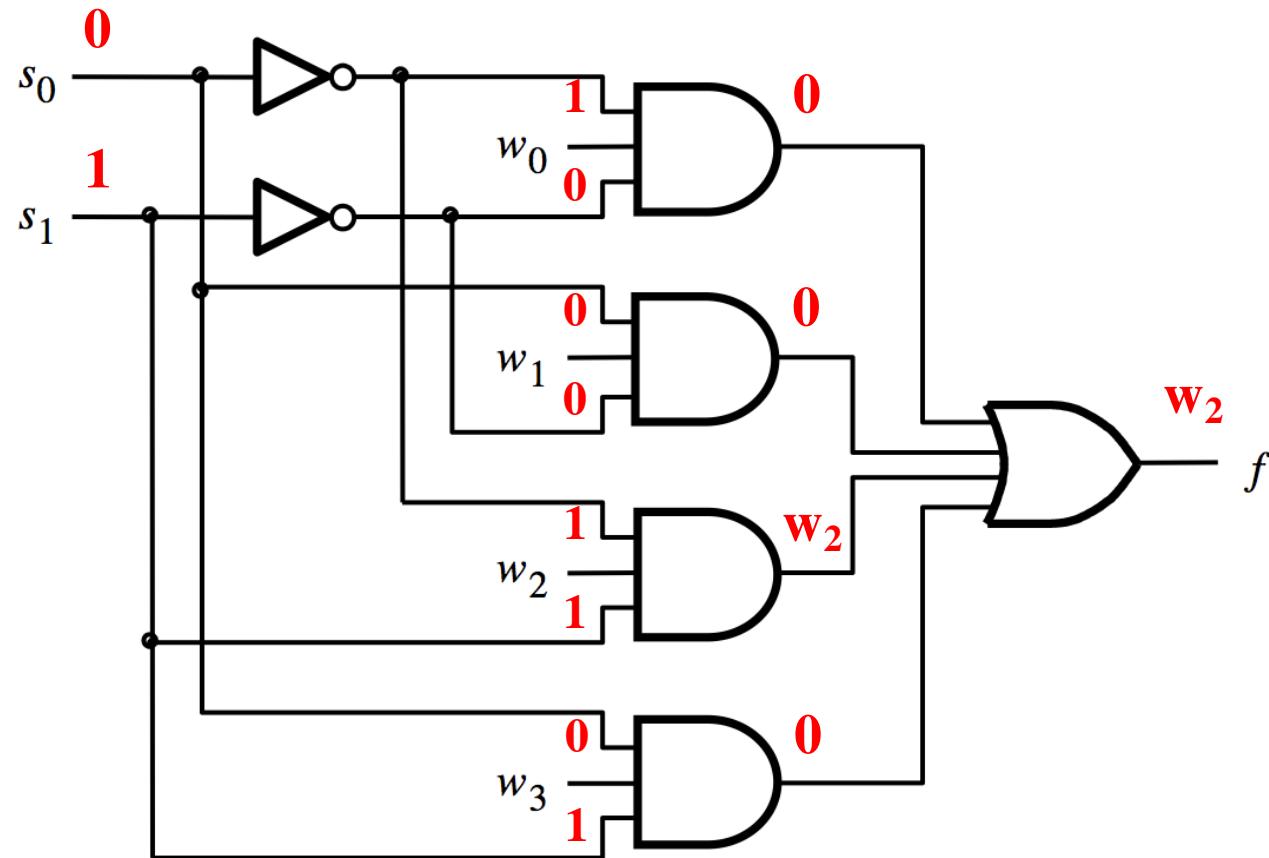
# Analysis of the 4-to-1 Multiplexer ( $s_1=0$ and $s_0=0$ )



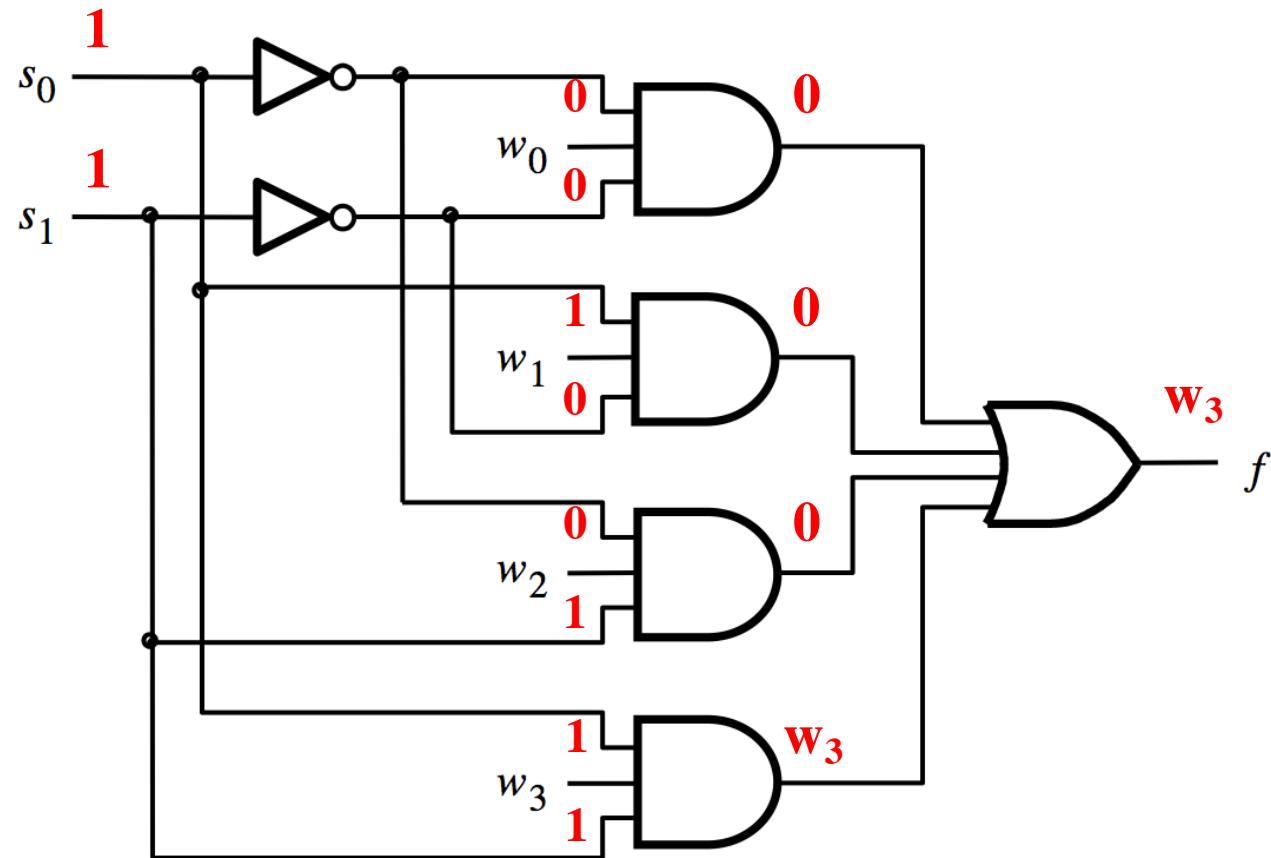
# Analysis of the 4-to-1 Multiplexer ( $s_1=0$ and $s_0=1$ )



# Analysis of the 4-to-1 Multiplexer ( $s_1=1$ and $s_0=0$ )

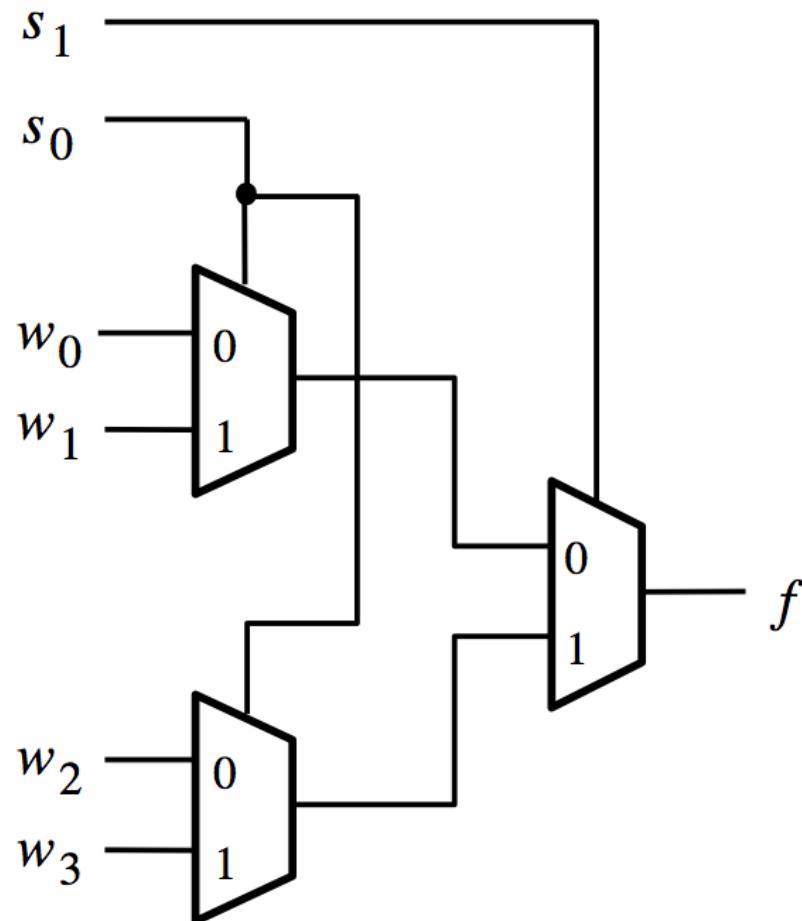


# Analysis of the 4-to-1 Multiplexer ( $s_1=1$ and $s_0=1$ )



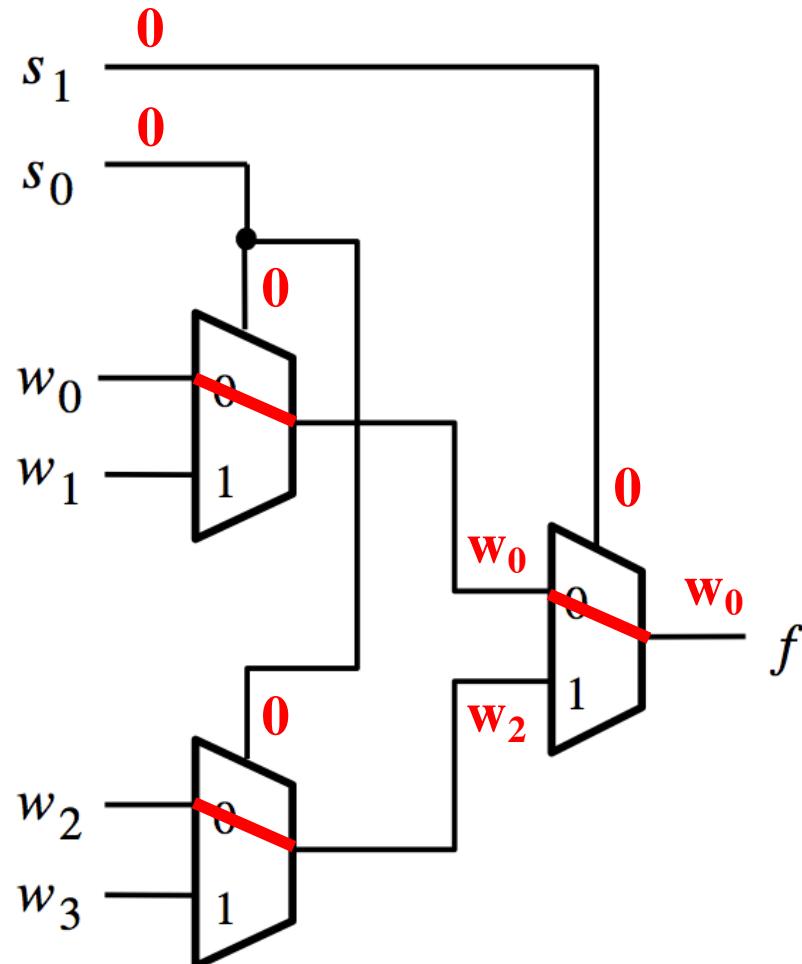
# **4-to-1 Multiplexer (alternative implementation)**

# Using three 2-to-1 multiplexers to build one 4-to-1 multiplexer



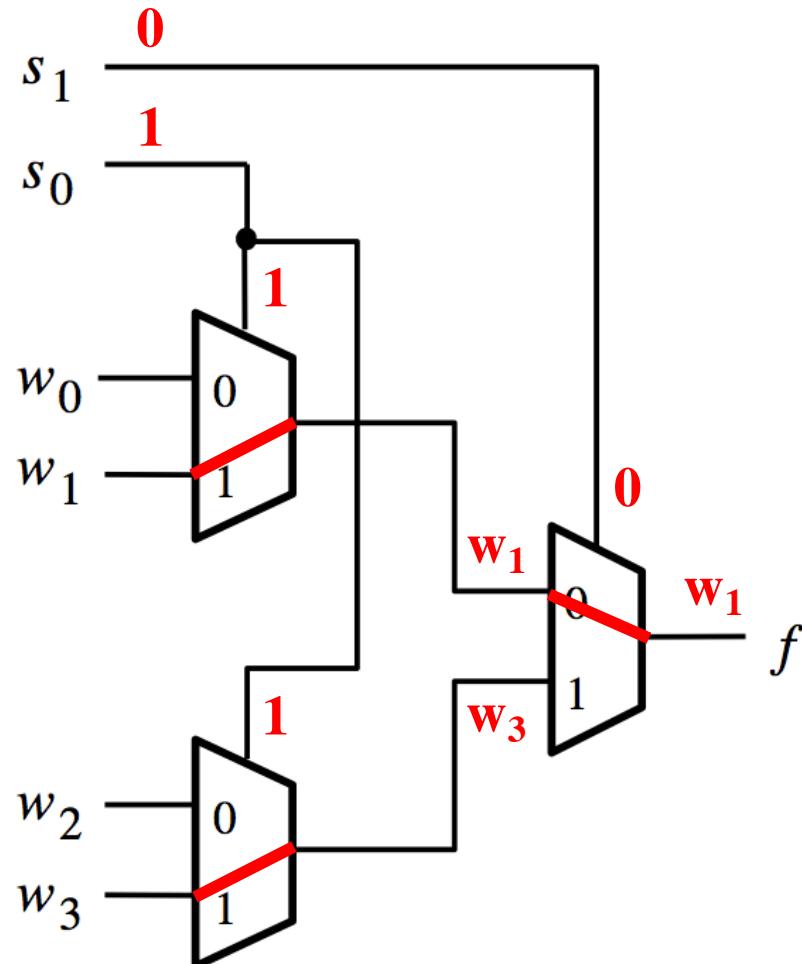
[ Figure 4.3 from the textbook ]

# Analysis of the Hierarchical Implementation ( $s_1=0$ and $s_0=0$ )



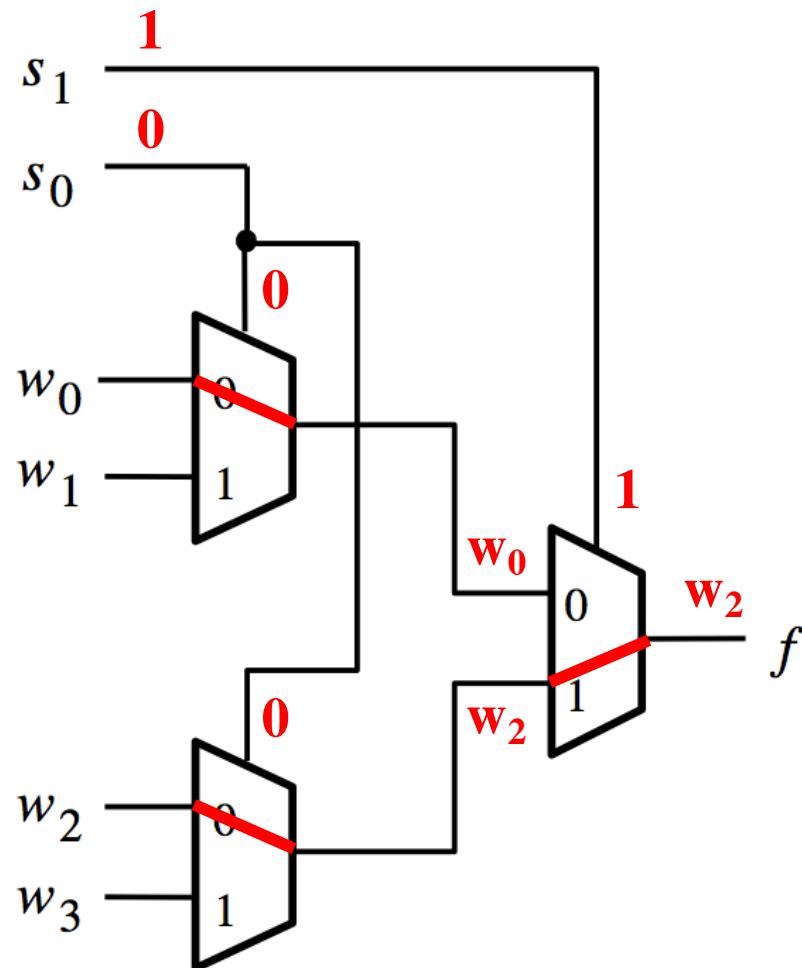
[ Figure 4.3 from the textbook ]

# Analysis of the Hierarchical Implementation ( $s_1=0$ and $s_0=1$ )



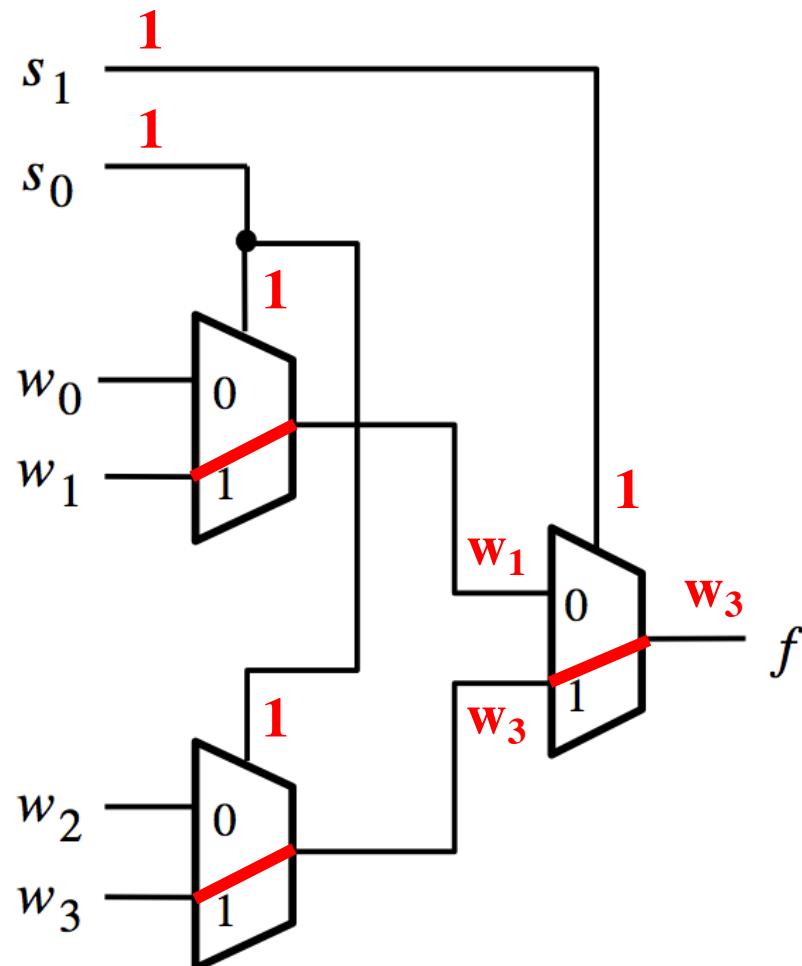
[ Figure 4.3 from the textbook ]

# Analysis of the Hierarchical Implementation ( $s_1=1$ and $s_0=0$ )



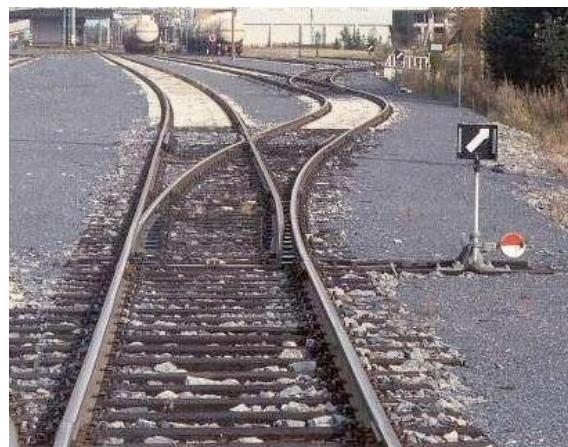
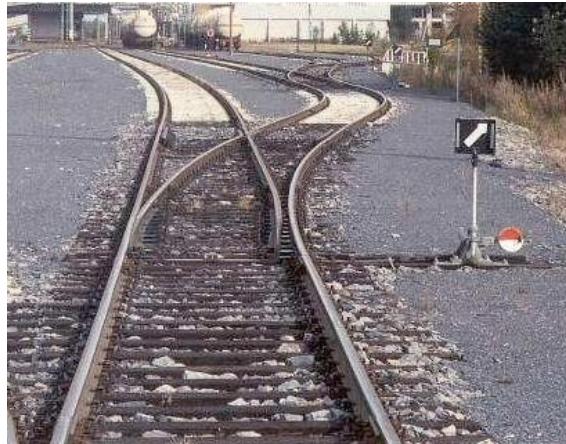
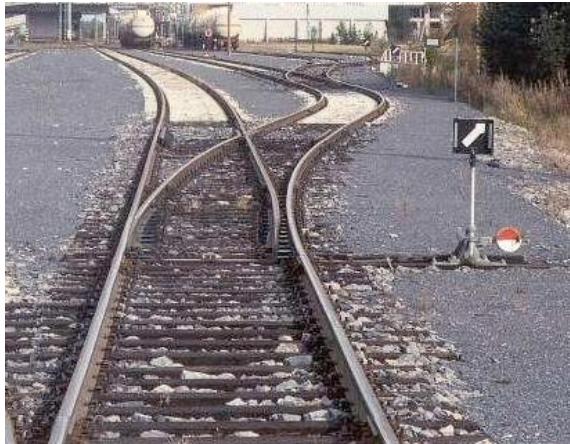
[ Figure 4.3 from the textbook ]

# Analysis of the Hierarchical Implementation ( $s_1=1$ and $s_0=1$ )



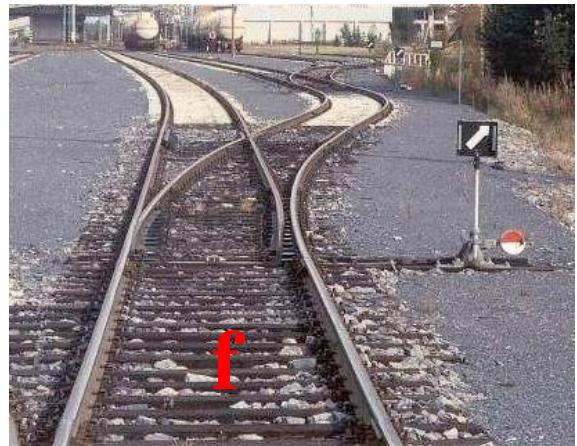
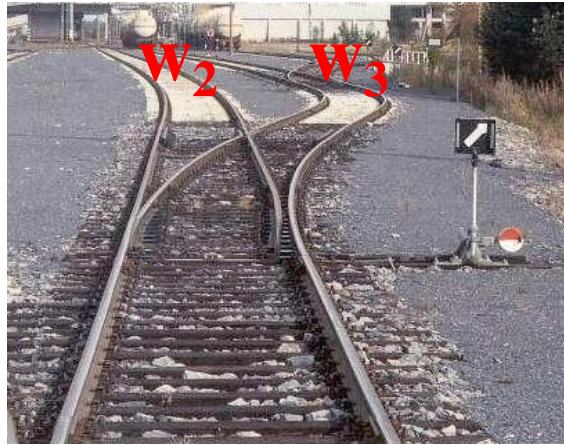
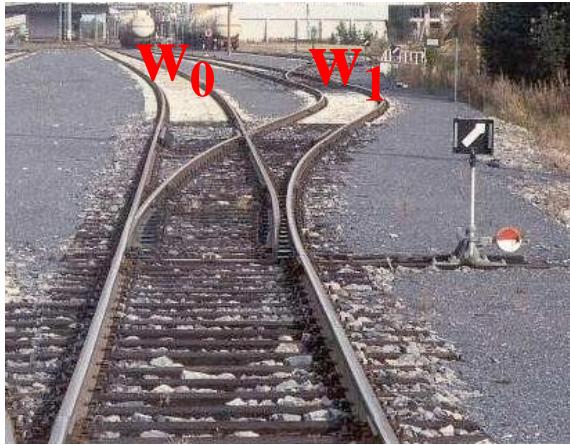
[ Figure 4.3 from the textbook ]

# Analogy: Railroad Switches



[http://en.wikipedia.org/wiki/Railroad\\_switch\]](http://en.wikipedia.org/wiki/Railroad_switch)

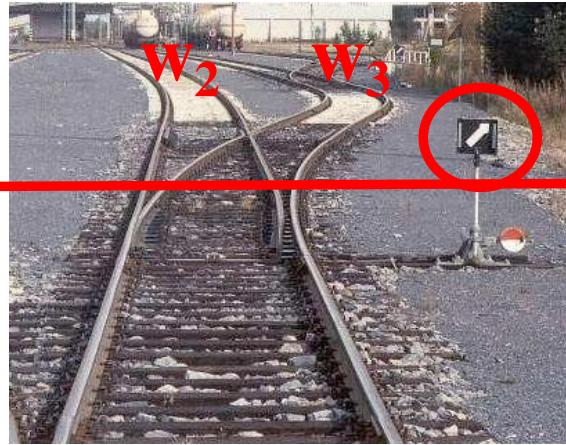
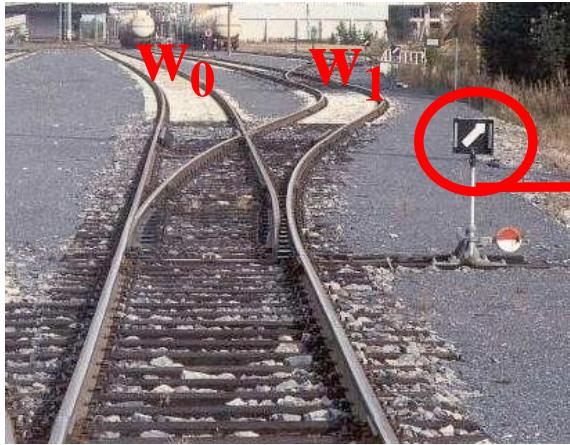
# Analogy: Railroad Switches



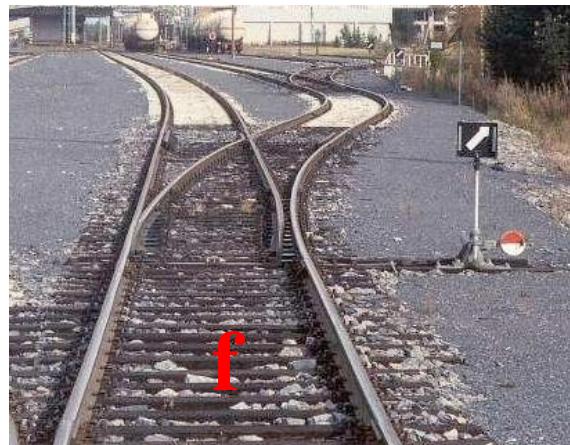
s<sub>1</sub>

[http://en.wikipedia.org/wiki/Railroad\\_switch\]](http://en.wikipedia.org/wiki/Railroad_switch)

# Analogy: Railroad Switches



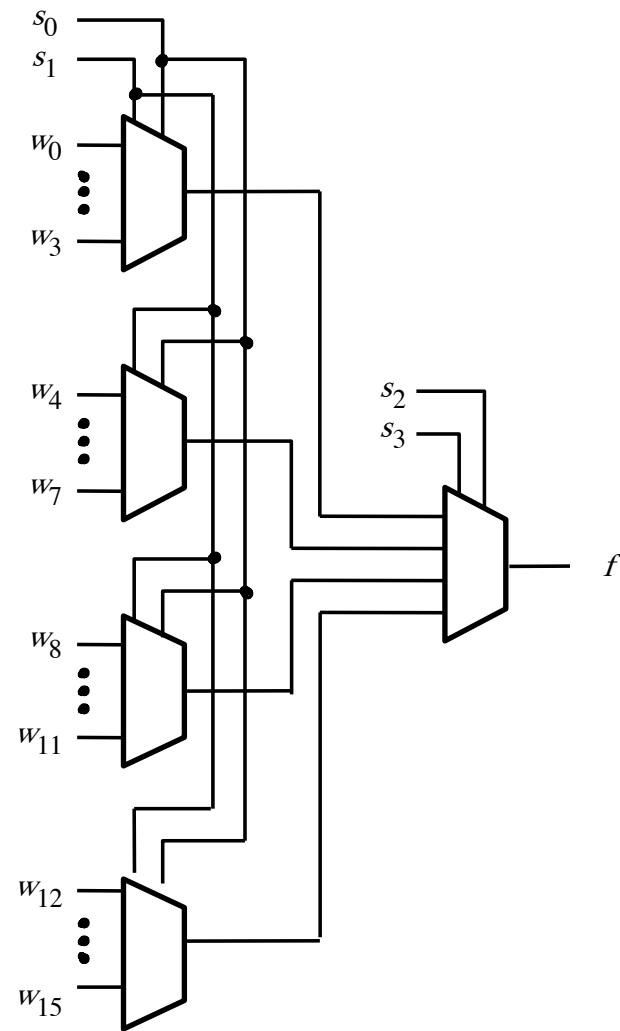
these two  
switches are  
controlled  
together



[http://en.wikipedia.org/wiki/Railroad\\_switch\]](http://en.wikipedia.org/wiki/Railroad_switch)

# **16-to-1 Multiplexer**

# 16-to-1 Multiplexer



[ Figure 4.4 from the textbook ]



[<http://upload.wikimedia.org/wikipedia/commons/2/26/SunsetTracksCrop.JPG>]

# **Questions?**

**THE END**